Efficient TopK-Query Algorithm for Minerva

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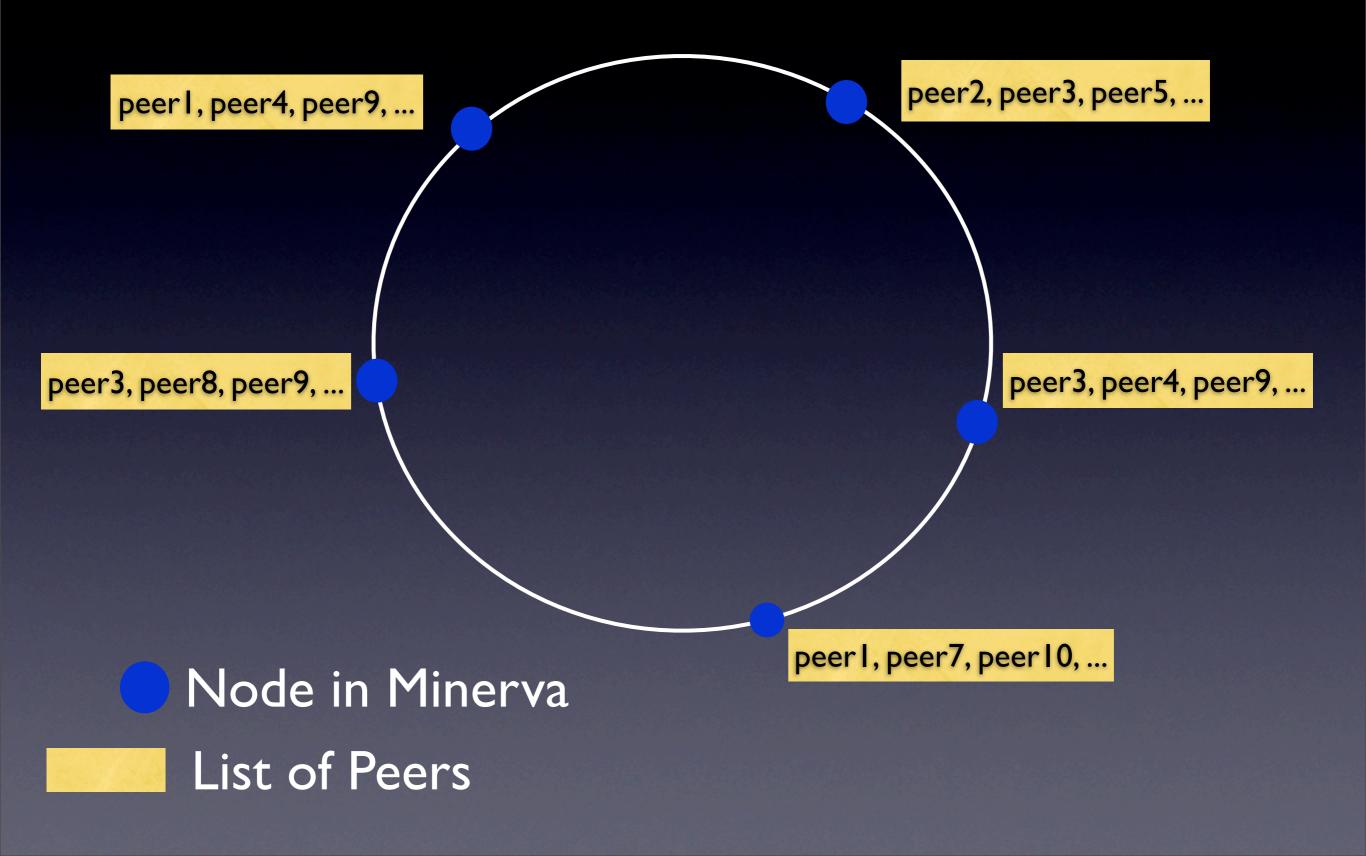
Motivation

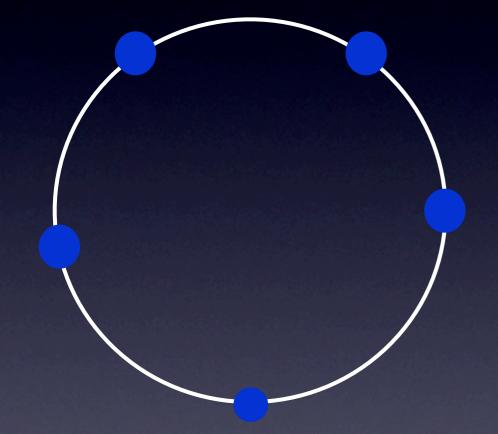
- New search technique: P2P Web Search
- Implementation of technique in Minerva
- Allows to search for multiple terms
- Problem of overlap in selection of peers
- Efficient selection of peers

Minerva Design

Node in Minerva

Minerva Design





"Britney"

"Britney"

"Spears"

"Britney"



"Britney"

"Britney Spears"



"Britney"

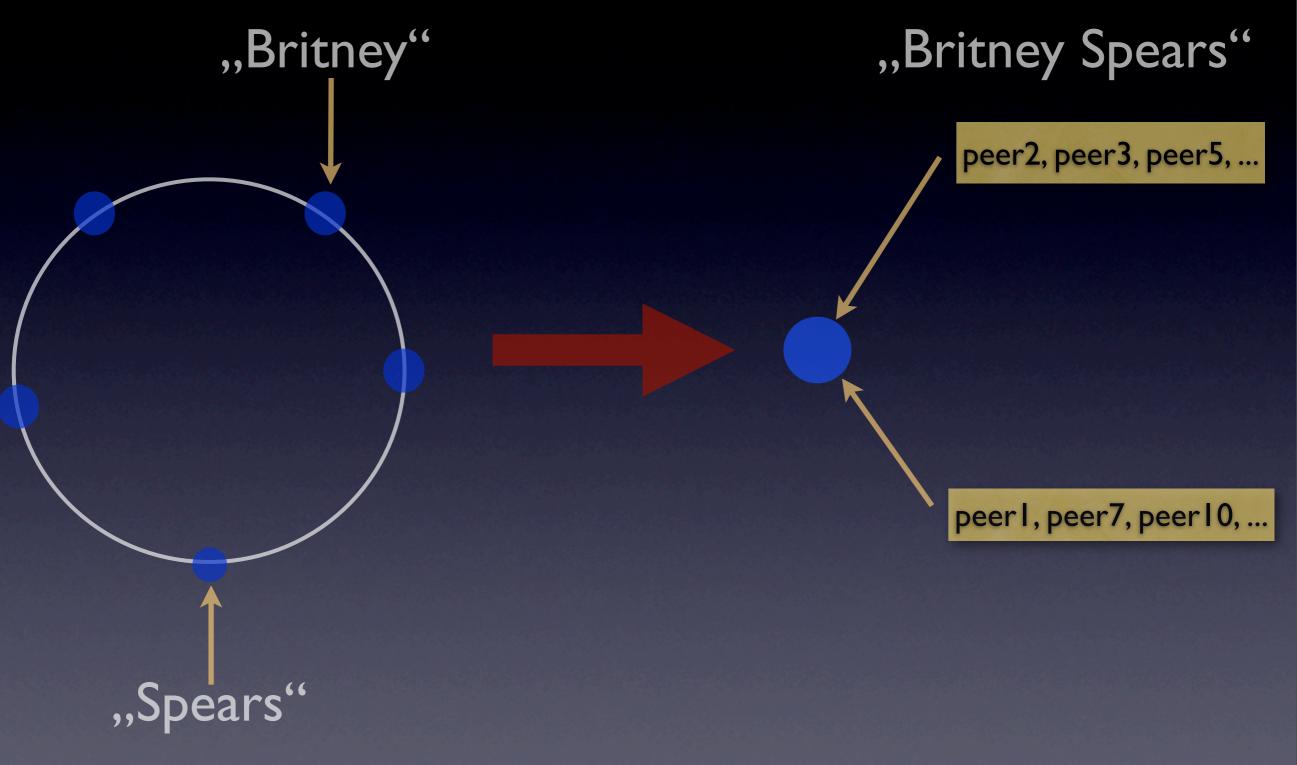
"Britney Spears"

peer2, peer3, peer5, ...

peerl, peer7, peerl0, ...

