



Effizientes Routing in P2P Netzwerken

Chord: A Scalable Peer-to-peer Lookup Protocol for Internet Applications

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Outline

- Short introduction to Chord
- Experiments



Introduction to Chord

Data addressed by a Hashtable

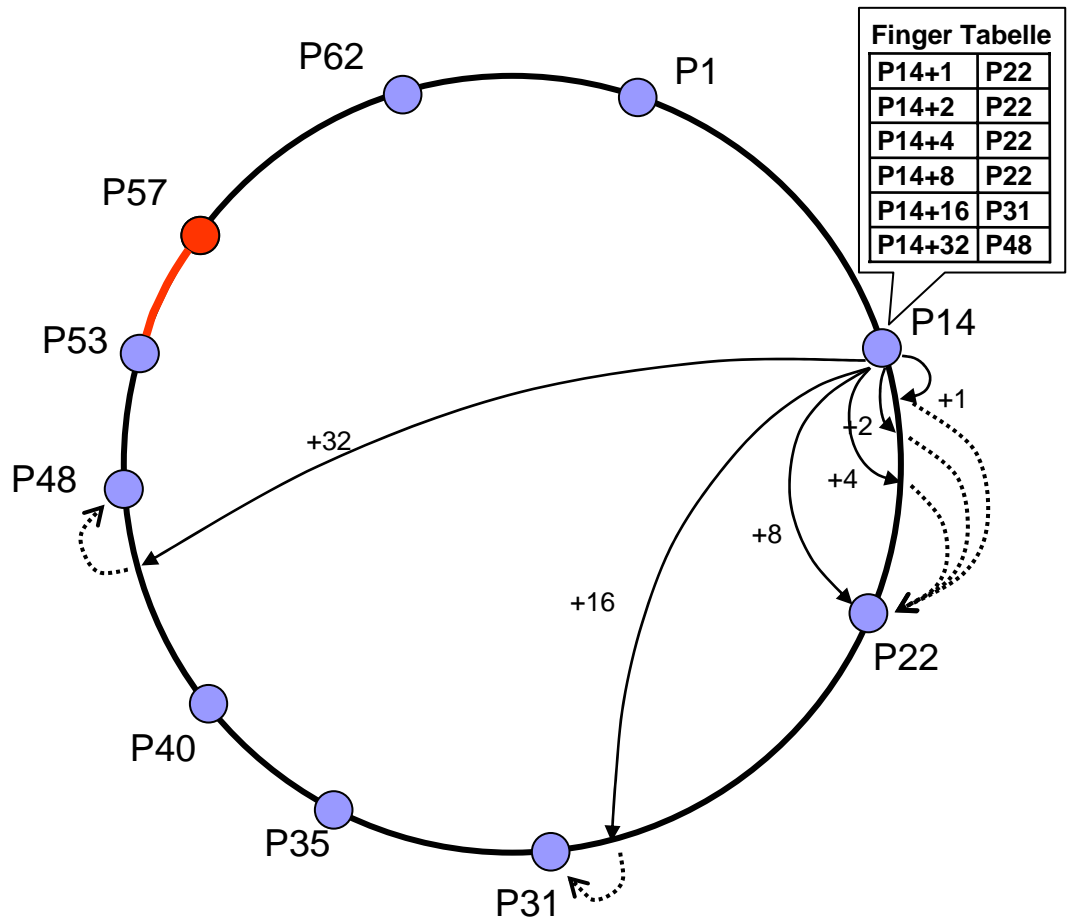
- calculating hashvalues for keys
- storing reference to data identified by hashvalues

A distributed hashtable (DHT)

