## Topological Data Analysis - II



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## Plan

(). Yesterday:

- Motivation
- Topology
- Simplicial Complexes
- Invariants
- Homology
- Algebraic Complexes
- Today
- Geometric Complexes
- Persistent Homology
- The Persistence Algorithm
- Application to Natural Images


## Outline

- Geometric Complexes
- Voronoi Diagram
- Delaunay Triangulation
- Alpha Complex
- Witness Complex
- Summary
- Persistent Homology
- The Persistence Algorithm
- Application to Natural Images


## Recall

- Procedure
- Cover points to get approximation of underlying space
- Take nerve to get combinatorial representation
- Example: $\varepsilon$-balls around points as cover
- Idea: Use geometry of embedding space to generate cover


## Voronoi Diagram



## Voronoi Diagram

- $p \in M \subseteq \mathbb{R}^{2}$
- Voronoi cell $\mathrm{V}(\mathrm{p})$ : closest points to p in $\mathbb{R}^{2}$

$$
V(p)=\left\{x \in \mathbb{R}^{2} \mid \mathrm{d}(x, p) \leq \mathrm{d}(x, y), \forall y \in M\right\}
$$



- Voronoi Diagram: Decomposition of $\mathbb{R}^{2}$ into Voronoi cells
- Voronoi (1868-1908)
- Idea: Use Voronoi cells as cover!



## Delaunay Triangulation



## Delaunay Triangulation

- Delaunay Triangulation: nerve of Voronoi cover
- Computational Geometry
- General position assumption
- no events with probability o
- no $k+1$ points on $(k-1)$-sphere
- has to be handled in practice
- Fast algorithms for $\mathbb{R}^{3}$
- Delaunay (1890-1980)



## Restricted Voronoi



## Alpha Complex



## Delaunay Subcomplex



## Alpha Complex

- Alpha cell: $A_{\varepsilon}(p)=B_{\varepsilon}(p) \cap V(p)$
- Alpha shape: union of alpha cells
- Alpha complex: nerve of alpha shape

$$
A_{\epsilon}(M)=\left\{\operatorname{conv} T \mid T \subseteq M, \bigcap_{p \in T} A_{\epsilon}(p) \neq \emptyset\right\}
$$



- Let D be the Delaunay triangulation
- $\mathrm{A}_{\mathrm{o}}=\emptyset$
- $\mathrm{A}_{\varepsilon} \subseteq \mathrm{D}$
- $\mathrm{A}_{\infty}=\mathrm{D}$
- $\mathrm{A}_{\varepsilon} \simeq \mathrm{C}_{\varepsilon}$
- [Edelsbrunner, Kirkpatrick, and Seidel '83], et al.


## Strong Witness



## Strong Witness

- Given: Point set $M \in \mathbb{R}^{\text {d }}$
- Strong witness: $\mathrm{x} \in \mathbb{R}^{\mathrm{d}}$
- $x$ is equidistant from $v_{0}, \ldots, v_{k} \in M$
- $x$ has no closer neighbor in $M$
- x witnesses k -simplex $\left\{\mathrm{v}_{\mathrm{o}}, \ldots, \mathrm{v}_{\mathrm{k}}\right\}$
- Idea: Sample for witnesses
- Problem: Prob(strong witness) =o for discrete set $M$


## Weak Witness

- Weak witness: $x \in \mathbb{R}^{d}$
- $\left|x-v_{i}\right| \leq|x-v|$ for $i=0, \ldots, k$ and $v \in M \backslash\left\{v_{o}, \ldots, v_{k}\right\}$
- x 's closest $\mathrm{k}+1$ neighbors are $\mathrm{v}_{\mathrm{o}}, \ldots, \mathrm{v}_{\mathrm{k}}$
- x witnesses k -simplex $\left\{\mathrm{v}_{\mathrm{o}}, \ldots, \mathrm{v}_{\mathrm{k}}\right\}$ weakly
- Strong witness $\Rightarrow$ weak witness
- (Theorem [de Silva])

A simplex has a strong witness iff all its faces have weak witnesses.