"Spears"

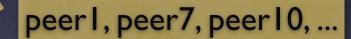






"Britney"



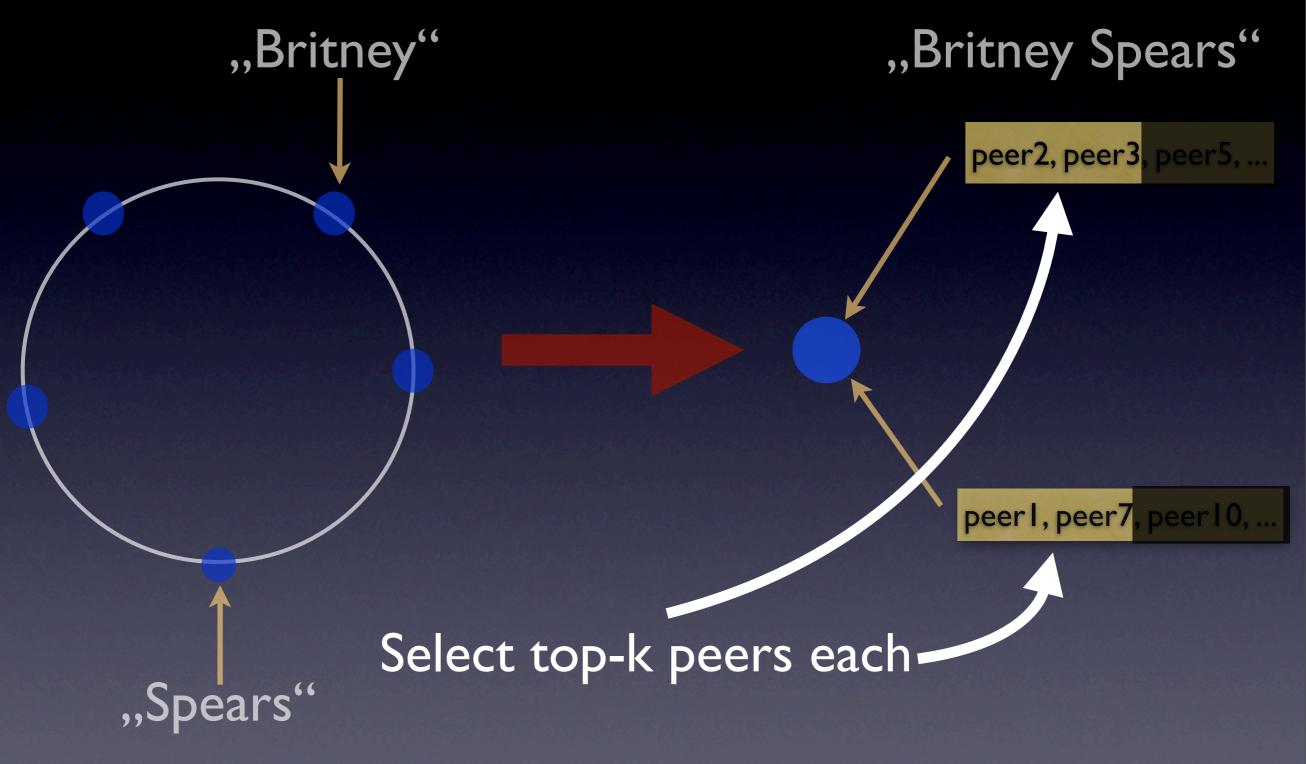


peer2, peer3, peer5, ...



Select top-k peers each-

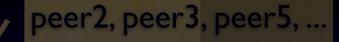


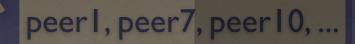




"Britney"









"Britney"

"Britney Spears"

3, peer5, ...

7, peer 10, ...

Fetching all peers from nodes is too expensive !!

"Spears"

"Britney"

Idea:

Usage of Distributed Top-K algorithm (TPUT, Klee, Fagin)

"Britney Spears"

3, peer5, ...

, peer 10, ...



TPUT

- Peerlists ranked in descending order
- Phase I:
 - Select from all peerlists top-k peers
 - Compute aggregated value for every peer
 - k'th top value is T₁ (phase-1 bottom)
- <u>Phase 2</u>:
 - Set threshold $T = (\tau_1 / m)$, m is number of nodes
 - Select k highest peers $\geq T$
 - k'th highest is T_2 (phase-2 bottom, where $T_1 \leq T_2 \leq T$)
 - Peers $< \tau_2$ are eliminated, remaining ones in set S

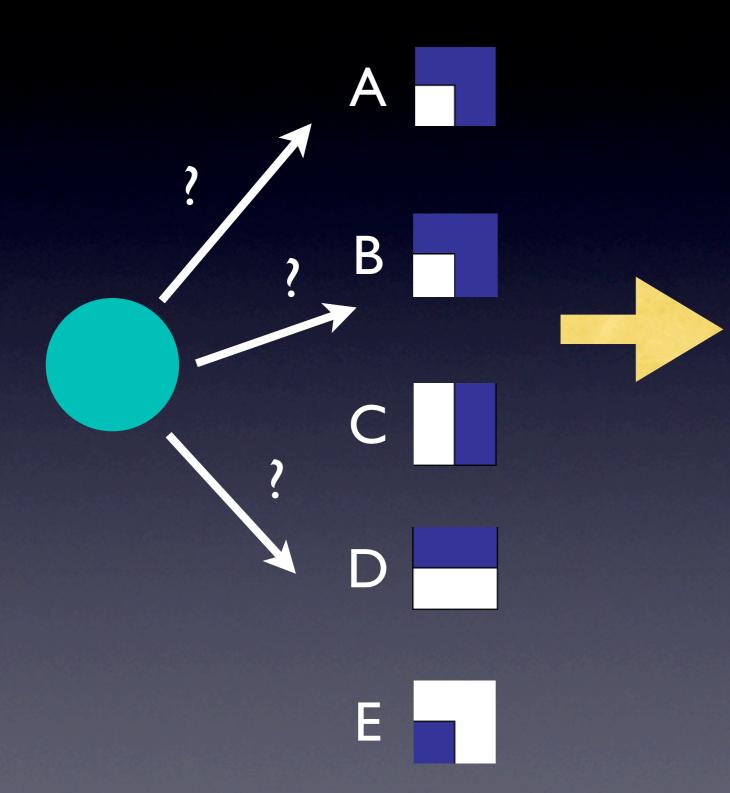
TPUT

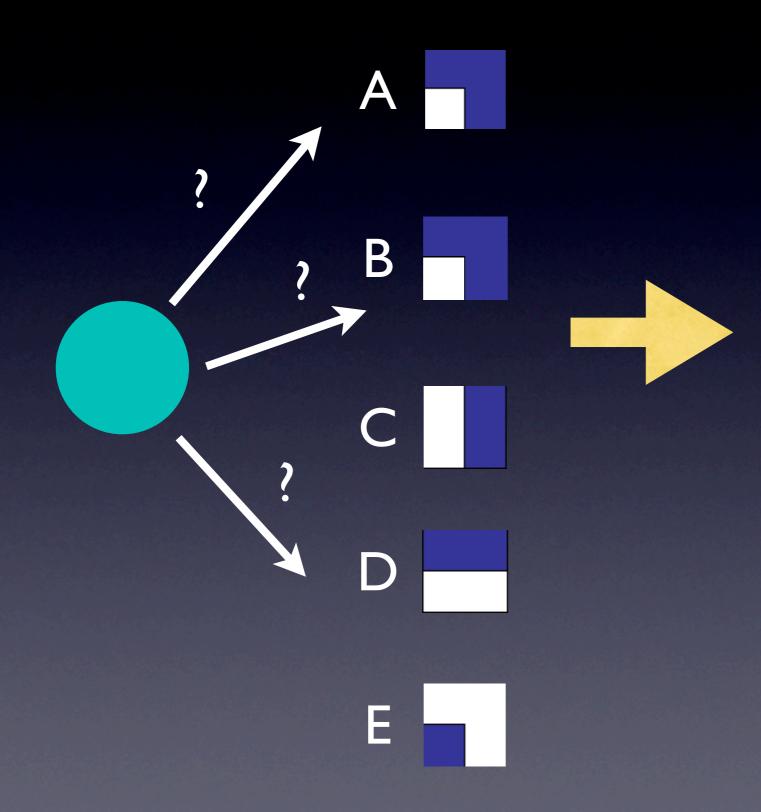
• <u>Phase 3:</u>

- Set S sended to all nodes
- Each node sends values of peers back
- Calculate exact sum of peers in S
- Choose top-k peers
- True top-k peers

