

Universität des Saarlandes FR 6.2 Informatik



Antonios Antoniadis and Martin Hoefer

SS 2014

Excercises Online Algorithms

http://www.mpi-inf.mpg.de/departments/d1/teaching/ss14/OnlineAlgos/

Sheet 1

Deadline: 1.05.2014

Rules: Until the end of the semester you have to reach 50% of the achievable points to be admitted to the exam.

Exercise 1 (*10 points*) It can be shown, that no deterministic online algorithm for the cowpath problem can obtain a competitive ratio strictly less than 9. Your task is to prove a (possibly weaker) lower bound ℓ on the competitive ratio that can be obtained by any deterministic online algorithm for the cow-path problem.

Hint: An algorithm can be described by a sequence of turns. Let f(i) be the number of steps that the cow makes between last crossing the origin and turn *i*.

Grading: Assuming that your proof is correct you will be assigned:

3 points,	$\text{if } 3 \leq \ell < 5,$
10 points,	$\text{if } 5 \leq \ell < 8.99,$
15 points(10 + 5 bonus points),	if $\ell \ge 8.99$.

Exercise 2 (10 points)

Describe and analyze, a randomized online algorithm for the cow-path problem, that has a competitive ratio of at most 7.

Exercise 3 (4+6 *points*)

Prove that algorithms LIFO (Last In First Out) and LFU (Least Frequently Used), both have an unbounded competitive ratio for the paging problem.

Exercise 4 (5+5 points)

The *Online Bin Packing* problem: we have an infinite supply of unit-size bins, and items of sizes $a_1, a_2, \ldots a_n$ arrive over time (we may assume that for all $i, 0 < a_i \le 1$). The goal is to pack the items into bins, so that the capacity is not exceeded on any bin, and the number of used bins is minimized.

The *First Fit* algorithm for the online bin packing problem: Consider the bins in some fixed prespecified order, and once an item arrives pack it to the first bin it fits.

Prove that:

- (i) First Fit has a competitive ratio of 2, and
- (ii) First Fit has a competitive ratio of at least 3/2, i.e., give a sequence of items for which First Fit uses at least 3/2 times the number of bins of an optimal solution.