## Isomap



- We want to capture the underlying space, not the embedding space
- Idea: Restrict witnesses to given points M


## Witness Complex

- Given: N points M
- Choose: n landmarks L
- M \L will act as witnesses

- $\mathrm{D}=\mathrm{n} \times \mathrm{N}$, distance matrix
- Construct $\varepsilon$-graph on L: Edge $[a b] \in \mathrm{W}_{\varepsilon}(\mathrm{M})$ iff there exists a witness with $\max (\mathrm{D}(\mathrm{a}, \mathrm{i}), \mathrm{D}(\mathrm{b}, \mathrm{i})) \leq \varepsilon$
- Do Vietoris-Rips Expansion


## Complex Summary

| Complex | Name | Idea | Scales? | Extends? |
| :--- | :---: | :--- | :---: | :---: |
| Cech | $\mathrm{C}_{\varepsilon}$ | Nerve of $\varepsilon$-balls | 1 K | $\sim$ |
| Vietoris-Rips | $\mathrm{V}_{\varepsilon}$ | Pairwise dist $<\varepsilon$ | 1 K | Y |
| Alpha | $\mathrm{A}_{\varepsilon}$ | Nerve of restricted Voronoi | 500 K | $\mathrm{~d} \leq 3$ |
| Witness | $\mathrm{W}_{\varepsilon}$ | Landmarks and witnesses | 1 K | Y |

- Conformal Alpha - no global scale parameter
- Flow - stable manifolds of distance function
- Cubical - rasterize, usually interpretation of images


## Outline

(-) Geometric Complexes

- Persistent Homology
- Filtrations
- Algebraic Result
- Simple Examples
- The Persistence Algorithm
- Application to Natural Images


## The Question of Scale



## Filtration

- A filtration of a space $X$ is a nested sequence of subspaces:

$$
\emptyset=X^{0} \subseteq X^{1} \subseteq \cdots \subseteq X^{l} \subseteq \cdots \subseteq X^{m}=X
$$

- $\mathrm{C}_{\varepsilon} \subseteq \mathrm{C}_{\varepsilon^{\prime}}$ if $\varepsilon \leq \varepsilon^{\prime} \quad$ (Also true for $\mathrm{V}_{\varepsilon}, \mathrm{A}_{\varepsilon}$, and $\mathrm{W}_{\varepsilon}$ )
- Simplices are always added, never removed
- Implies partial order on simplices
- Full order: sequence of simplices
- $\mathrm{K}_{\mathrm{i}}=$ union of first i simplices in sequence


Witness Complex

## Inductive Systems



Idea: Follow basis elements from birth to death
Problem: Need a compatible basis!

## Persistent Homology



- Persistence barcode: multiset of intervals



## Algebraic Result

1. Correspondence

- Input: Filtration
- Structure of homology: graded $k[t]$-module

2. Classification

- $k$, a field $\Rightarrow k[t]$ is a PID
- Structure theorem for graded PIDs

3. Parameterization

- $n$ half-infinite
- $m$ finite

$$
\left(\bigoplus_{i=1}^{n} \Sigma^{\alpha_{i}} k[t]\right) \oplus\left(\bigoplus_{j=1}^{m} \Sigma^{\gamma_{j}} k[t] /\left(t^{n_{j}}\right)\right)
$$

- Barcode: multiset of $n+m$ intervals (birth, death)
- Complete discrete invariant!


## Deconstructing the Graph (2D)



Torus!


## Discovering 3D Structure



## Outline

© Geometric Complexes
() Persistent Homology

- The Persistence Algorithm
- Adding a Simplex
- Example Filtration
- Application to Natural Images


## Adding a Simplex

- Given: Filtered complex K
- $\mathrm{K}_{\mathrm{i}}=\mathrm{K}_{\mathrm{i}-1} \cup \sigma$, where $\sigma$ is a k -simplex
- Let $\mathrm{c}=\partial \sigma$. c is $\mathrm{a}(\mathrm{k}-1)$-chain.
- (Lemma) c is a cycle.
- Proof: $\partial c=\partial \partial \sigma=0$.
- (Lemma) c is in $\mathrm{K}_{\mathrm{i}-1}$.
- Proof: $K_{i}$ is a simplicial complex.