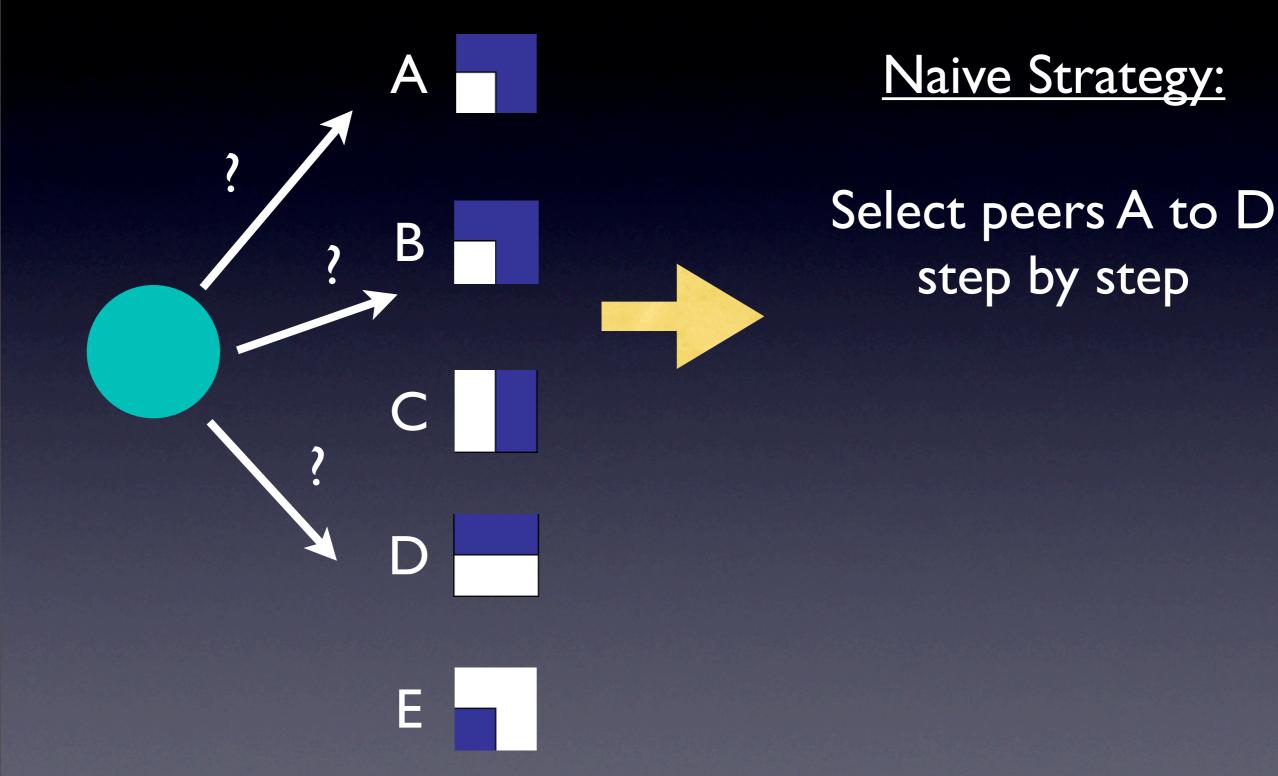
TPUT

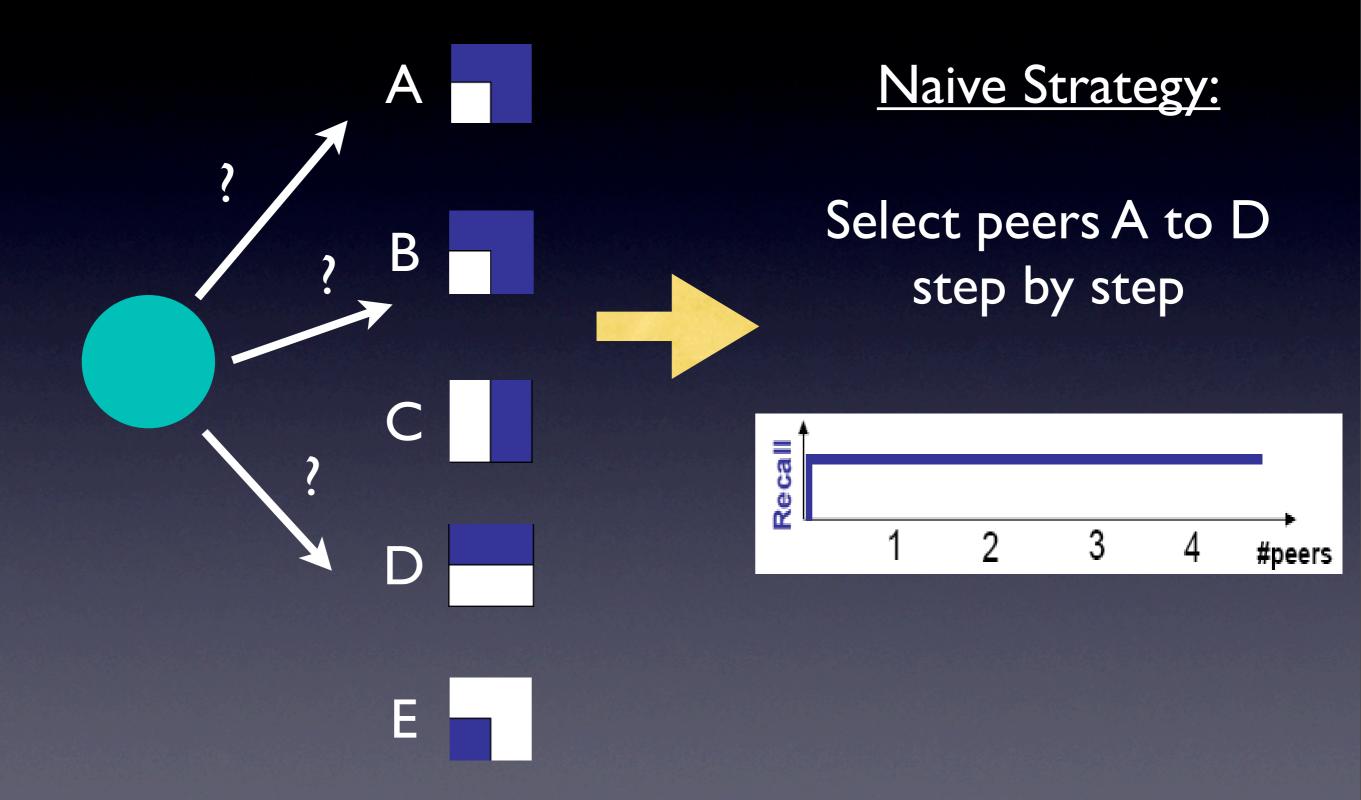
- Algorithm computes best k peers (for constant ranking /scores)
- <u>Now:</u> Adapt algorithm for Minerva peer selection
- Problem of Overlap Aware
- Goal: Analyze how peer selection in Minerva influences the result of Top-K