- partitioned
- distributed over peers

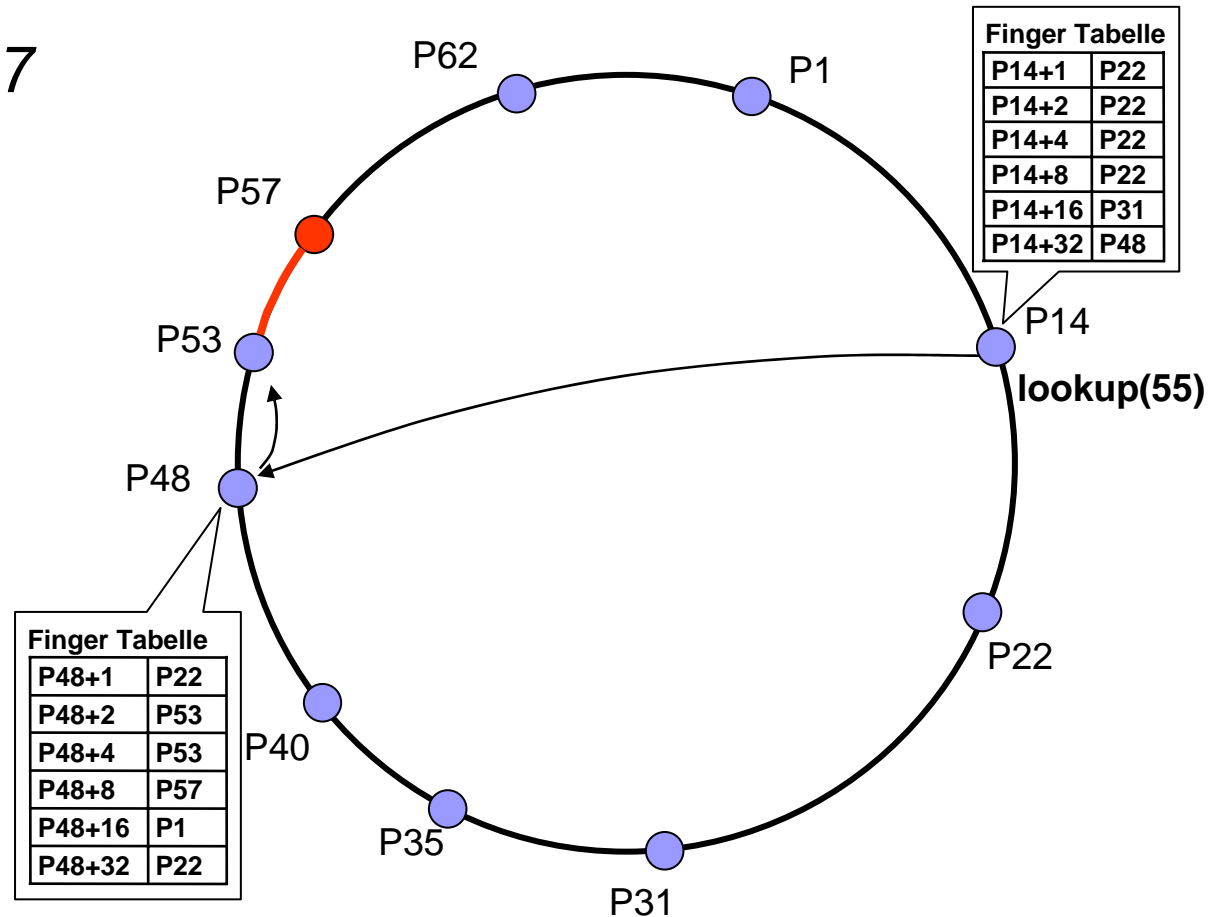
Introduction to Chord

- $lookup(key) = x$,
first node x
following key with
 $x.ID \geq key$
- $finger[i] =$
 $lookup(x.ID + 2^{i-1})$
- periodic update of
neighbours and
fingers



