

Matching

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Match agents to items.

- public housing
- dorm assignment
- school choice
- course allocation
- organ transplant
- food banks

Goals:

- high value
- respect priority
- stability of self
- minimizing waste
- transparent
- easy to optimize
- minimize regret

Model

- set of N agents
- set of I items
- defn Prefs \succ_i : strict total orders over I ,
 $a \succ_i b$
- defn util $v_i: I \rightarrow [0, \infty]$
 where $v_i(a) \geq v_i(b)$ iff $a \succeq_i b$

Ex. items

	a	b	c	d
agents	100	3	2	1
2	100	10	2	1
3	100	10	18	1
4	100	99	98	17

defn An allocation x

- $x_{ia} = 1$ if i gets a
- $x_{ia} = 0$ if i doesn't get a

feasible if $x_{ia} = x_{ja} = 1$
 then $i = j$

defn An alloc x is

ex post Pareto Eff. (PE)
 if there is no alloc. y s.t.

$$\sum_{a \in I} v_{ia} y_{ia} \geq \sum_{a \in I} v_{ia} x_{ia}$$

\forall all $i \in N$, strict for some i .

Mechanisms

def. Map inputs (\succ_i) to alloc. x

def. Strategyproof if each agent's alloc. is optimized by reporting her true value.

def. PE if equal alloc. is PE

Serial Dictatorship

Pick an ordering of agents.
 For $i = 1$ to n ,

- let a be i 's favorite remaining item
- set $x_{ia} = 1$ and remove a

Claim ex post PE + strategyproof

defn A mech. has "equal-treatment-of-equals" if identical agents get the same alloc. in expectation.

Random Serial dictatorship (RSD)

- run serial dictatorship w/ random order.
- Claim Ex post PE, strategyproof, equal-treatment-of-equals

def a randomized alloc. is a set of allocs $\{x^1, \dots, x^k\}$ and a convex comb. of them $\{q_1, \dots, q_k\}$

def. A lottery is $\{P_{ia}\}$ s.t.

$$\sum_{i=1}^n P_{ia} = 1 \quad \forall a$$

$$\sum_{a \in I} P_{ia} = 1 \quad \forall i$$

$v_{ia} = 100, v_{ib} = 3, v_{ic} = 2, v_{id} = 1$

pref b w/prob. $\frac{1}{3}$ vs a w/prob. $\frac{1}{100}$?

defn (risk neutral)

util of lottery = exp. util of outcome

$$v_i, p_i \geq 0 \quad \forall i$$

$$\sum_{a \in I} p_i a = 1 \quad \forall i$$

util of lottery = exp. util of outcome

Assume risk neutrality.

Ex. PSD lottery $p_i a = 1/4$

$$E[V_i] = \sum_{a \in I} v_i a p_i \approx 25$$

defn. A lottery is ex ante PE if there is no lottery q that someone prefers and no one is harmed.

claim. ex ante PE \rightarrow ex post PE

Q.? Is PSD ex ante PE?

Ex. 8 agents, 8 items

values	a	2 copies b
4 copies of 1	100	1
4 copies of 2	100	99

RSD:

$$E[V_1^{RSD}] \approx \frac{1}{8} \times 100 = 12.5$$

$$E[V_2^{RSD}] \approx \frac{3}{8} \times 100 = 37.5$$

ALT: give a to a type 1
" b " " 2

$$E[V_1^{ALT}] \approx 25$$

$$E[V_2^{ALT}] = 49.5$$

Two-sided matching

Agents to agents.

- job markets, NRM
- school choice
- marriage markets

Goals

- PE
- no justified envy (stability)
- strategy proof

- set of men M
- set of women W
- prefs $a \succ_x b$: strict, total
- matching $\mu: M \rightarrow W$
 $\mu(m)$ is m 's match
 $\mu(w)$ is w 's match
 (m, w) if $\mu(m) = w$

defn PE: μ is PE if \nexists

ν st.

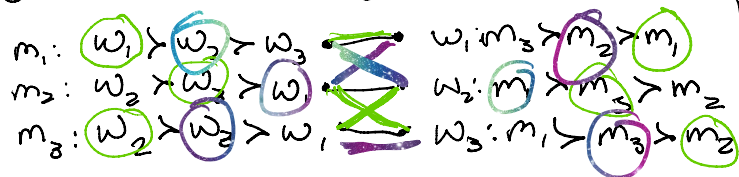
$$\nu(x) \succ_x \mu(x) \quad \forall x$$

$$\nu(x) \succ_x \mu(x) \text{ for some } x$$

defn. μ is stable if

- IR: x prefers $\mu(x)$ to being single
- no pair (m, w) st. $w \succ_m \mu(m)$ and $m \succ_w \mu(w)$

Goal: find PE stable μ



Claim Stable \Rightarrow PE

Deferred Acceptance

- Let $\mu(m) = m \quad \forall m$
- Let $S = \{m : \mu(m) = m\}$
- While $\exists m \in S$

a) m applies to favorite w who has not yet rejected him

Claim Stable \Rightarrow PE

m-DA on example

- 1) $m_1 \rightarrow w_1 : \mu = \{(m_1, w_1)\}$
- 2) $m_2 \rightarrow w_2 : \mu = \{(m_1, w_1), (m_2, w_2)\}$
- 3) $m_3 \rightarrow w_2 : \mu = \{(m_1, w_1), (m_3, w_2)\}$
- 4) $m_2 \rightarrow w_3 : \mu = \{(m_1, w_1), (m_2, w_3), (m_3, w_2)\}$

w-DA on example

- 1) $w_1 \rightarrow m_3 : \mu = \{(w_1, m_3)\}$
- 2) $w_2 \rightarrow m_1 : \mu = \{(w_1, m_3), (w_2, m_1)\}$
- 3) $w_3 \rightarrow m_3 : \mu = \{(w_2, m_1), (w_3, m_3)\}$
- 4) $w_1 \rightarrow m_2 : \mu = \{(w_1, m_2), (w_2, m_1), (w_3, m_3)\}$

- a) m applies to favorite w who has not yet rejected him
- b) let $m' = u(w)$. If $m \succ_w m'$,
 - $S := S \cup m'$
 - $u(w) = m$

Thm DA is stable.

Prf. (sketch)

- 1) DA terminate
- 2) Stability: men's options get worse
women's " better

Every person prefers outcome when they propose!

Claim: Unique man-opt / woman-pessimal SM
+ m-DA finds it (in fact, form a lattice)

Claim: m-DA is strategyproof for men
but not for women