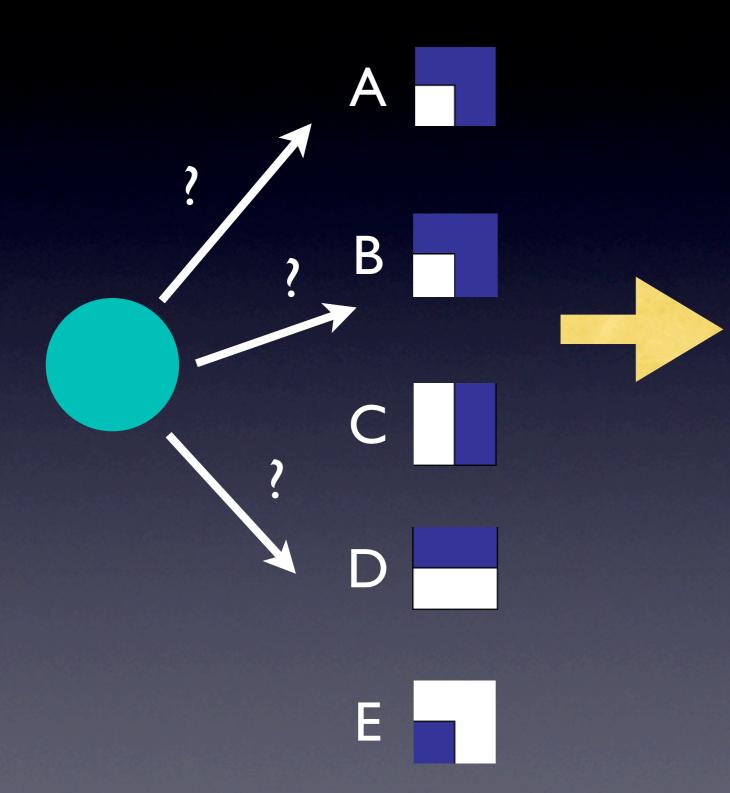


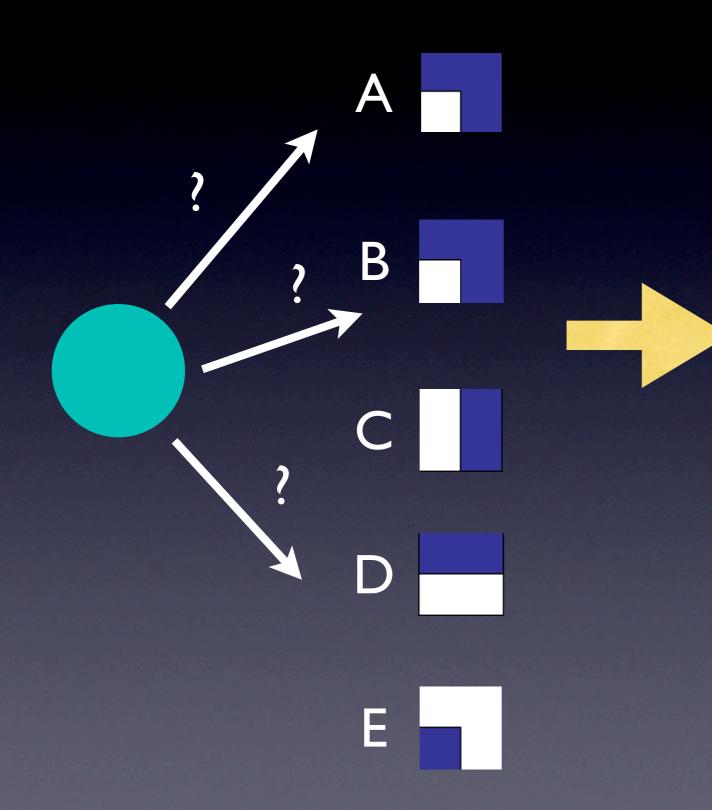


Naive Strategy:

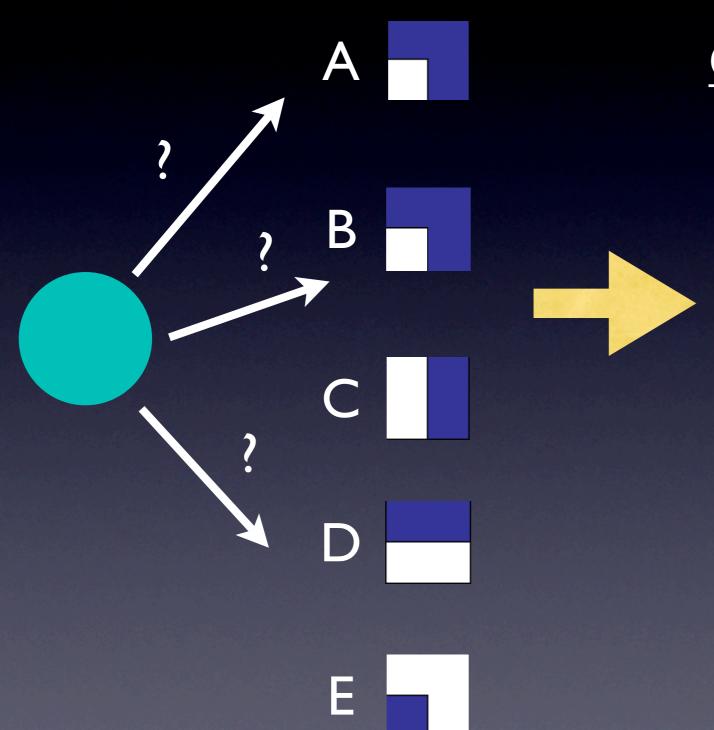






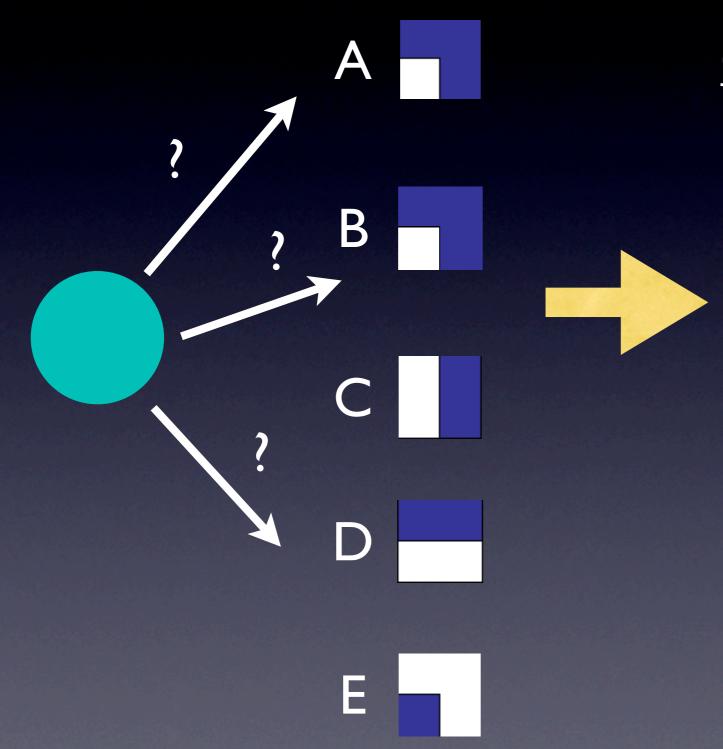


Overlap Aware Strategy:



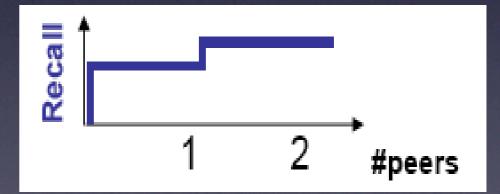
Overlap Aware Strategy:

Combine A and E for max. Recall by min. of peers



Overlap Aware Strategy:

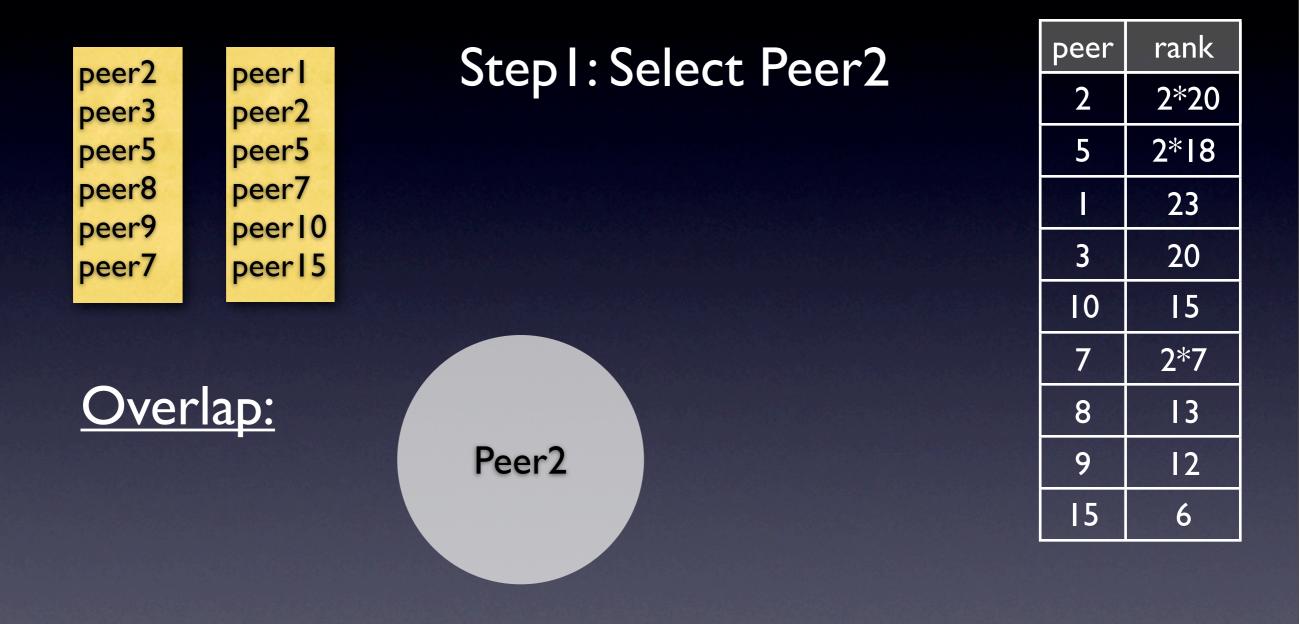
Combine A and E for max. Recall by min. of peers

