Universität des
Saarlandes
FR Informatik


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# Tutorials for "Automated Reasoning" <br> Exercise sheet 14 

Exercise 14.1: (4 P)
Use superposition to show that the following set of (implicitly universally quantified) clauses is not satisfiable:

$$
\begin{gathered}
f(a, x) \approx x \\
x \approx a \vee x \approx g(a) \\
x \not \approx g(x) \\
f(a, g(a)) \approx g(b) \\
b \not \approx a
\end{gathered}
$$

Use the LPO with precedence $f>g>a>b$. Compute only inferences that are required according to the ordering restrictions of the superposition calculus.

Exercise 14.2: (3 P)
Let $\succ$ be the LPO with the precedence $a>b>c>g>f$. Order the following ground equational clauses according to the clause ordering $\succ_{C}$ defined in the lecture.

$$
\begin{gathered}
f(a, c) \not \approx g(f(c, c)) \vee a \approx g(c) \\
a \not \approx b \vee f(f(a, a), f(a, b)) \not \approx f(g(a), g(b)) \\
g(a) \approx g(b) \vee g(f(c, a)) \approx f(a, c) \\
a \approx b \vee g(g(c)) \not \approx f(c, b) \\
g(c) \approx f(a, b)
\end{gathered}
$$

Exercise 14.3: (2 P)
Let $\succ$ be a reduction ordering that is total on ground terms and $E$ be a set of equations. Suppose that equations in $E$ can be oriented (i.e. for every $t \approx s \in E$ exactly one of $t \succ s$ or
$s \succ t$ holds). Show that the set of semi-critical pairs of $E$ with respect to $\succ$ is the same as the set of critical pairs of rewrite system $R$ where $R=\{s \rightarrow t \mid s \succ t \wedge s \approx t \in E\}$.

Exercise 14.4: (3 P)
Prove Theorem 6.4.

Exercise 14.5: (2 Bonus Points)
Is there any reduction ordering that is total on $T_{\Sigma}(X)$ ? Give a proof or counterexample.

Submit your solution in lecture hall 001 during the lecture on July 23. Please write your name and the date of your tutorial group on your solution.

Note: Joint solutions are not permitted (work in groups is encouraged).

