# Chapter IV: Link Analysis

Information Retrieval & Data Mining Universität des Saarlandes, Saarbrücken Wintersemester 2013/14

#### Friendship Networks, Citation Networks, ...

• Link analysis studies the relationships (e.g., friendship, citation) between objects (e.g., people, publications) to find out about their characteristics (e.g., popularity, impact)

- Social Network Analysis (e.g., on a friendship network)
  - Closeness centrality of a person v is the fraction of shortest paths between any two persons (u, w) that pass through v

- Bibliometrics (e.g., on a citation network)
  - Co-citation measures how many papers cite both u and v
  - Co-reference measures how many common papers both u and v refer to

#### ..., and the Web?

- World Wide Web can be seen as directed graph G(V, E)
  - web pages correspond to vertices (or, nodes) V
  - hyperlinks between them correspond to edges E
- Link analysis on the Web graph can give us clues about
  - which web pages are **important** and should thus be ranked higher
  - which pairs of web pages are similar to each other
  - which web pages are probably **spam** and should be ignored

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### **Chapter IV: Link Analysis**

## IV.1 The World Wide Web as a Graph Degree Distributions, Diameter, Bow-Tie Structure

## IV.2 PageRankRandom Surfer Model, Markov Chains

## IV.3 HITS Hyperlinked-Induced Topic Search

## IV.4 Topic-Specific and Personalized PageRank Biased Random Jumps, Linearity of PageRank

## IV.5 Online Link Analysis OPIC

## IV.6 Similarity Search SimRank, Random Walk with Restarts

## IV.7 Spam DetectionLink Spam, TrustRank, SpamRank

## IV.8 Social Networks SocialPageRank, TunkRank

### IV.1 The World Wide Web as a Graph

- 1. How Big is the Web?
- 2. Degree Distributions
- 3. Random-Graph Models
- 4. Bow-Tie Structure

Based on MRS Chapter 21

#### 1. How Big is the Web?

- How big is the entire World Wide Web?
  - quasi-infinite when you consider all (dynamic) URLs (e.g., of calendars)
- Indexed Web is a more reasonable notion to look at
  - [Gulli and Signori '05] estimated it as 11.5 billions (109) in 2005
  - Google claimed to know about more than 1 trillion (10<sup>12</sup>) URLs in 2008
  - <u>WorldWideWebSize.com</u> provides daily estimates obtained by extrapolating from the number of results returned by Google and Bing on the basis of Zipf's law (currently: 3.6 billion 38 billion)

#### 2. Degree Distributions

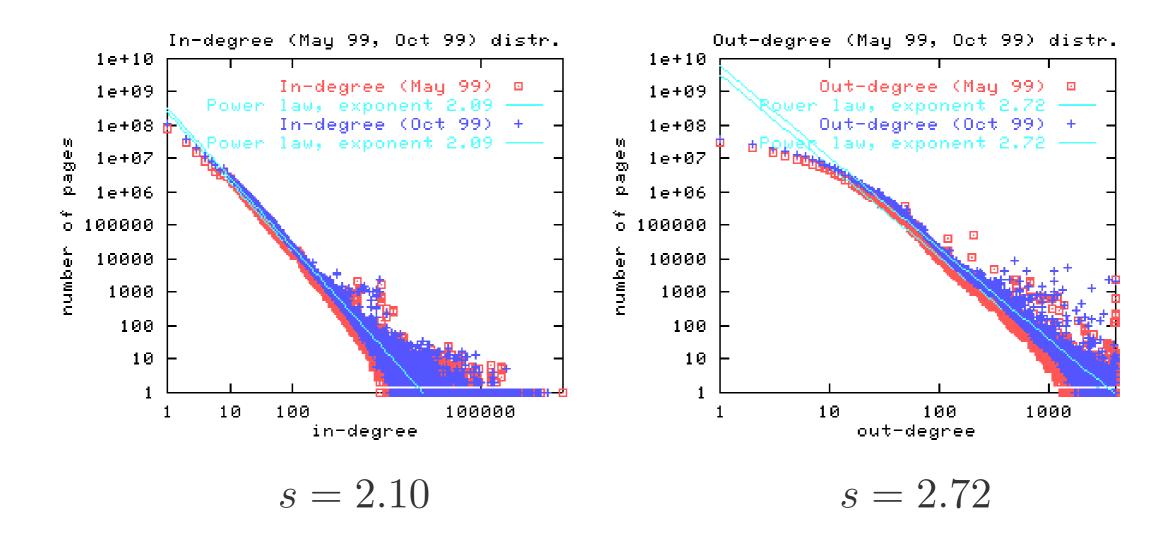
- What is the distribution of in-/out-degrees on the Web graph?
  - in-degree(v) of vertex v is the number of incoming edges (u, v)
  - out-degree(v) of vertex v is the number of outgoing edges (v, w)
- Zipfian distribution has probability mass function