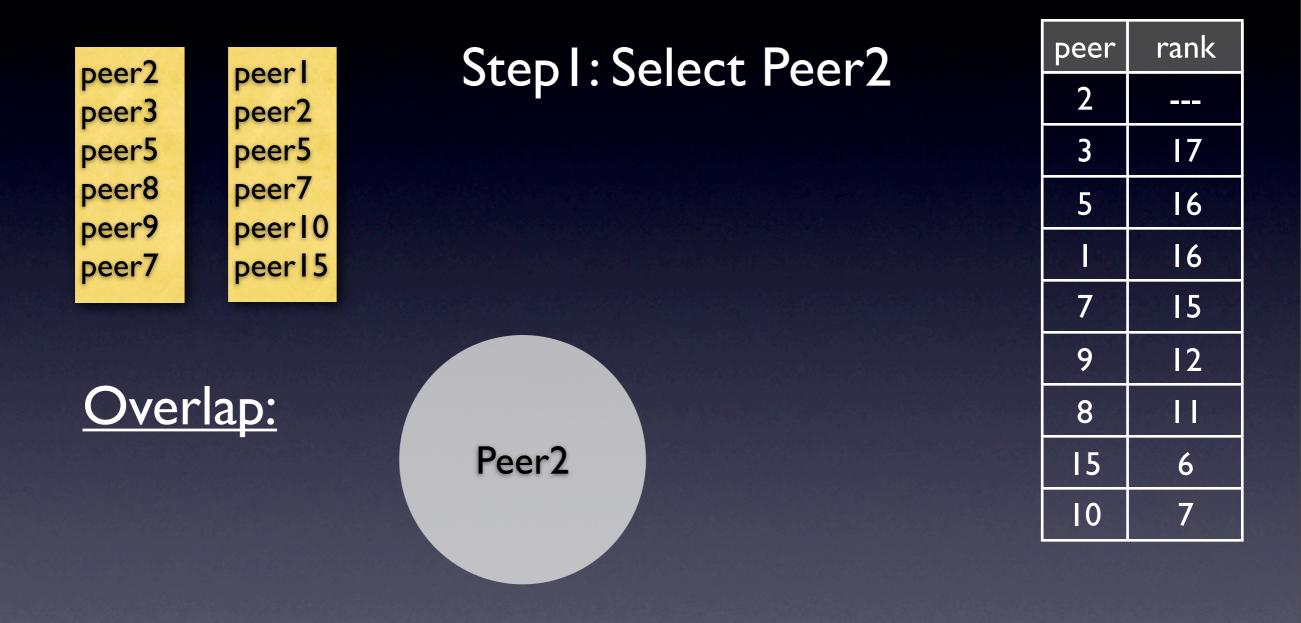


peer2	peerl
peer3	peer2
peer5	peer5
peer8	peer7
peer9	peer10
peer7	peer 15

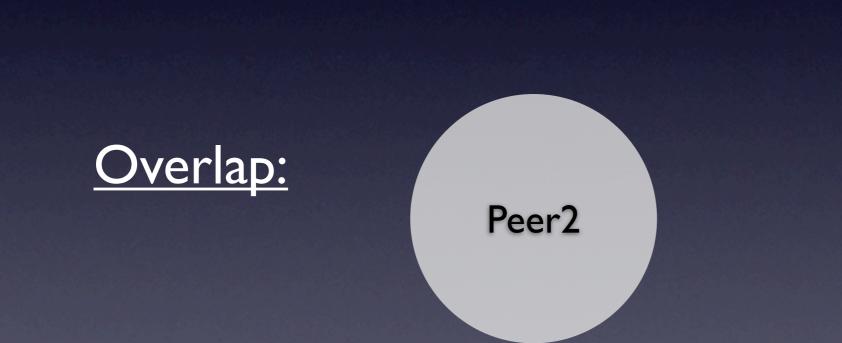


peer	rank
2	2*20
5	2*18
I	23
3	20
10	15
7	2*7
8	13
9	12
15	6

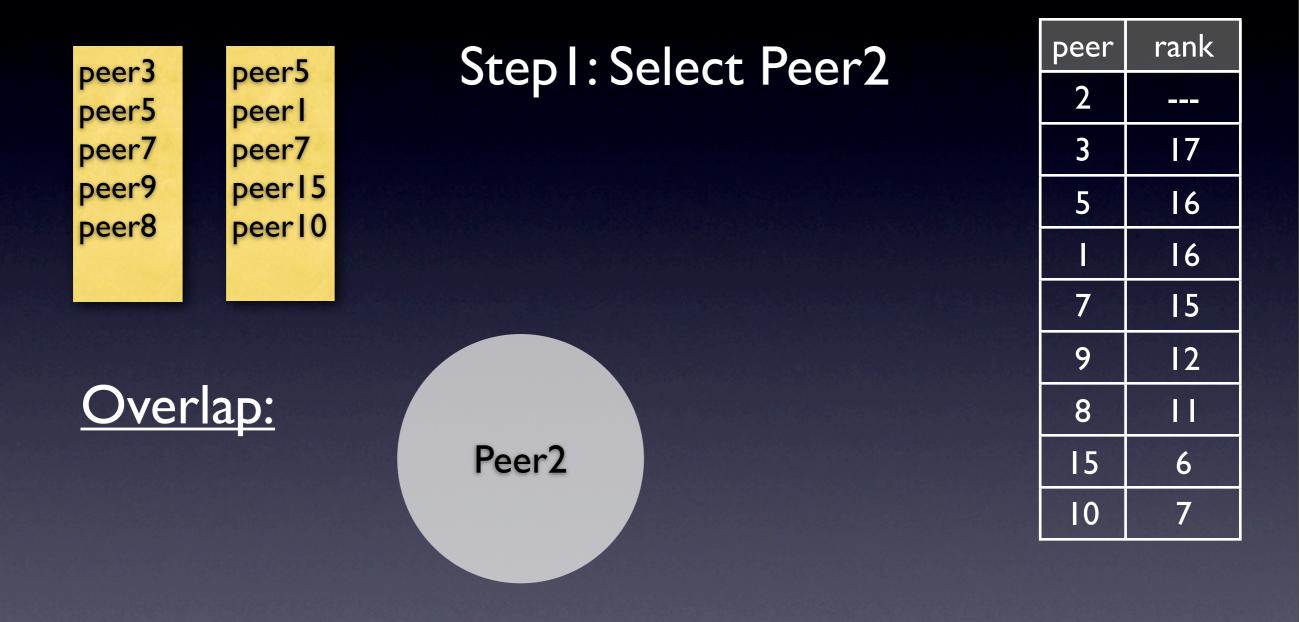


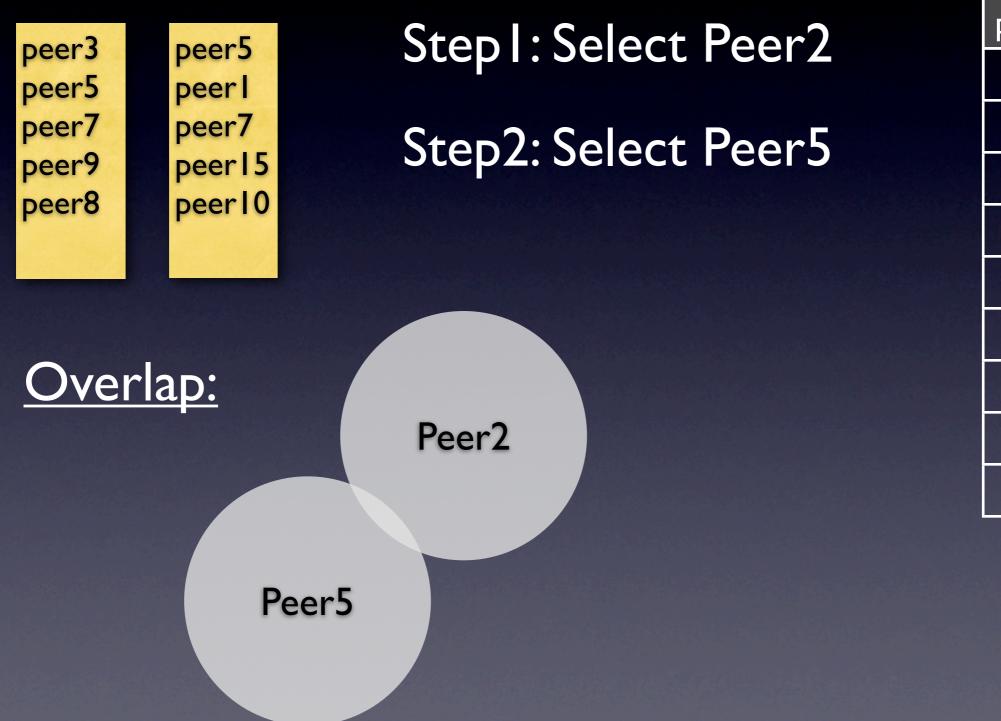