## Gaussian Elimination

- $\sigma$ is a $k$-simplex
- $\mathrm{c}=\partial \sigma$ is a $(\mathrm{k}-1)$-cycle in $\mathrm{K}_{\mathrm{i}-1}$
- Two cases: c is a boundary or not in $\mathrm{K}_{\mathrm{i}-1}$
- $M_{k}$ is matrix for $\partial_{k}$
- c is a boundary iff
- it is in range $\left(M_{k}\right)$
- we can write it in terms of a basis for $M_{k}$
- Gaussian elimination maintains a basis for range $\left(M_{k}\right)$
- Filtration and persistence imply ordering on pivots


## Case 1: c is a boundary in $\mathrm{K}_{\mathrm{i}-1}$

- If c is a boundary, then $\exists \mathrm{d} \in \mathrm{C}_{\mathrm{k}+1}\left(\mathrm{~K}_{\mathrm{i}-1}\right)$, such that $\mathrm{c}=\partial \mathrm{d}$
- (Lemma) $\sigma+\mathrm{d}$ is a k -cycle in $\mathrm{K}_{\mathrm{i}}$.

- Proof: $\partial(\sigma+\mathrm{d})=\partial \sigma+\partial \mathrm{d}=\mathrm{c}+\partial \mathrm{d}=0$.
- $\sigma$ creates a new k-cycle class
- $\sigma$ is a creator


## Case 2: c is not a boundary in $\mathrm{K}_{\mathrm{i}-1}$

- (Lemma) c becomes a boundary in $\mathrm{K}_{\mathrm{i}}$.
- Proof: c= $=\partial \sigma$.
- $\ln \mathrm{K}_{\mathrm{i}-1}$
- c is a cycle

- c is not a boundary
- c is in a non-boundary homology class
- In $\mathrm{K}_{\mathrm{i}}$ : c is a boundary, so its homology class is trivial.
- $\sigma$ destroys a ( $\mathrm{k}-1$ )-dimensional class
- $\sigma$ is a destroyer
- Suppose $\tau$ created that class that $\sigma$ destroyed
- We pair $(\tau, \sigma)$ to get the lifetime interval


## Example



## Filtration

- Initially, cascade $=\sigma_{i}$

cascade
(cascade)


## Vertices a, b, c, d

- $\partial \sigma=o$ for all vertices $\sigma$



## ab

- We sort $\partial \mathrm{ab}=\mathbf{b}+\mathbf{a}$ by youngest
- Since $b$ is unpaired, pair with $a b$



## bc, cd

- $\partial \mathrm{bc}=\mathbf{c}+\mathrm{b}$
- $\partial c d=\mathbf{d}+c$



## ad

- $\quad$ $\quad$ ad $=(d+a) \sim(d+a)+(d+c)=c+a$

$$
\sim(c+a)+(c+b)=b+a \sim(b+a)+(b+a)=0
$$

| filtration |  | b |  | c | d |  | ab | bc | cd | ad | ac | abc | acd |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  |  | a |  | bc | c |  | b | c | d |  |  |  |  |



| $\partial$ (cascade) | o | o | 0 | $\circ$ | $b$ | $c$ | $d$ | $\circ$ |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
|  |  |  |  |  | $a$ | $b$ | $c$ |  |

## ac

- $\partial \mathrm{ac}=(\mathrm{c}+\mathrm{a}) \sim(\mathrm{c}+\mathrm{a})+(\mathrm{c}+\mathrm{b})=\mathrm{b}+\mathrm{a}$
$\sim(b+a)+(b+a)=0$

| filtration |  | b |  | c | d | a | b | bc | cd | ad | ac | abc | acd |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  |  | a |  | bc | cd | b | b | c | d |  |  |  |  |


| cascade | a | $b$ | $c$ | $d$ | $a b$ | $b c$ | $c d$ | $a d$ | $a c$ |  |
| :---: | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
|  |  |  |  |  |  |  |  | $c d$ | $b c$ |  |
|  |  |  |  |  |  |  |  |  | $b c$ | $a b$ |
|  |  |  |  |  |  |  |  | $a b$ |  |  |


| $\partial$ (cascade) | o | o | o | 0 | $b$ | $c$ | $d$ | $o$ | 0 |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
|  |  |  |  |  | $a$ | $b$ | $c$ |  |  |