$$f(k; s, N) = \frac{1/k^s}{\sum_{n=1}^{N} 1/n^s}$$

with rank k, parameter s, and total number of objects N

- provides good model of **many real-world phenomena**, e.g., word frequencies, city populations, corporation sizes, income rankings
- appear as **straight line** with slope -s in **log-log-plot**

### Degree Distributions



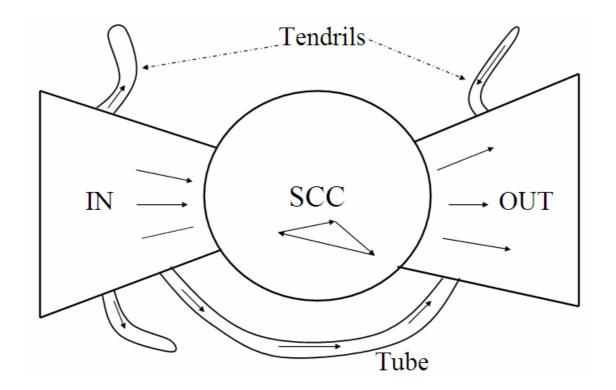
• Full details: [Broder et al. '00]

#### 3. Random-Graph Models

- Generative models of undirected or undirected graphs
- Erdös-Renyi Model G(n, p) generates a graph consisting of n vertices; each possible edge (u, w) exists with probability p
- Barabási-Albert Model generates a graph by successively adding vertices u with m edges; the edge (u, v) attaches to vertex v with probability proportional to deg(v)
- Preferential attachment ("the rich get richer") in the Barabási-Albert Model yields graphs with properties similar to Web graph
- Full details: [Barabási and Albert '99]

#### 4. Bow-Tie Structure

• The Web graph looks a lot like a **bow tie** [Broder et al. '00]



- Strongly Connected Component (SCC) of web pages that are reachable from each other by following a few hyperlinks
- IN consisting of web pages from which SCC is reachable
- OUT consisting of web pages reachable from SCC

#### Additional Literature for IV.1

- A.-L. Barabási and R. Albert: Emergence of Scaling in Random Networks, Science 1999
- A. Broder, R. Kumar, F. Maghoul, P. Raghavan, S. Rajagopalan, R. Stata, A. Tomkins, and J. L. Wiener: *Graph Structure in the Web*, Computer Networks 33:309-320, 2000
- A. Gulli and A. Signori: The Indexable Web is More than 11.5 Billion Pages, WWW 2005
- R. Meusel, O. Lehmberg, C. Bizer: *Topology of the WDC Hyperlink Graph* http://webdatacommons.org/hyperlinkgraph/topology.html, 2013

### IV.2 PageRank

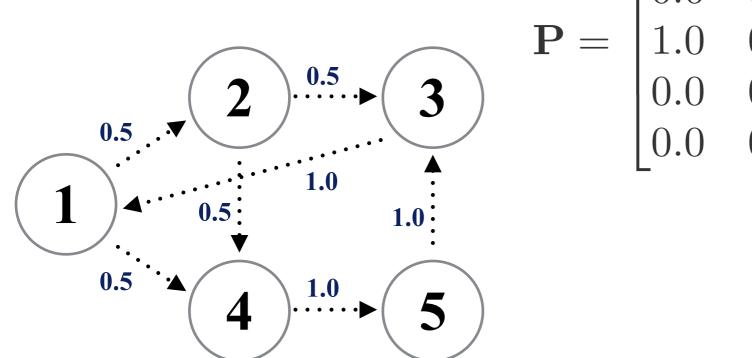
- Hyperlinks distinguish the Web from other document collections and can be interpreted as endorsements for the target web page
- In-degree as a measure of the importance/authority/popularity of a web page v is easy to manipulate and does not consider the importance of the source web pages
- PageRank considers a web page v important if many important web pages link to it
- Random surfer model