Step I: Select Peer2

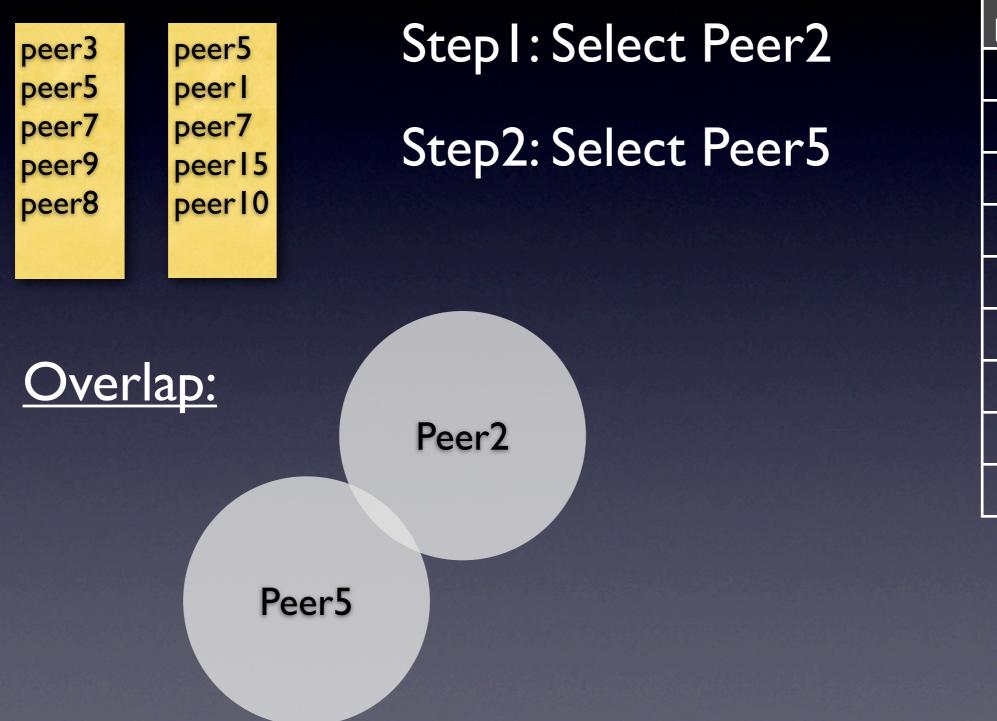


peer	rank
2	
3	17
5	16
	16
7	15
9	12
8	11
15	6
10	7





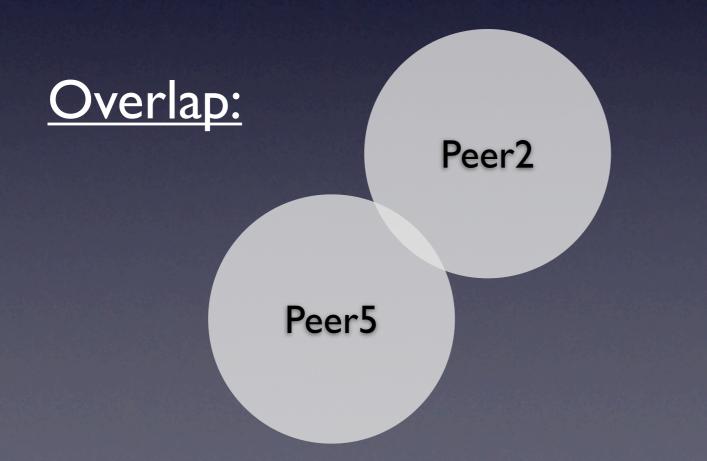
peer	rank
2	
3	17
5	16
	16
7	15
9	12
8	11
15	6
10	7



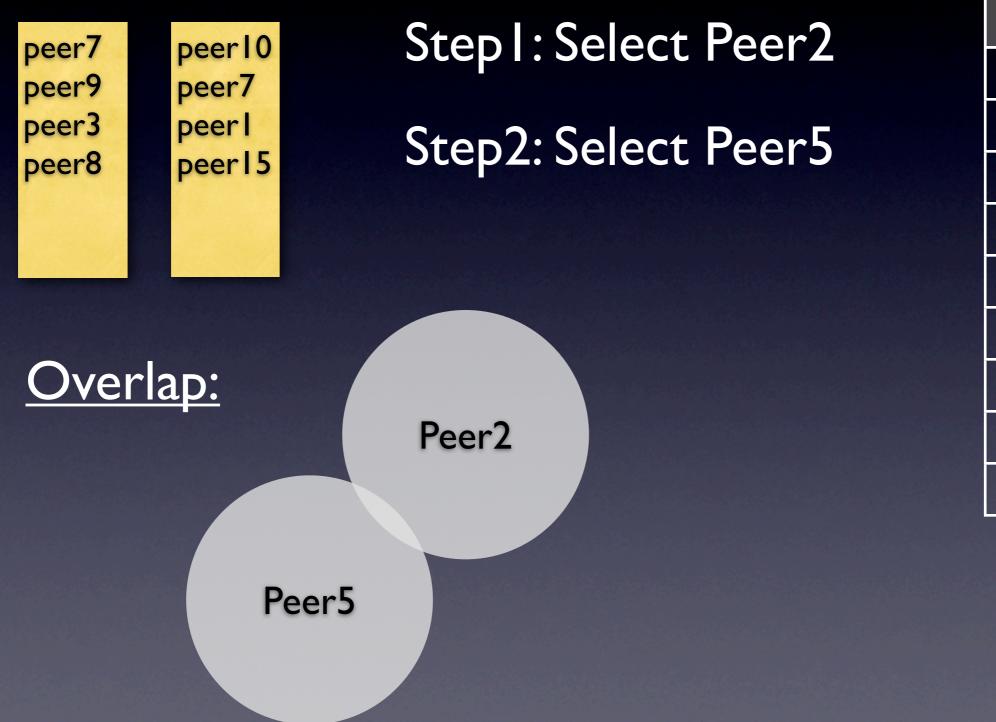
peer	rank
2	
5	
10	14
7	12
I	12
9	12
3	11
8	6
15	7

