By random variables or discrete random variables we mean random variables taking either finitely many values or countably infinite values.

Short problems. Each of the following short problems comprise of five credits.

- 1. Let X be a random variable.
 - (a) For all even integer $k \ge 2$, show that $\mathbf{E}[X^k] \ge \mathbf{E}[X]^k$.
 - (b) If X is a non-negative integer valued random variable then $\mathbf{Pr}[X > 0] \leq \mathbf{E}[X]$.
 - (c) Show that $\mathbf{Pr}[X=0] \leq \frac{\mathbb{V}ar[X]}{\mathbf{E}[X]^2}$.
- 2. Let a coin be such that heads comes up independently with probability p on each flip. What is the expected number of flips until you get k heads?
- 3. Suppose cards are being drawn at random with replacement from a deck of n cards.
 - (a) What is the expected number of cards we must draw until we have seen all n cards?
 - (b) Suppose we have drawn only 2n times. What is the expected number of different cards chosen from the deck?
- 4. Consider the random walk on the integer numbers starting at 0. Let F(n) be the expected number of steps for this walk to reach either -n or +n. Give exact expression for F(n).
- 5. For $n \ge 4$ and let H = (V, E) be an *n*-uniform hypergraph.
 - (a) If $|E| \leq 4^{n-1}$, then there exists a coloring of vertices of H such that very edge in H contains vertices with at least two different colors.
 - (b) If $|E| \leq \frac{4^{n-1}}{3^n}$, then show that there exists a coloring of vertices of H such that very edge in H contains vertices of all four colors.
- 6. We saw in class that $(X,d) \xrightarrow{O(\log n)} l_1^{O(\log^2 n)}$ where *n* is the number of points in the metric space (X,d), i.e., #X = n. Using this result and Holder's inequality show that $(X,d) \xrightarrow{O(\log n)} l_p^{O(\log^2 n)}$.

Hölder's inequality: $||x||_p \cdot ||y||_q \ge \langle x, y \rangle$ where $\frac{1}{p} + \frac{1}{q} = 1$.

7. Let $\mathcal{F} = \{S_1, \ldots, S_m\}$ be a collection of m sets with $S_i \subset [n] = \{1, \ldots, n\}$. Given an assignment $\chi : [n] \to \{1, -1\}$, the *discrepancy* of a set $S \subseteq [n]$ is defined as

$$\operatorname{disc}_{\chi}(S) = \mid \sum_{x \in S} \chi(x) \mid .$$

Discrepancy of \mathcal{F} wrt χ is defined as

$$\operatorname{disc}_{\chi}(\mathcal{F}) = \max_{S_i \in \mathcal{F}} \operatorname{disc}_{\chi}(S_i).$$

Using a random assignment $\chi : [n] \to \{1, -1\}$ and Chernoff bound show existence of a coloring such that the discrepancy of \mathcal{F} wrt to that coloring is $O(\sqrt{n \log m})$.

8. Let C_n denotes a cycle (graph) with *n*-vertices, and d(,) denotes the shortest distance metric between the vertices of C_n (note that we assume edges in C_n and unit length). Show that C_n can be embedded in a famility of trees \mathcal{F}_n on *n*-vertices such that

$$1 \le \frac{E_{T \leftarrow \mathcal{F}}[d_T(x, y)]}{d(x, y)} < 2,$$

where $d_T(x, y)$ denotes the shorest distance metric in the tree T.

- 9. Let p be a prime number, and m, n are positive integers. Given a set of m linear equalities over n variables modp show that there exists an assignment to the variables that will satisfy at least $\frac{1}{p}$ of linear equalities. What happens if p is not a prime?
- 10. Consider the general setting of the Lovász local lemma, i.e. probability space Ω , and "bad" events A_1, \ldots, A_m which we want to forbid. Now let B be another event in the same probability space. Suppose that B depends on the set of events $\Gamma(B) \subseteq \{A_1, \ldots, A_m\}$. Then prove that,

$$\mathbf{Pr}\left[B \mid \prod_{i=1}^{m} \bar{A}_i\right] \leq \frac{\mathbf{Pr}[B]}{\prod_{i \in \Gamma(B)} (1-x_i)}$$

where x_i 's are the reals corresponding to the events A_i in the usual statement of the Lovász Local Lemma.

Long problems. Each of the following long problems comprise of 10 credits.