Larry Page & Sergey Brin

- follows a uniform random outgoing link with probability (1- $\varepsilon$ )
- jumps to a uniform random web page with probability  $\varepsilon$
- Intuition: Important web pages are the ones that are visited often

#### Markov Chains

 $S = \{1, \dots, 5\}$ 



$$\mathbf{P} = \begin{bmatrix} 0.0 & 0.5 & 0.0 & 0.5 & 0.0 \\ 0.0 & 0.0 & 0.5 & 0.5 & 0.0 \\ 1.0 & 0.0 & 0.0 & 0.0 & 0.0 \\ 0.0 & 0.0 & 0.0 & 0.0 & 1.0 \\ 0.0 & 0.0 & 1.0 & 0.0 & 0.0 \end{bmatrix}$$
1.0

#### Stochastic Processes & Markov Chains

• Discrete stochastic process is a family of random variables

$$\{X_t \mid t \in T\}$$

with  $T = \{0, 1, 2 ...\}$  as discrete time domain

• Stochastic process is a Markov chain if

$$P[X_t = x \mid X_{t-1} = w, \dots, X_0 = a]$$

$$= P[X_t = x \mid X_{t-1} = w]$$

holds, i.e., it is memoryless

• Markov chain is **time-homogeneous** if for all times t

$$P[X_{t+1} = x \mid X_t = w] = P[X_t = x \mid X_{t-1} = w]$$

holds, i.e., transition probabilities do not depend on time

### State Space & Transition Probability Matrix

- State space of a Markov chain  $\{X_t | t \in T\}$  is the countable set S of all values that  $X_t$  can assume
  - $X_t: \Omega \to S$
  - Markov chain is in state s at time t if  $X_t = s$
  - Markov chain  $\{X_t \mid t \in T\}$  is **finite** if it has a finite state space
- If a Markov chain  $\{X_t | t \in T\}$  is finite and time-homogeneous, its transition probabilities can be described as a matrix  $\mathbf{P} = (p_{ij})$

$$p_{ij} = P[X_t = j \mid X_{t-1} = i]$$

• For |S| = n the transition probability matrix **P** is a n-by-n right-stochastic matrix (i.e., its rows sum up to 1)

$$\forall i : \sum_{j} p_{ij} = 1$$

#### Properties of Markov Chains

- State *i* is **reachable** from state *j* if there exists a  $n \ge 0$  such that  $(\mathbf{P}^n)_{ij} > 0$  (with  $\mathbf{P}^n = \mathbf{P} \times ... \times \mathbf{P}$  as *n*-th exponent of  $\mathbf{P}$ )
- States *i* and *j* communicate if *i* is reachable from *j* and vice versa
- Markov chain is **irreducible** if all states  $i, j \in S$  communicate
- Markov chain is **positive recurrent** if the **recurrence probability** is 1 and the **mean recurrence time** is finite for every state *i*

$$\sum_{k=1}^{\infty} P[X_k = i \land \forall 1 \le j < k : X_j \ne i \mid X_0 = i] = 1$$

$$\sum_{k=1}^{\infty} k P[X_k = i \land \forall 1 \le j < k : X_j \ne i \mid X_0 = i] < \infty$$

#### Properties of Markov Chains

• Markov chain is **aperiodic** if every state *i* has period 1 defined as

$$gcd \{ k : P[X_k = i \land \forall 1 \le j < k : X_j \ne i \mid X_0 = i] > 0 \}$$