Step I: Select Peer2

Step2: Select Peer5

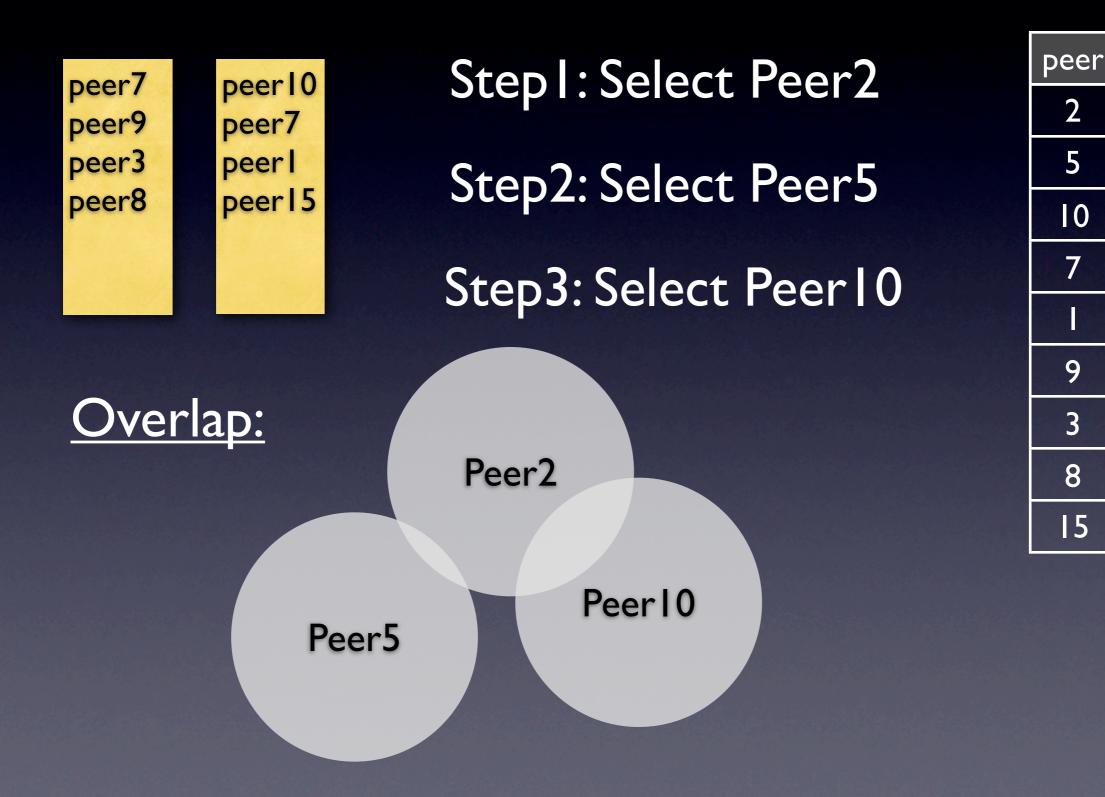


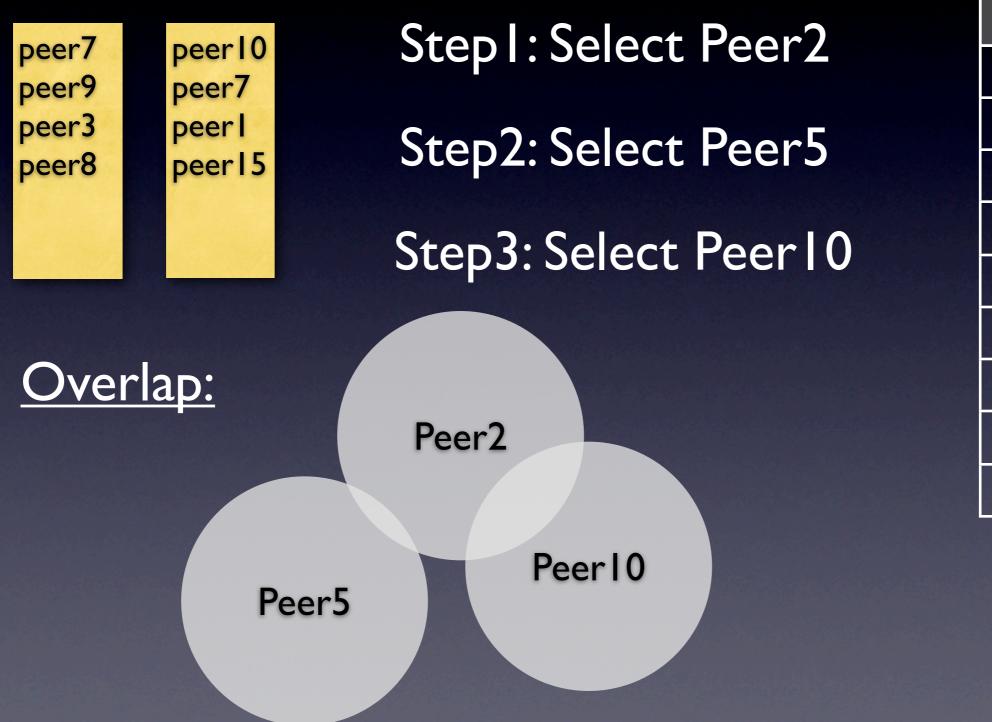
peer	rank
2	
5	
10	14
7	12
Ι	12
9	12
3	11
8	6
15	7



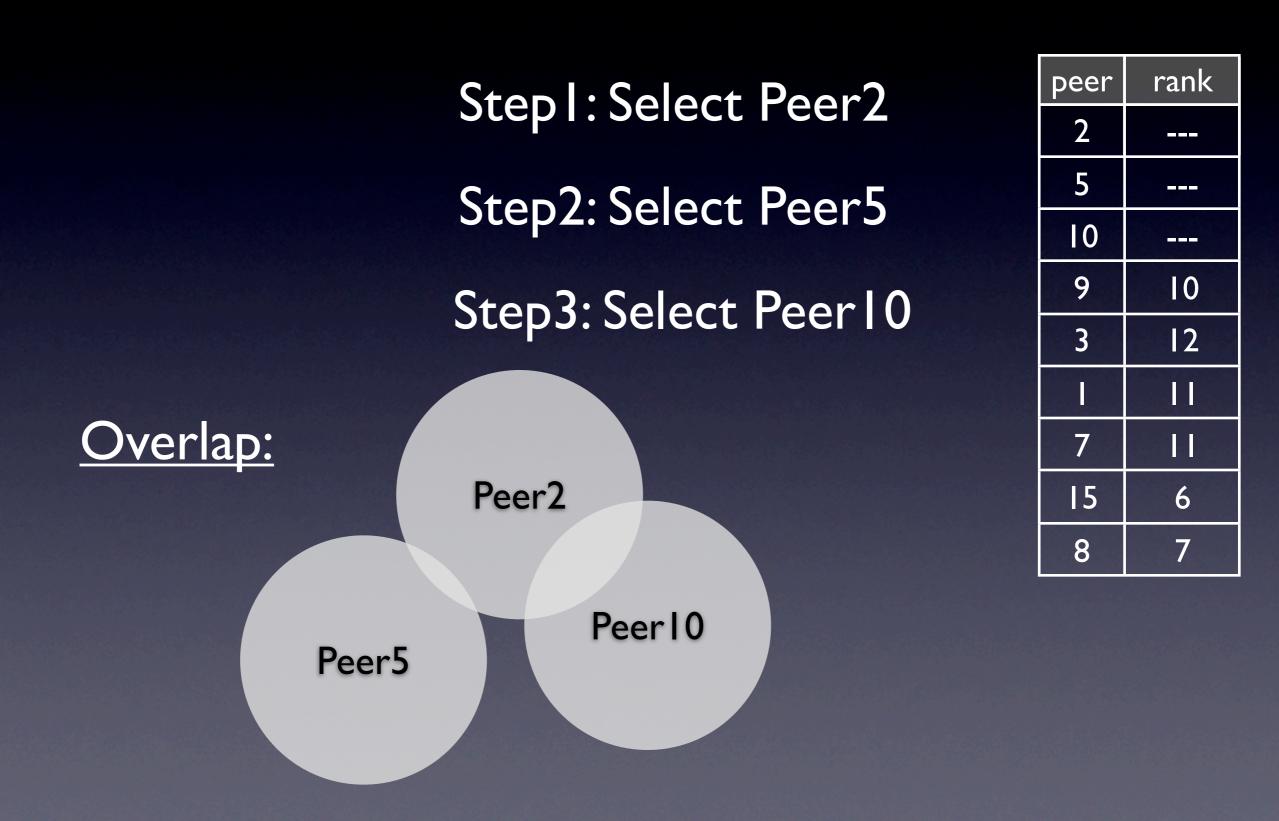
peer	rank
2	
5	
10	14
7	12
Ι	12
9	12
3	11
8	6
15	7

