

max planck institut informatik



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# Assignment 2 for Approximation Algorithms and Hardness of Approximation Discussion: Thursday, 8 May 2014, 14 pm

## Assignment 1 (PTAS for Knapsack)

In the last tutorial, we discussed the following two observations about GREEDYKNAPSACK:

1. If  $p_i \leq \varepsilon \text{OPT}$  for all *i*, GREEDYKNAPSACK gives a  $(1 - \varepsilon)$ -approximation.

2. There are at most  $\lceil \frac{1}{\varepsilon} \rceil$  items with profit at least  $\varepsilon$ OPT in any optimal solution.

Assignment 2 (Tightness of TSP approximation algorithms)

Find examples for which

- 1. METRICTSPVIAMST does not find a  $\alpha$ -approximation for any constant  $\alpha < 2$ .
- 2. METRICTSPVIAMSTANDMATCHING does not find a  $\alpha$ -approximation for any constant  $\alpha < \frac{3}{2}$ .

### Assignment 3 (Asymmetric TSP)

We are given a directed graph G with edge costs  $d: V \times V \to \mathbb{R}_{\geq 0}$  that satisfy the directed triangle inequality, i.e.,  $d(u, v) \leq d(u, w) + d(w, v)$  for all  $u, v, w \in V$ . Give a  $O(\log n)$ -approximation for the problem of finding a shortest tour visiting all vertices.

Use the minimum-cost cycle cover. Shrink the cycles and recurse.

### Assignment 4 (Approximating maximum compression)

Let X be a set of strings and let ||X|| denote the sum of lengths of the strings in X. Consider the *maximum compression* problem: find a superstring s of all strings in X that maximizes ||X|| - |s|. Call such an optimal string  $s^*$ . Show that in polynomial time, one can compute a  $\frac{1}{2}$ -approximation, i.e., a superstring s of all strings in X with compression

$$||X|| - |s| \ge \frac{1}{2} (||X|| - |s^*|).$$

Then use an approximation algorithm for maxTSP.

Hint: Reduce this problem to finding a longest traveling salesman path in a suitably constructed graph.

#### Open problem

Consider the greedy algorithm for the shortest superstring problem: Take two strings with maximum overlap and replace them by overlapping them as much as possible. Repeat until a single string is left. Does this algorithm achieve a 2-approximation?