- Markov chain is **ergodic** if it is time-homogeneous, irreducible, positive recurrent, and aperiodic
- The 1-by-*n* vector  $\boldsymbol{\pi}$  is the **stationary state distribution** of the Markov chain described by  $\mathbf{P}$  if  $\pi_i \geq 0$ ,  $\Sigma \pi_i = 1$ , and

$$\pi \, \mathrm{P} = \pi$$

- $\pi_i$  is the limit probability that Markov chain is in state i
- $1/\pi_i$  reflects the average time until the Markov chain returns to state *i*
- Theorem: If a Markov chain is finite and ergodic, then there exists a unique stationary state distribution  $\pi$

$$\mathbf{P} = \begin{bmatrix} 0.0 & 0.5 & 0.0 & 0.5 & 0.0 \\ 0.0 & 0.0 & 0.5 & 0.5 & 0.0 \\ 1.0 & 0.0 & 0.0 & 0.0 & 0.0 \\ 0.0 & 0.0 & 0.0 & 0.0 & 0.0 \\ 0.0 & 0.0 & 1.0 & 0.0 & 0.0 \end{bmatrix}$$

$$\boldsymbol{\pi}^{0} = \begin{bmatrix} 1.0 & 0.0 & 0.0 & 0.0 & 0.0 \end{bmatrix}$$

$$S = \{1, \dots, 5\}$$

$$\mathbf{P} = \begin{bmatrix} 0.0 & 0.5 & 0.0 & 0.5 & 0.0 \\ 0.0 & 0.0 & 0.5 & 0.5 & 0.0 \\ 1.0 & 0.0 & 0.0 & 0.0 & 0.0 \\ 0.0 & 0.0 & 0.0 & 0.0 & 0.0 \\ 0.0 & 0.0 & 1.0 & 0.0 & 0.0 \end{bmatrix}$$

$$\boldsymbol{\pi}^{0} = \begin{bmatrix} 1.0 & 0.0 & 0.0 & 0.0 & 0.0 \\ 1.0 & 0.0 & 0.0 & 0.0 & 0.0 \end{bmatrix}$$

$$\boldsymbol{\pi}^{1} = \begin{bmatrix} 0.0 & 0.5 & 0.0 & 0.5 & 0.0 \end{bmatrix}$$

$$S = \{1, \dots, 5\}$$

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$$\mathbf{P} = \begin{bmatrix} 0.0 & 0.5 & 0.0 & 0.5 & 0.0 \\ 0.0 & 0.0 & 0.5 & 0.5 & 0.0 \\ 1.0 & 0.0 & 0.0 & 0.0 & 0.0 \\ 0.0 & 0.0 & 0.0 & 0.0 & 0.0 \\ 0.0 & 0.0 & 1.0 & 0.0 & 0.0 \end{bmatrix}$$

$$\boldsymbol{\pi}^0 = \begin{bmatrix} 1.0 & 0.0 & 0.0 & 0.0 & 0.0 \\ 0.0 & 0.5 & 0.0 & 0.5 & 0.0 \end{bmatrix}$$

$$\boldsymbol{\pi}^1 = \begin{bmatrix} 0.0 & 0.5 & 0.0 & 0.5 & 0.0 \\ 0.0 & 0.0 & 0.25 & 0.25 & 0.5 \end{bmatrix}$$

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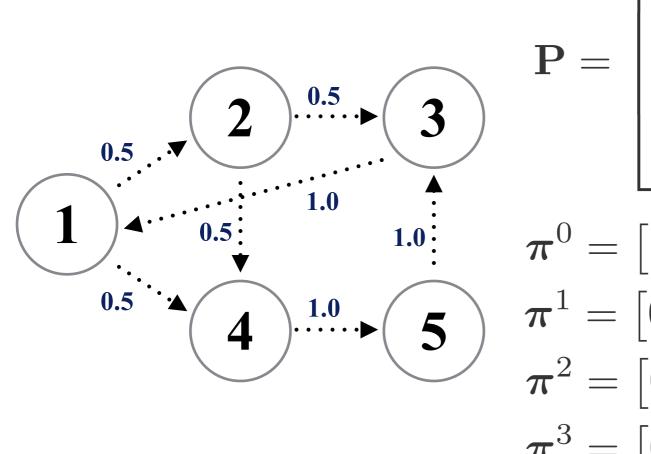
$$\mathbf{P} = \begin{bmatrix} 0.0 & 0.5 & 0.0 & 0.5 & 0.0 \\ 0.0 & 0.0 & 0.5 & 0.5 & 0.0 \\ 1.0 & 0.0 & 0.0 & 0.0 & 0.0 \\ 0.0 & 0.0 & 0.0 & 0.0 & 0.0 \\ 0.0 & 0.0 & 1.0 & 0.0 & 0.0 \end{bmatrix}$$

