

- This problemset has *three* questions.
- For other questions, either send the solutions to `gawry1+aos@gmail.com`, or leave them in the envelope attached to the doors of my office (room 321).

1. You are given a LZW parse (so, a sequence of n blocks, each block being either a single letter, or a previously defined block concatenated with a single character). Construct a small structure allowing you to access any letter of the corresponding text efficiently. For full credit, construct a structure of size $\mathcal{O}(n)$ allowing answering any such query in $\mathcal{O}(\log n)$ time, for partial credit the size of your structure can be larger, but should be $o(n^2)$.
2. Prove that for **any** text of length n over a fixed finite alphabet Σ , its LZW parse consists of at most $\mathcal{O}(\frac{n}{\log n})$ blocks (the constant hidden under the big-O depends on $|\Sigma|$, though).
3. A primitive square is a word of the form xx with x being primitive, which means that it is not possible to write $x = y^k$ with $k > 1$. For instance `abaaba` is a primitive square, but `abababab` is not. We want to get a good bound on the maximal number of subwords of a word of length n that are primitive squares. Note that if the same primitive square occurs multiple times in the word, we count it multiple times.
 - (a) Prove a simplified version of the “three squares” lemma: if there are three primitive words x, y, z such that yy is a proper prefix of xx , and zz is a proper prefix of yy , then $|x| \geq 2|z|$.
Hint: assume that $|x| < 2|z|$, draw a picture, then try to apply the periodicity lemma to deduce that z is actually not primitive.
 - (b) Prove the following “three squares” lemma: if there are three primitive words x, y, z such that yy is a proper prefix of xx , and zz is a proper prefix of yy , then $|x| \geq |y| + |z|$.
Hint: this is tricky and for extra credit.
 - (c) Use the above lemma (in either version) to bound the number of primitive squares that all begin at the same position. Observe that multiplying this bound by n gives you a bound on the total number of primitive squares.
- (SPOJ) 4. Given a word $w[1..n]$ find the largest k such that w contains a substring of the form u^k , for some nonempty word u .
Hint: $\mathcal{O}(n \log n)$ is enough. Guess $|u|$ and try to compute the largest k in $\mathcal{O}(\frac{n}{|u|})$ time. using some longest common prefix/suffix queries.