

Universität des
Saarlandes
FR 6.2 Informatik

Dr. Ernst Althaus, Dr. Benjamin Doerr, David Steurer
SS 2005

## Exercises for Optimization

1. Assignment

Due 29.04.2005

Exercise $1(9 \times 1+2$ Points)
Let $A=\left(a_{1}, \ldots, a_{n}\right)$ be a non-singular $n \times n$ real matrix with columns $a_{i}$ and let $\operatorname{adj}(A)$ denote the $n \times n$ matrix with $\operatorname{adj}(A)_{j i}=(-1)^{i+j} \operatorname{det}\left(m_{i j}(A)\right)$ where $m_{i j}(A)$ is the $(n-1) \times(n-1)$ submatrix obtained by deleting row $i$ and column $j$ from $A$. The matrix $\operatorname{adj}(A)$ is called adjugate ("Adjungierte") of $A$.
a) Give the definition of linear independence.
b) State "the" two formulae for computing the determinant of a matrix.
c) Which of the following statements are true for all $A, B \in \mathbb{R}^{n \times n}$ ? (without proof)
(a) $\operatorname{det} A=0 \Leftrightarrow$ columns (rows) of $A$ are linearly dependent
(b) $\operatorname{det}(A+B)=\operatorname{det} A+\operatorname{det} B$
(c) $\operatorname{det}(A \cdot B)=\operatorname{det} A \cdot \operatorname{det} B$
d) Show $|\operatorname{det} A| \leq n!\|A\|_{\infty}^{n}$. ( $\|\cdot\|_{\infty}$ denotes the maximum absolute value in a matrix or vector.)
e) Let $b \in \mathbb{R}^{n}$. Show that $\operatorname{adj}(A) \cdot b=\left(\begin{array}{c}\operatorname{det}\left(m_{1}(A, b)\right) \\ \vdots \\ \operatorname{det}\left(m_{n}(A, b)\right)\end{array}\right)$ where $m_{i}(A, b)$ is obtained from $A$ by replacing column $i$ by vector $b$.
f) Show $\operatorname{adj}(A) \cdot A=I \cdot \operatorname{det} A$ where $I \in \mathbb{R}^{n \times n}$ is the identity matrix.
g) State and prove Cramer's Rule. (Hint. Deduce $\frac{1}{\operatorname{det} A} A \cdot \operatorname{adj}(A)=I$ from f).)
h) Suppose $A$ is an integer matrix. Show $A^{-1}$ is an integer matrix if and only if $|\operatorname{det} A|=1$.
(Remark. $A$ is called unimodular if $A, A^{-1} \in \mathbb{Z}^{n \times n}$.)
i) Suppose $A$ and $b$ have integer entries. Let $x$ be such that $A x=b$. Show that each $x_{i}$ is rational. Let $x_{i}=p_{i} / q_{i}$ for integers $p_{i}, q_{i}$ with $\operatorname{gcd}\left(p_{i}, q_{i}\right)=1$. Give upper bounds for $\left|p_{i}\right|$ and $\left|q_{i}\right|$ in terms of $n,\|A\|_{\infty}$ and $\|b\|_{\infty}$.
(Remark. This shows that the representation size of $x$ is polynomial in the representation size of $A$ and $b$.)
j) Show that for any matrix $A \in \mathbb{R}^{n \times m}$ there is a matrix $Y$ such that

$$
\left\{A x \mid x \in R^{m}\right\}=\left\{b \in \mathbb{R}^{n} \mid Y^{T} b=0\right\} .
$$

Can you find a geometric interpretation? (Hint. This is not related to Cramer's Rule.)

## Exercise $2(4 \times 1$ Points $)$

For the following exercises only use that the dual of a linear program of the form

$$
(P): \text { maximize } c^{T} x \text { subject to } A x \leq b, x \geq 0
$$

is the linear program

$$
(D): \text { minimze } y^{T} b \text { subject to } y^{T} A \geq c^{T}, y \geq 0
$$

a) Let $x$ and $y$ be feasible solutions to $(P)$ and $(D)$, respectively. Prove algebraically that $c^{T} x$ is at most $y^{T} b$. (Remark. This fact is called weak duality.)
b) Transform the problem $(D)$ to an equivalent problem $\left(D^{\prime}\right)$ of the same form than $(P)$.
c) Determine the dual of $\left(D^{\prime}\right)$. Call this problem $\left(P^{\prime}\right)$.
d) What is the relation between $(P)$ and $\left(P^{\prime}\right)$ ?

Exercise $3(1+1+2+1$ Points $)$
Consider the linear problem

$$
(P): \text { maximize } c x_{1}+d x_{2} \text { subject to } x_{1} \leq 1, x_{2} \leq 1, x_{1}+x_{2} \leq 1
$$

a) Graph the feasible region.
b) Determine the dual problem $(D)$ of $(P)$. (Hint. $\min _{y^{T} A=c^{T}, y \geq 0} y^{T} b$ is dual to $\max _{A x \leq b} c^{T} x$.)
c) For each $c, d$ show whether $(P)$
(a) is infeasible,
(b) is unbounded,
(c) has exactly one optimum solution or
(d) has more than one optimum solution.

In case of $c$ ) or $d$ ), compute an optimum solution and prove optimality using the dual.
d) Consider optimum solutions $x$ and $y$ for $(P)$ and $(D)$, respectively. Can you observe a connection between the positive entries of $y$ and the constraints in $(P)$ that are fulfilled with equality by $x$ ?
What happens to $(D)$ if $(P)$ is unbounded?