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$$\boldsymbol{\pi}^{1} = \begin{bmatrix} 0.0 & 0.5 & 0.0 & 0.5 & 0.0 \\ 0.0 & 0.0 & 0.25 & 0.25 & 0.5 \end{bmatrix}$$

$$\boldsymbol{\pi}^{2} = \begin{bmatrix} 0.0 & 0.0 & 0.25 & 0.25 & 0.5 \end{bmatrix}$$

$$\boldsymbol{\pi}^{3} = \begin{bmatrix} 0.25 & 0.0 & 0.5 & 0.0 & 0.25 \end{bmatrix}$$



 $S = \{1, \dots, 5\}$ 

$$\mathbf{7} = \begin{bmatrix} 0.0 & 0.5 & 0.0 & 0.5 & 0.0 \\ 0.0 & 0.0 & 0.5 & 0.5 & 0.0 \\ 1.0 & 0.0 & 0.0 & 0.0 & 0.0 \\ 0.0 & 0.0 & 0.0 & 0.0 & 1.0 \\ 0.0 & 0.0 & 1.0 & 0.0 & 0.0 \end{bmatrix}$$

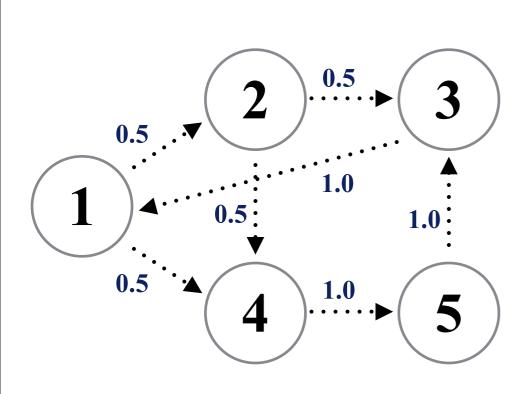
$$\mathbf{7}^{0} = \begin{bmatrix} 1.0 & 0.0 & 0.0 & 0.0 & 0.0 \\ \mathbf{7}^{0} = \begin{bmatrix} 0.0 & 0.5 & 0.0 & 0.5 & 0.0 \end{bmatrix}$$

$$\mathbf{7}^{1} = \begin{bmatrix} 0.0 & 0.5 & 0.0 & 0.5 & 0.0 \end{bmatrix}$$

$$\mathbf{7}^{2} = \begin{bmatrix} 0.0 & 0.0 & 0.25 & 0.25 & 0.5 \end{bmatrix}$$

$$\mathbf{7}^{3} = \begin{bmatrix} 0.25 & 0.0 & 0.5 & 0.0 & 0.25 \end{bmatrix}$$

$$\mathbf{7}^{4} = \begin{bmatrix} 0.5 & 0.125 & 0.25 & 0.125 & 0 \end{bmatrix}$$



 $S = \{1, \dots, 5\}$ 

$$\mathbf{P} = \begin{bmatrix} 0.0 & 0.5 & 0.0 & 0.5 & 0.0 \\ 0.0 & 0.0 & 0.5 & 0.5 & 0.0 \\ 1.0 & 0.0 & 0.0 & 0.0 & 0.0 \\ 0.0 & 0.0 & 0.0 & 0.0 & 1.0 \\ 0.0 & 0.0 & 1.0 & 0.0 & 0.0 \end{bmatrix}$$

$$\boldsymbol{\pi}^{0} = \begin{bmatrix} 1.0 & 0.0 & 0.0 & 0.0 & 0.0 \end{bmatrix}$$