Introduction to Chord

- $lookup(55) = P57$





Experimental Setup

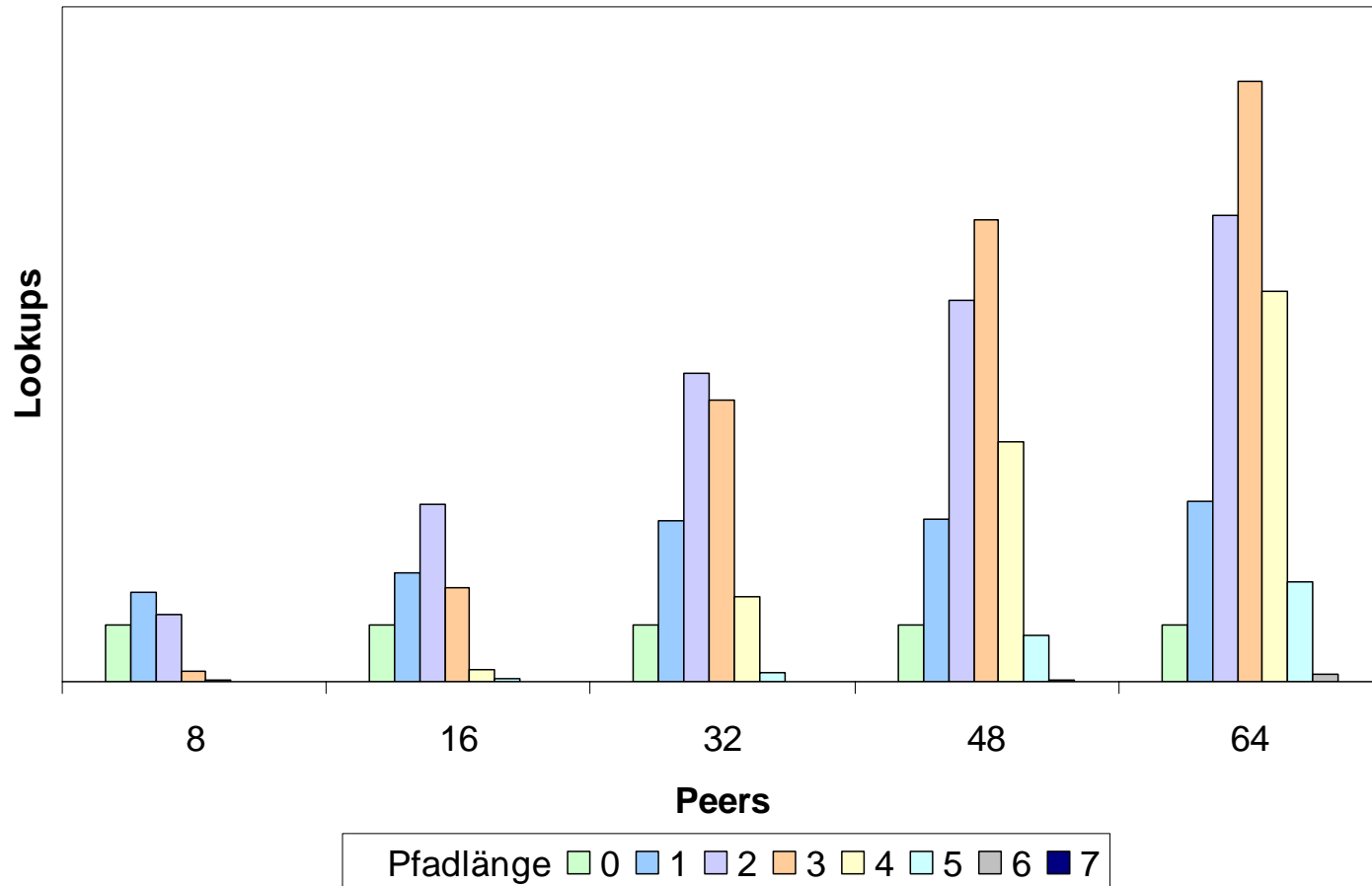
- Java-Implementation of Chord
 - ca. 4500 lines
- Experiments
 - one Java-VM
 - Network device: localhost