## abc

- $\partial \mathrm{abc}=\mathrm{ac}+\mathrm{bc}+\mathrm{ab}$




## acd

- $\quad$ acd $=\mathrm{ac}+\mathrm{ad}+\mathrm{cd} \sim$

$$
(a c+a d+c d)+(a c+b c+a b)=a d+c d+b c+a b
$$

| filtration |  |  | b | c |  | d | ab | bc | cd | ad | ac | abc | acd |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  |  |  | ab | bc |  | cd | b | c | d | acd | abc | ac | ac | ac |




$$
\begin{array}{|c|cccccccccccc|}
\hline \partial \text { (cascade) } & 0 & 0 & 0 & 0 & b & c & d & 0 & 0 & \text { ac } & \text { ad } \\
\hline & & & & & a & b & c & & & & b c & c d \\
& & & & & & & & & & a b & b c \\
& & & & & & & & & & & & a b
\end{array}
$$

## Barcode

| $a_{0} b^{b}$ |  |  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 0 a, b | $1{ }^{\text {a }}$ ab, ${ }^{\text {c, }}$ | $2 \mathrm{~cd}, \mathrm{ad}$ | 3 ac | 4 abc | 5 | acd |

- $\beta_{\mathrm{o}}$ : a is unpaired $\Rightarrow[0, \infty)$
- $\beta_{o}:(b, a b) \Rightarrow[0,1)$
- $\beta_{o}:(c, b c) \Rightarrow \emptyset$
- $\beta_{o}:(d, c d) \Rightarrow[1,2)$
- $\beta_{1}:(\mathrm{ad}, \mathrm{acd}) \Rightarrow[2,5)$
- $\beta_{1}:(\mathrm{ac}, \mathrm{abc}) \Rightarrow[4,5)$


## Outline

© Geometric Complexes
() Persistent Homology
() The Persistence Algorithm

- Application to Natural Images


## Natural Images


J. H. van Hateren, Neurobiophysics, U. Groningen

## Local Structure: $3 \times 3$ Patches


$(0.81,0.62,0.64,0.82,0.65,0.64,0.83,0.66,0.65) \in \mathbb{R}^{9}$

## Mumford Dataset

- David Mumford (Brown)
- $3 \times 3$ patches $\left(\mathbb{R}^{9}\right)$
- Subtract mean intensity $\left(\mathbb{R}^{8}\right)$
- Remove low contrast patches
- Rescale to unit length $(\mathbb{S})$
- 2.5 million points on $\mathbb{S} 7$
- What is its structure?
- Examine dense areas


## Space of Idealized Lines

- Lines in natural images
- Rasterized in $3 \times 3$ patches

- Parameterization
- Distance to center: $\mathbb{I}$
- Angle: $\mathbb{S}^{1}$
- Space is annulus: $\mathbb{I} \times \mathbb{S}^{1}$


Demo

## Graph Structure



2D Structure


## The Klein Bottle



- Can we design a compression algorithm that uses the Klein bottle?


## Software

- PLEX: comptop.stanford.edu/programs/plex
- Cech
- Vietoris-Rips
- Witness
- Persistence
- CgAL: www.cgal.org
- Alpha
- Persistence (?)
- CHomP: chomp.rutgers.edu
- Alpha Shapes: biogeometry.duke.edu/software/alphashapes
- GGobi: www.ggobi.org


## Conclusion

- We are flooded by point set data and need to find structure in them
- Topology studies connectivity of spaces
- Topological analysis may be viewed as generalization of clustering
- To analyze point sets, we require a combinatorial representation approximating the original space
- Homology focuses on the structure of cycles
- Persistent homology analyzes the relationship of structures at multiple scales