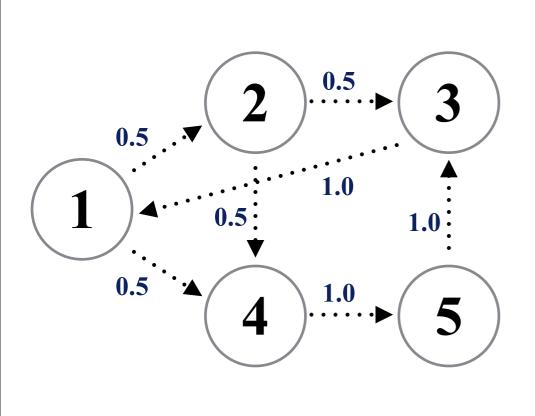
$$\boldsymbol{\pi}^{1} = \begin{bmatrix} 0.0 & 0.5 & 0.0 & 0.5 & 0.0 \end{bmatrix}$$

$$\boldsymbol{\pi}^{2} = \begin{bmatrix} 0.0 & 0.5 & 0.0 & 0.5 & 0.0 \end{bmatrix}$$

$$\boldsymbol{\pi}^{3} = \begin{bmatrix} 0.25 & 0.0 & 0.5 & 0.0 & 0.25 \end{bmatrix}$$

$$\boldsymbol{\pi}^{4} = \begin{bmatrix} 0.5 & 0.125 & 0.25 & 0.125 & 0 \end{bmatrix}$$

 $\boldsymbol{\pi}^5 = \begin{bmatrix} 0.25 & 0.25 & 0.0625 & 0.3125 & 0.125 \end{bmatrix}$ 



 $S = \{1, \dots, 5\}$ 

$$\mathbf{P} = \begin{bmatrix} 0.0 & 0.5 & 0.0 & 0.5 & 0.0 \\ 0.0 & 0.0 & 0.5 & 0.5 & 0.0 \\ 1.0 & 0.0 & 0.0 & 0.0 & 0.0 \\ 0.0 & 0.0 & 0.0 & 0.0 & 1.0 \\ 0.0 & 0.0 & 1.0 & 0.0 & 0.0 \end{bmatrix}$$

$$\boldsymbol{\pi}^0 = \begin{bmatrix} 1.0 & 0.0 & 0.0 & 0.0 & 0.0 \\ \boldsymbol{\pi}^1 = \begin{bmatrix} 0.0 & 0.5 & 0.0 & 0.5 & 0.0 \end{bmatrix}$$

$$\boldsymbol{\pi}^2 = \begin{bmatrix} 0.0 & 0.5 & 0.0 & 0.5 & 0.0 \end{bmatrix}$$

$$\boldsymbol{\pi}^3 = \begin{bmatrix} 0.25 & 0.0 & 0.5 & 0.0 & 0.25 \end{bmatrix}$$

$$\boldsymbol{\pi}^4 = \begin{bmatrix} 0.5 & 0.125 & 0.25 & 0.125 & 0 \end{bmatrix}$$

$$\boldsymbol{\pi}^5 = \begin{bmatrix} 0.25 & 0.025 & 0.0625 & 0.3125 & 0.125 \end{bmatrix}$$

$$\boldsymbol{\pi}^5 = \begin{bmatrix} 0.25 & 0.125 & 0.25 & 0.1875 & 0.1875 \end{bmatrix}$$

## Computing $\pi$ (Method 1)

- Stationary state distribution is the limit distribution
- Idea: Compute k-step state probabilities  $\pi^k$  until they converge

#### • Power (iteration) method

- select arbitrary initial state probability distribution  $\pi^0$
- compute  $\pi^k = \pi^{k-1} \mathbf{P}$  until they converge (e.g.,  $|\pi^k \pi^{k-1}| < \varepsilon$ )
- report last  $\pi^k$  as stationary state distribution  $\pi$

• Power (iteration) method basically **simulates the Markov chain** and is the **method of choice in practice** when dealing with huge state spaces, exploiting that **matrix-vector multiplication** is **easy to parallelize** 