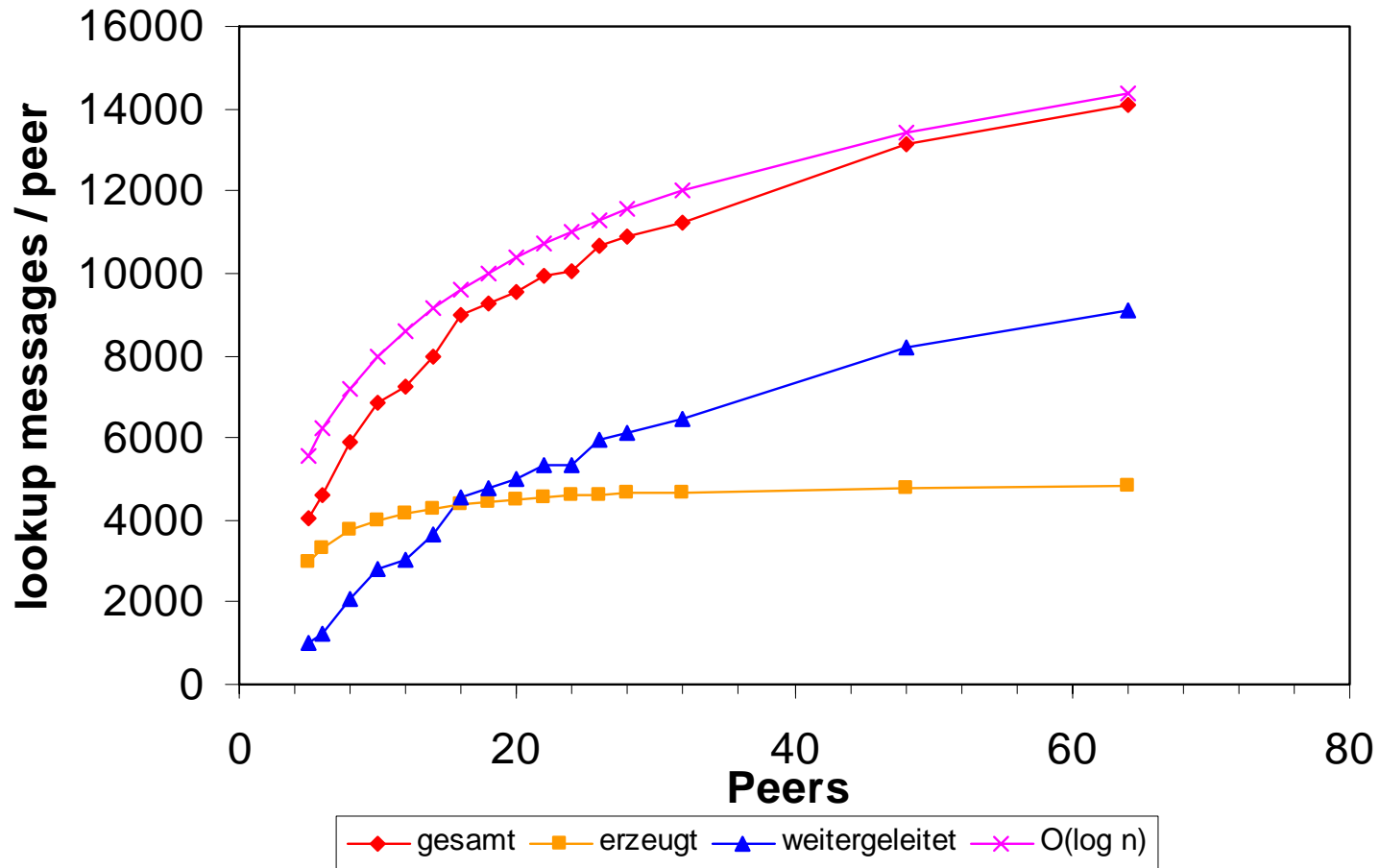
Experimental Setup

- Path length
- Number of routing messages
- Lookup duration & timeouts
- Bandwidth used
 - Lookups
 - Stabilization

