

Algorithmic Game Theory

Summer 2015

Exercise Set 1

Please hand in your solutions at the beginning of the lecture on May 4.
 You may work in groups of up to three students.

Exercise 1: (5 + 3 Points)

Given the game with the following costs (lower left corner of a cell refers to the row player; upper right corner to the column player):

	A	B	C	D
A	4 6	6 5	5 4	6 5
B	5 6	1 5	2 1	5 5
C	6 6	5 3	3 3	1 2
D	7 6	7 3	6 5	7 7

- a) What are the pure Nash equilibria? Explain why exactly those states are pure Nash equilibria and why no other state is one.
- b) Why can no other strategy in this game be part of any mixed Nash equilibrium?
 Or: What are the strategies chosen in a mixed Nash equilibrium and why?

Exercise 2: (4 Points)

The Hit-Below Game is given as follows. Each of n players announces a number in $\{1, \dots, k\}$. A prize of 1 Euro is split equally among the players whose chosen number is closest to $2/3$ times the average number. Show that the Hit-Below Game has a unique Nash equilibrium and that in this equilibrium each player picks a pure strategy.

Exercise 3: (4 Points)

Consider the following four sets:

$$\begin{aligned}
 A &= \{x \in \mathbb{R}^2 \mid x_1, x_2 \geq 0 \text{ and } x_1^2 + x_2^2 = 1\} \\
 B &= \{x \in \mathbb{R}^2 \mid x_1, x_2 \geq 0 \text{ and } 1 \leq x_1^2 + x_2^2 \leq 2\} \\
 C &= \{x \in \mathbb{R}^2 \mid x_1^2 + x_2^2 \leq 2\} \\
 D &= \{x \in \mathbb{R}^2 \mid x_1^2 + x_2^2 < 1\}
 \end{aligned}$$

To which of those sets does Brouwer's fixed point theorem apply? For the other ones give a continuous function that has no fixed point.

Hint: In one dimension you need a closed set D because $f: (0, 1] \rightarrow (0, 1]$ with $f(x) = \frac{x}{2}$ does not have a fixed point. You need a bounded set because $f: [0, \infty) \rightarrow [0, \infty)$ with $f(x) = x+1$ does not have a fixed point.

Exercise 4:

(4 Points)

Note: In the exercise, we consider *utilities* (or payoffs) rather than *costs*. Equilibria etc. are defined as if utilities were negative costs.

Consider the *single-item first-price auction game* with two players. Each player $i \in \{1, 2\}$ has strategies $[0, 1]$; his utility is parameterized by a value $v_i \in [0, 1]$. In state (b_1, b_2) , the utilities are given as

$$u_1(b_1, b_2) = \begin{cases} v_1 - b_1 & \text{if } b_1 \geq b_2 \\ 0 & \text{otherwise} \end{cases} \quad \text{and} \quad u_2(b_1, b_2) = \begin{cases} 0 & \text{if } b_1 \geq b_2 \\ v_2 - b_2 & \text{otherwise} \end{cases} .$$

Under what condition (regarding v_1 and v_2) is there a pure Nash equilibrium? State all pure Nash equilibria.