### Computing $\pi$ (Method 2)

• Stationary state distribution  $\pi$  fulfills  $\pi = \pi P$ , which can be cast into a system of linear equations

$$\mathbf{P} = \begin{bmatrix} 0.0 & 0.5 & 0.0 & 0.5 & 0.0 \\ 0.0 & 0.0 & 0.5 & 0.5 & 0.0 \\ 1.0 & 0.0 & 0.0 & 0.0 & 0.0 \\ 0.0 & 0.0 & 0.0 & 0.0 & 1.0 \\ 0.0 & 0.0 & 1.0 & 0.0 & 0.0 \end{bmatrix}$$

$$\pi_1 = 0.0 \times \pi_1 + 0.0 \times \pi_2 + 1.0 \times \pi_3 + 0.0 \times \pi_4 + 0.0 \times \pi_5$$
 $\pi_2 = 0.5 \times \pi_1 + 0.0 \times \pi_2 + 0.0 \times \pi_3 + 0.0 \times \pi_4 + 0.0 \times \pi_5$ 
 $\pi_3 = 0.0 \times \pi_1 + 0.5 \times \pi_2 + 0.0 \times \pi_3 + 0.0 \times \pi_4 + 1.0 \times \pi_5$ 
 $\pi_4 = 0.5 \times \pi_1 + 0.5 \times \pi_2 + 0.0 \times \pi_3 + 0.0 \times \pi_4 + 0.0 \times \pi_5$ 
 $\pi_5 = 0.0 \times \pi_1 + 0.0 \times \pi_2 + 0.0 \times \pi_3 + 1.0 \times \pi_4 + 0.0 \times \pi_5$ 
 $1 = 1.0 \times \pi_1 + 1.0 \times \pi_2 + 1.0 \times \pi_3 + 1.0 \times \pi_4 + 1.0 \times \pi_5$ 

• Solutions can be found, e.g., using Gauss elimination

### Computing $\pi$ (Method 3)

• Stationary state probability distribution  $\pi$  is the left eigenvector of the transition probability matrix P for the eigenvalue  $\lambda = 1$ 

$$\pi P = \lambda \pi$$

• Can be computed using the characteristic polynomial

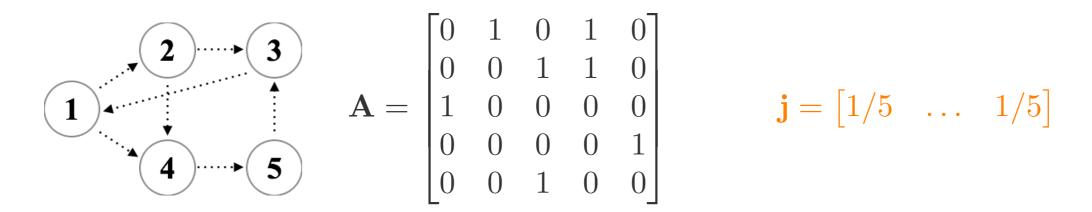
$$(\mathbf{P} - \lambda \mathbf{I}) \; \boldsymbol{\pi} = \mathbf{0}$$

- Random surfer model
  - follows a uniform random outgoing link with probability (1- $\varepsilon$ )
  - jumps to a uniform random web page with probability  $\varepsilon$
- Let A be the adjacency matrix of the Web graph, matrix T captures following of a uniform random outgoing link

$$\mathbf{T}_{ij} = \begin{cases} 1/out(i) & : & (i,j) \in E \\ 0 & : & \text{otherwise} \end{cases}$$

- Random surfer model
  - follows a uniform random outgoing link with probability (1- $\varepsilon$ )
  - jumps to a uniform random web page with probability  $\varepsilon$
- Vector **j** captures jumping to a uniform random web page