- 1. Consider a fair die, i.e., when we roll the die the probability we get i, for $i \in \{1, 2, 3, 4, 5, 6\}$, is equal to 1/6. What is the expected number of rolls until the first pair of consecutive sixes appears?
- 2. Given a graph G = (V; E) with V = n, a dominating set for G is a subset $D \subseteq V$ such that each vertex $v \in V$ is either in D or has a neighbor in D.
 - (a) Show that any graph with minimum degree δ has a dominating set of size at most $\frac{n \log n}{\delta + 1}$.
 - (b) Improve the bound for dominating set to $\frac{n(1+\log(\delta+1))}{\delta+1}$
- 3. See the definition of discrepancy given above in problem-(7). In that problem we were interested in upper bounding discrepancy of a family of sets \mathcal{F} . In this problem we want to lower bound discrepancy, i.e., we want to show there exists a family of sets \mathcal{F} with n subsets of [n] such that for every coloring $\chi: [n] \to \{1, -1\}, \operatorname{disc}_{\chi}(\mathcal{F}) > \Omega(\sqrt{n})$. Complete the proof by proving the following subproblems:
 - (a) For a fixed $\chi : [n] \to \{1, -1\}$, pick a subset $S \subseteq [n]$ by including each element in S with probability $\frac{1}{2}$. Show that there exists a constant c > 0 such that

$$\mathbf{Pr}\left[\operatorname{disc}_{chi}(S) > \sqrt{n}/c\right] > 1/2.$$

- (b) Let \mathcal{F} consists of n sets picked independently as above. Show that for any fixed assignment $\chi: [n] \to \{1, -1\}, \operatorname{disc}_{\chi}(\mathcal{F}) > \sqrt{n/c}$ with probability $> 1 \frac{1}{2^n}$.
- (c) Using union bound over all 2^n assignments to show existence of a family \mathcal{F} with n sets for which all assignments $\chi : [n] \to \{1, -1\}$ have discrepancy $> \sqrt{n/c}$.
- 4. Let F be an infinite family of graphs. Let $G \in \mathcal{G}(n, p)$, where p = p(n) and $\mathcal{G}(n, p)$ denotes the Erdős-Reńyi random graph model. Suppose F, p are such that the expected number of subgraphs on k vertices, belonging to the family F, in G is t > 1, t fixed. Prove that there exists a graph on n vertices which does not contain any subgraph of size k from the family F.

5. Let G = (V, E) be an undirected graph and suppose each $v \in V$ is associated with a set S(v) of 8r colors, where $r \geq 1$. Suppose, in addition that for each $v \in V$ and $c \in S(v)$ there are at most r neighbours u of v such that c lies in S(u). Prove that there is a proper coloring of G assigning to each vertex v a color from its class S(v) such that, for any edge $(u, v) \in E$, the colors assigned to u and v are different.

[Hint: Consider the family of events $A_{u,v,c}$, such that u and v are both colored with color c.]

6. Consider the Lollipop graph given below:



The top part is a complete graph K_n on n vertices, and the lower tail is a path of length n. Compute tight bounds for the following quantities:

- (a) The expected time of the random walk starting at u to arrive at v.
- (b) The expected time of the random walk starting at v to arrive at u.
- (c) The expected time of the random walk starting at v to visit all the vertices in the graph.
- (d) The expected time of the random walk starting at u to visit all the vertices in the graph.
- 7. Show that for a *n* point set $P \subset \mathbb{R}^D$ and $d = O\left(\frac{\log n}{\varepsilon^2}\right)$ there exists a linear map $f : \mathbb{R}^D \to \mathbb{R}^d$ such that for all $p_i, p_j, p_k \in P$,

$$|(p_j - p_i)^{\mathrm{T}}(p_k - p_i) - (f(p_j) - f(p_i))^{\mathrm{T}}(f(p_k) - f(p_i))| \le \varepsilon ||p_j - p_i|| ||p_k - p_i||.$$

[Hint: See the proof of Johnson-Lindenstrauss lemma covered in class.]

- 8. Suppose we are given a discrete probability distribution X on [n], i.e., we are given $p_i \ge 0$ such that $\mathbf{Pr}[X=i] = p_i$ and $\sum_{i=1}^{n} p_i = 1$. We want to preprocess this distribution in O(n) time such that given an oracle that uniformly generates random number from the interval [0,1], we should be able to sample the distribution X in O(1) time for each query. Assume that we are working in the real RAM model of computation.
- 9. We are given n points distributed uniformly at random within the unit square in a plane. Each point connects to the k-closest points. Let us denote the resulting graph as $\Gamma(n,k)$. Show that there exists α such that if $k \ge \alpha \log n$, then $\Gamma(n,k)$ is connected with probability at least 1 1/n.

10. A random variable X taking on positive integer values is said to be a Poisson random variable with parameter $\lambda > 0$ if

$$\mathbf{Pr}[X=i] = \exp\left(-\lambda\right)\frac{\lambda^i}{i!}$$

for $i = 0, 1, 2, \ldots$

- (a) Show that $\mathbf{E}[X] = \lambda$.
- (b) Let X_1, \ldots, X_n be identically distributed Poisson random variable with parameter λ , and let $X = \sum_{i=1}^n X_i$. Show that
 - i. Show that X is a Poisson random variable with parameter $n\lambda$.
 - ii. $\mathbf{Pr}[X > (1+\epsilon)n\lambda] \le \exp(-\epsilon^2 n\lambda).$