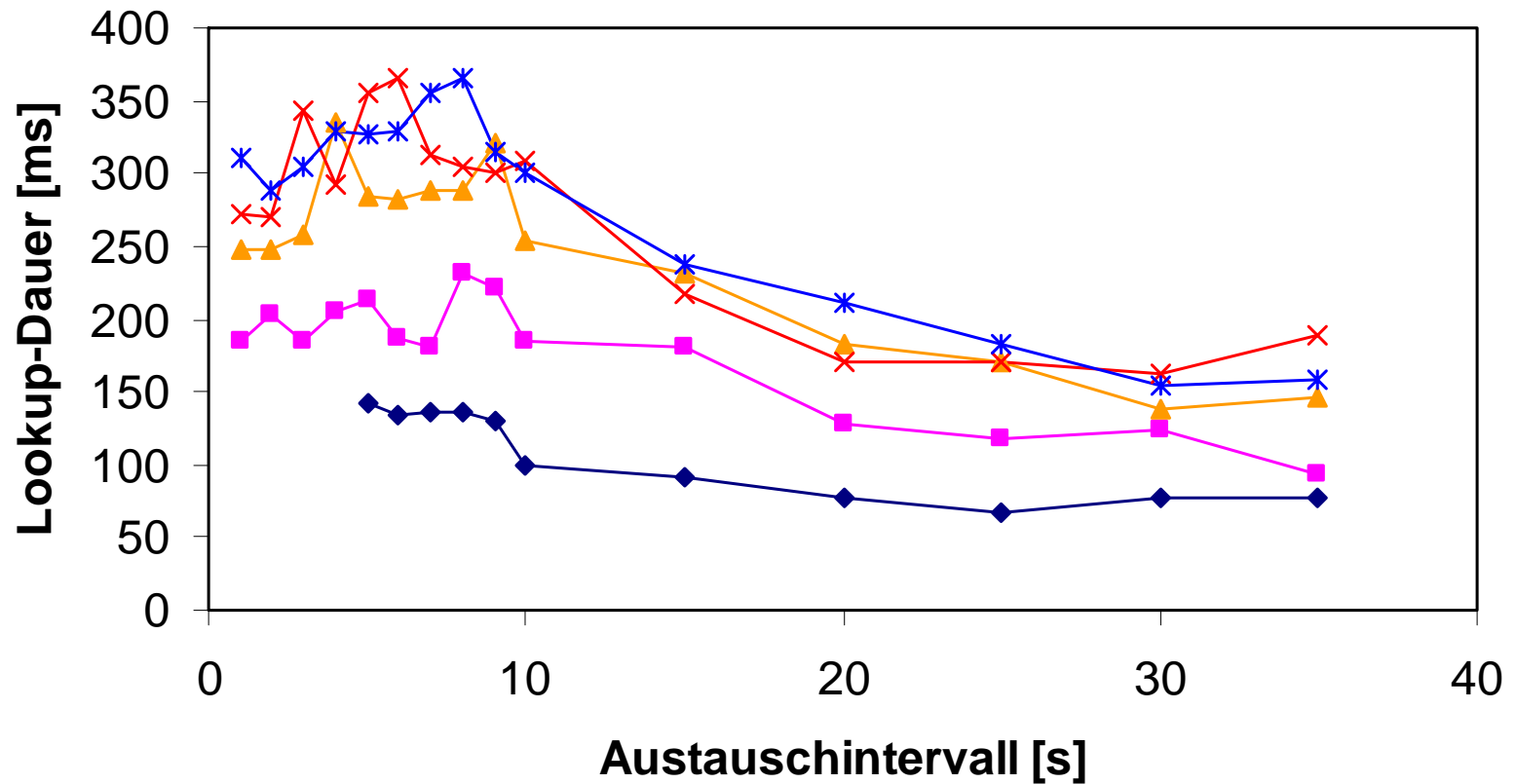
Path Lengths



Routing Messages

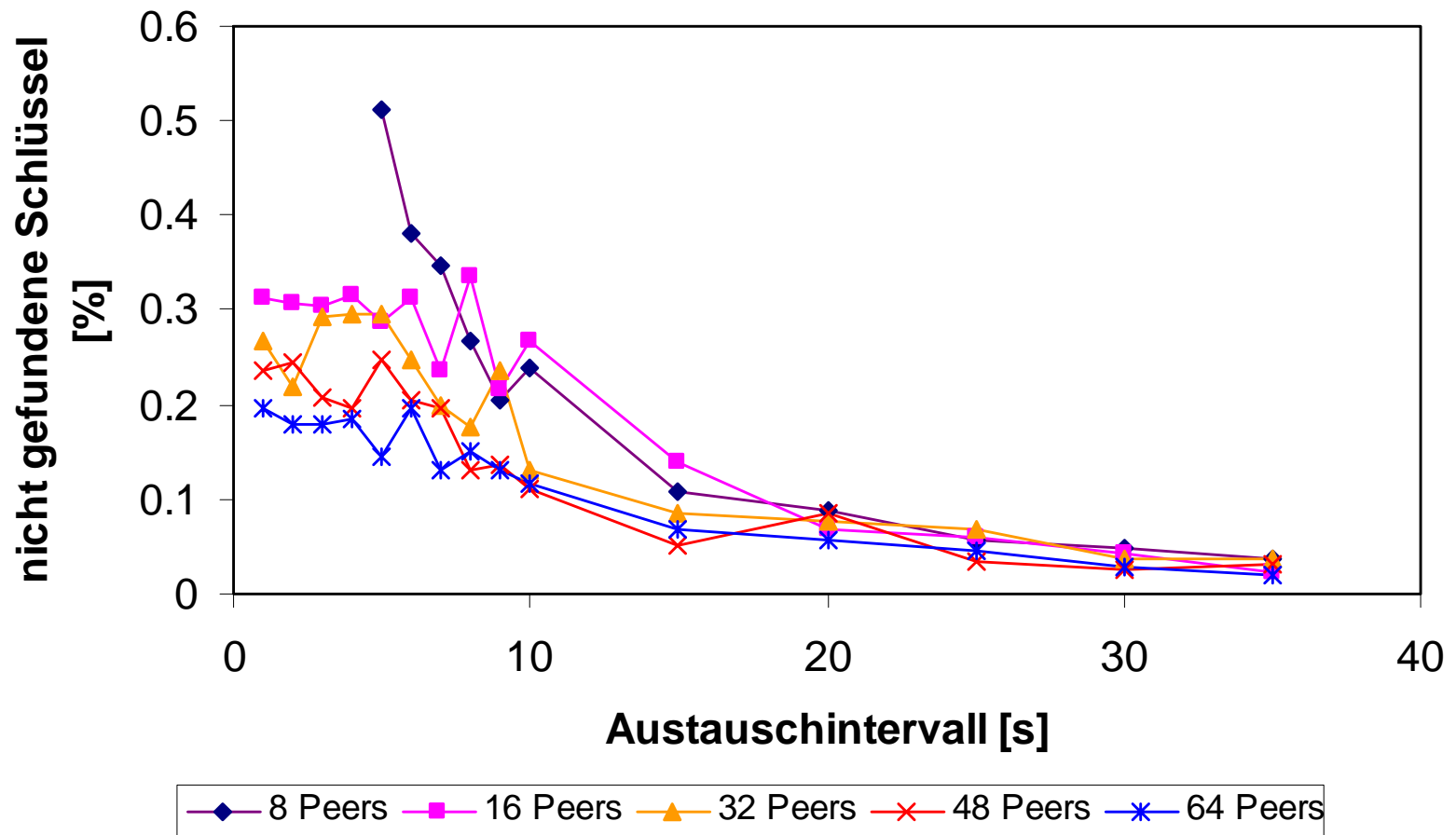