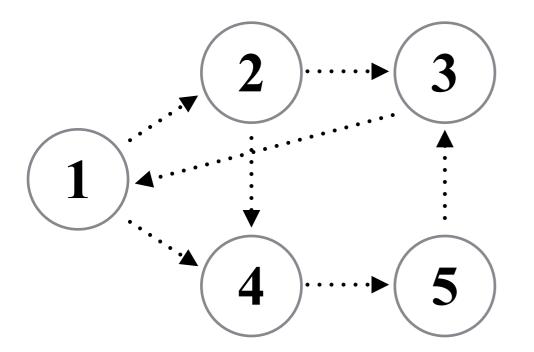
$$\mathbf{j}_i = 1/|V|$$



• Transition probability matrix of Markov chain then obtained as

$$\mathbf{P} = (1 - \epsilon) \mathbf{T} + \epsilon \begin{bmatrix} 1 & \dots & 1 \end{bmatrix}^T \mathbf{j}$$

• With  $\varepsilon = 0.15$  we obtain



$$\mathbf{P} = \begin{bmatrix} 0.030 & 0.455 & 0.030 & 0.455 & 0.030 \\ 0.030 & 0.030 & 0.455 & 0.455 & 0.030 \\ 0.880 & 0.030 & 0.030 & 0.030 & 0.030 \\ 0.030 & 0.030 & 0.030 & 0.030 & 0.880 \\ 0.030 & 0.030 & 0.880 & 0.030 & 0.030 \end{bmatrix}$$

$$\boldsymbol{\pi} = \begin{bmatrix} 0.24079 & 0.13234 & 0.24799 & 0.18858 & 0.19029 \end{bmatrix}$$

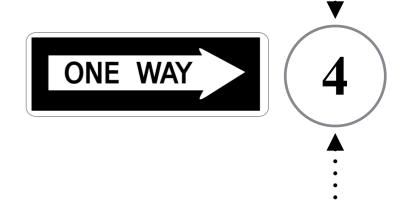
• Transition probability matrix of Markov chain then obtained as

$$\mathbf{P} = (1 - \epsilon) \mathbf{T} + \epsilon \left[1 \dots 1\right]^{T} \mathbf{j}$$

$$\pi_{i} = (1 - \epsilon) \sum_{(j,i) \in E} \frac{\pi_{j}}{out(j)} + \frac{\epsilon}{|V|}$$

- We need to deal with dangling nodes (having out-degree zero)
  - Re-normalize  $\pi^k$  such that  $|\pi^k| = 1$  after every iteration of power method
  - Make P truly right stochastic by defining matrix T as

$$\mathbf{T}_{ij} = \begin{cases} 1/out(i) & : & (i,j) \in E \\ 1/|V| & : & out(i) = 0 \\ 0 & : & \text{otherwise} \end{cases}$$



### PageRank as a Markov Chain (Is It Ergodic?)

- Markov chain defined by transition probability matrix T is
  - finite (only finite number of web pages)
  - time-homogeneous (by design)
  - **irreducible** (random surfer can jump from every state *i* to every state *j*)
  - **positive recurrent** (random surfer can "jump up" on state *i*)
  - aperiodic (period of every state is 1 because of "jump up" on state i)

...it is thus **ergodic** and unique stationary state probabilities  $\pi$  exist

• Random jump is essential to make the Markov chain ergodic

### PageRank & Queries

- Random jump probability typically set as  $\varepsilon = 0.15$  (i.e., random surfer follows on average about seven links in a row)
- PageRank determines a static global ranking of web pages, is query-independent, and orthogonal to textual relevance
- Combination of PageRank score and retrieval models, e.g., as
  - linear combination of cosine similarity and PageRank score

$$\alpha \times sim(q,d) + (1-\alpha) \times pr(d)$$

• document prior in a query-likelihood language model

$$P(q|d) \times P(d)$$

• together with many other features in machine-learned ranking model

### Summary of IV.2

#### Markov chains

as a kind of stochastic process useful to describe random walks

#### Stationary state distribution

is guaranteed to exist if the Markov chain is finite and ergodic can be computed using (i) power iteration (ii) solving a system of linear equations or (iii) determining an eigenvector of a matrix

#### • PageRank

as Google's initial secret of success is based on a random surfer model can be described as a finite and ergodic Markov chain yields a static query-independent importance score

#### Additional Literature for IV.2

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