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SS 2005

Exercises for Optimization

3. Assignment

Due 13.05.2005

Exercise 1 (3 Points)

Given the linear program

$$(P) : \text{minimize } c^T x \text{ subject to } Ax = b,$$

use the dual program to show that exactly one of the following 3 possibilities holds:

- (P) is infeasible,
- (P) is unbounded,
- all feasible solutions of (P) are optimal.

Exercise 2 (3 Points)

Derive and prove a Theorem of the Alternatives for the system

$$Ax < b.$$

(*Hint.* Consider the problem to maximize t subject to $Ax + t\mathbb{1} \leq b$ and its dual problem. $\mathbb{1}$ is the all-ones vector.)

Exercise 3 (3 Points)

Is $x^* = (0, 4/3, 2/3, 5/3, 0)^T$ an optimal solution of (P) ? (*Hint.* Use complementary slackness.)

$$(P) \quad \begin{array}{ll} \max & 7x_1 + 6x_2 + 5x_3 - 2x_4 + 3x_5 \\ \text{s.t.} & x_1 + 3x_2 + 5x_3 - 2x_4 + 2x_5 \leq 4 \\ & 4x_1 + 2x_2 - 2x_3 + x_4 + x_5 \leq 3 \\ & 2x_1 + 4x_2 + 4x_3 - 2x_4 + 5x_5 \leq 5 \\ & 3x_1 + x_2 + 2x_3 - x_4 - 2x_5 \leq 1 \\ & x_i \geq 0 \end{array}$$

Exercise 4 (3 Points)

Formulate and prove the theorem of complementary slackness for the following linear program and its dual program:

$$\begin{array}{ll} \max & c^T x + d^T y \\ \text{s.t.} & Ax + By \leq a \\ & Cx + Dy = b \\ & x \geq 0 \end{array}$$

You may use the rules for forming the dual and the theorem of strong duality.

Exercise 5 (2 + 2 Points)

We are given the following optimization problem:

A factory is producing two types of (waist-)belts A and B. A is of better quality than B. The profit is 2 for type A and 1,50 for type B. The production of one belt of type A needs 2 units of time whereas only 1 time unit is needed to produce one belt of type B. There is a daily capacity of 1000 units of time. The leather delivery allows a production of 800 belts per day. There are at most 400 belt buckles for type A per day available and 700 for type B.

Which amount of each type A and B should be produced to maximize the profit?

- a) Formulate this problem as a linear program.
- b) Solve this problem graphically.

Exercise 6 (2 + 2 + 2* + 2* + 2* Points)

Let $D = (V, A)$ be a directed graph with capacities $c : A \rightarrow \mathbb{R}_{\geq 0}$, source $s \in V$ and sink $t \in V$.

Suppose each arc $a \in A$ can carry a non-negative amount of flow of at most $c(a)$. The total amount of flow entering a vertex $v \neq s, t$ should be equal to the total amount of flow leaving v .

The *maximum flow problem* asks for maximizing the net amount of flow leaving s , that is the difference between the amount of flow leaving s and the amount of flow entering s .

The following notation may be useful.

$$\delta^-(U) := \text{set of arcs entering } U \subseteq V \text{ in } D$$

$$\delta^+(U) := \text{set of arcs leaving } U \subseteq V \text{ in } D$$

$$\delta^\pm(u) := \delta^\pm(\{u\}) \text{ for } u \in V$$

You earn bonus points for c)–e).

- a) Formulate the maximum flow problem as a linear program (P).
- b) Determine the corresponding dual program (D).
- c) Let x^* and y^* be optimal solutions for (P) and (D), respectively.
Formulate the complementary slackness conditions for x^* and y^* .
- d) Show that there exists W with $s \in W \subseteq V - t$ such that all arcs $\delta^-(W)$ carry no flow and each arc $a \in \delta^+(W)$ carries flow $c(a)$ in x^* . (*Hint.* Use complementary slackness.)
- e) Let OPT denote the objective function value of x^* . Now show

$$OPT = \sum_{a \in \delta^+(W)} c(a) \leq \sum_{a \in \delta^+(U)} c(a) \text{ for all } U \text{ with } s \in U \subseteq V - t.$$

Remark. This fact is called ‘max-flow min-cut’ theorem.