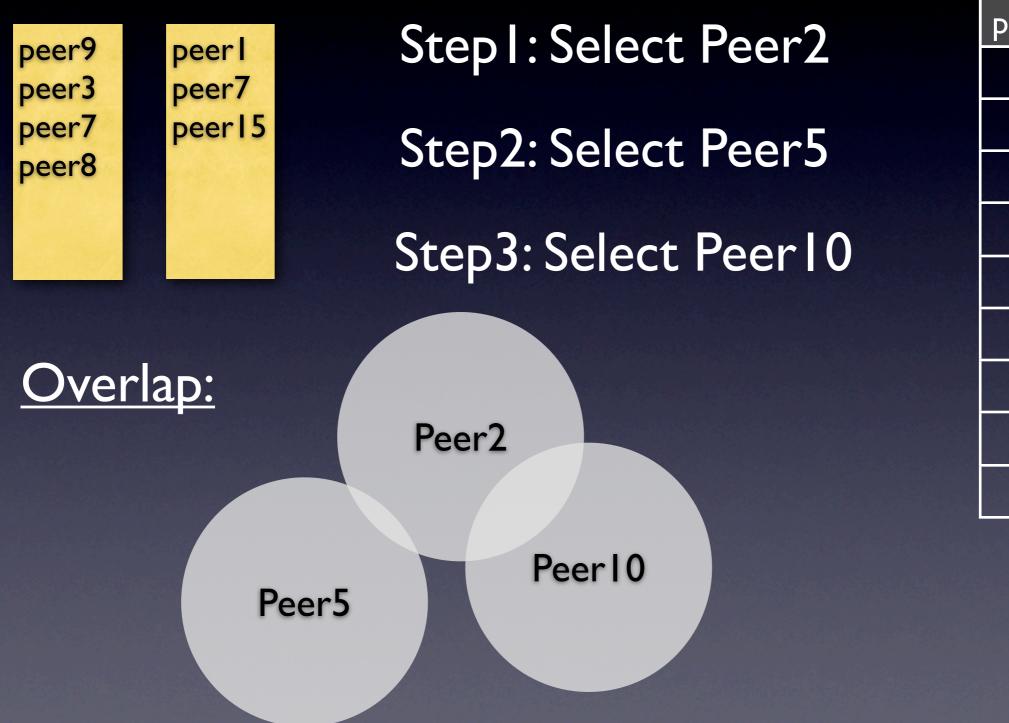
rank



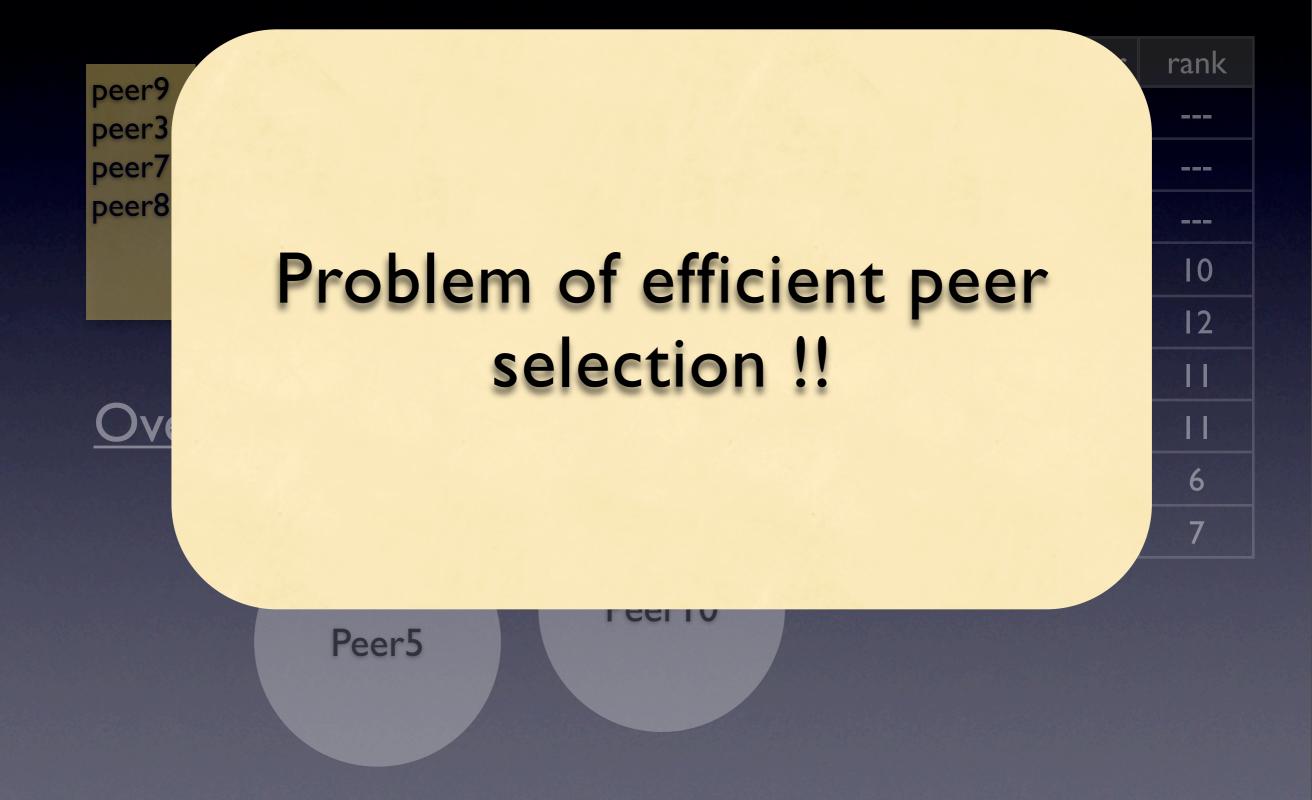


peer	rank
2	
5	dis
10	
9	10
3	12
I	
7	11
15	6
8	7





peer	rank
2	
5	
10	
9	10
3	12
I	Ш
7	11
15	6
8	7



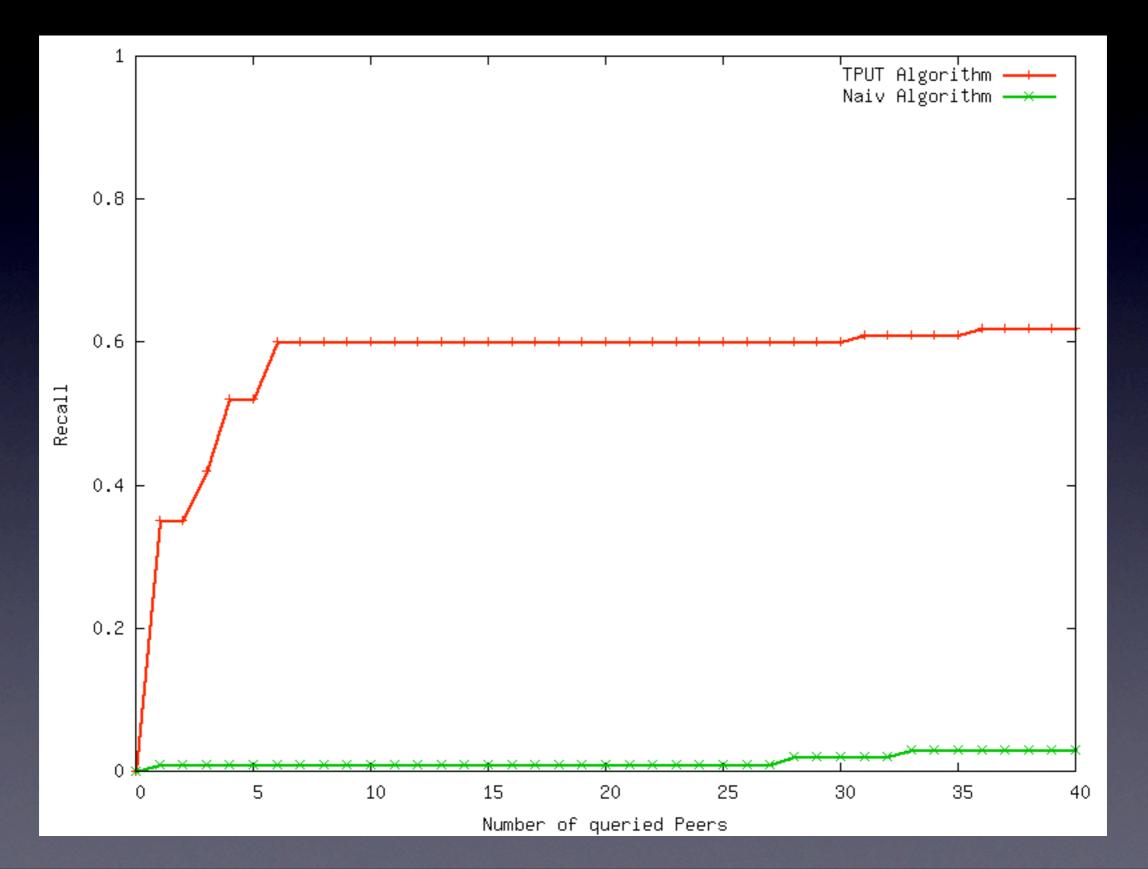
Problems

- TPUT not applicable for Overlap Aware Routing, needs a stable rank
- Naiv Algorithm could be to iterate over all peers, but too expensive
- "Batching": Select parts of documentlists from every peer, but what exactly and from which peer ?

Experimental Setup

- Wikipedia dump (approx. 5.1GB raw data) as test data in PostgreSQL Database
- Linux on Intel 3GHz and IGB memory
- Data clustered in 1.000 different peers (with overlaps)
- Data Tables:
 - Table for statistics, includes: peerid, term
 - Main table with all informations (term, score, termfrequency, docid, peerid, doclength, ...)
 - A lot of indices

Results



Conclusion / Ongoing Work

- Problem in peer selection if rank / score changes in every iteration step
- Framework running, queries working, TPUT fully implemented
- <u>Now:</u> Find algorithm that computes peer selection better and efficient for dynamical ranking condition

End

Thank you