Lookup Duration

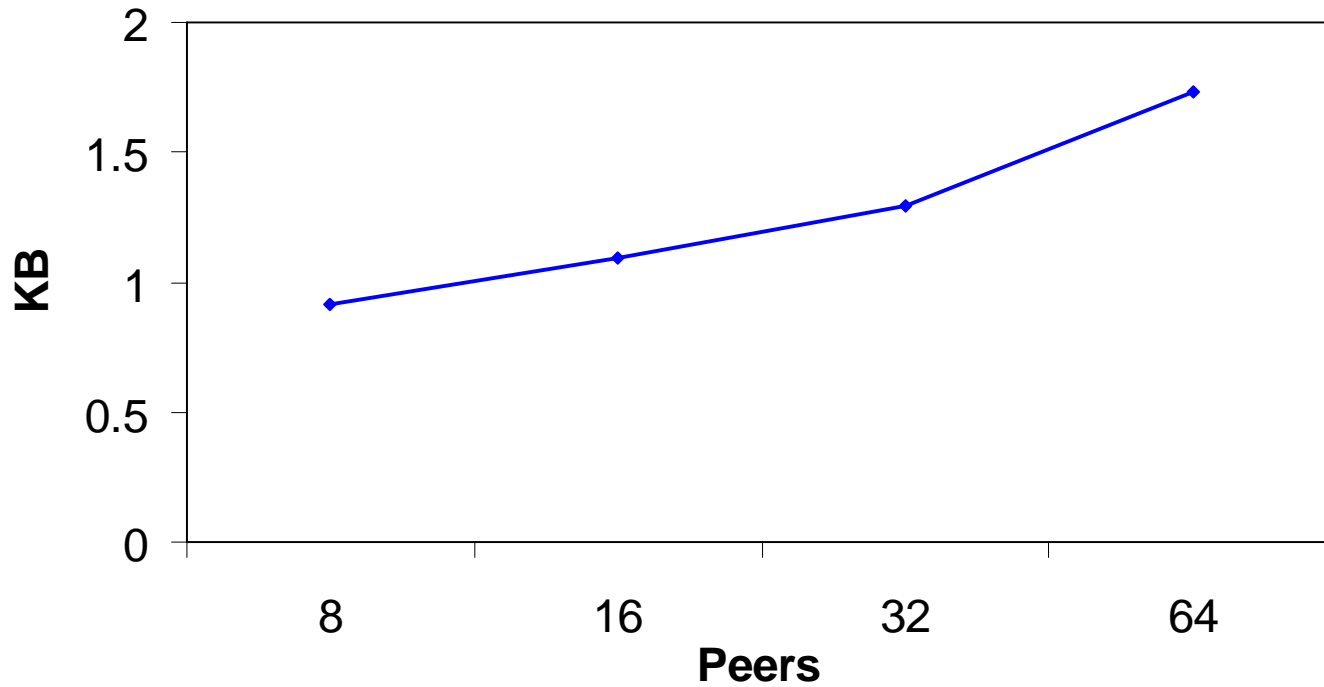


◆ 8 Peers ■ 16 Peers ▲ 32 Peers × 48 Peers * 64 Peers

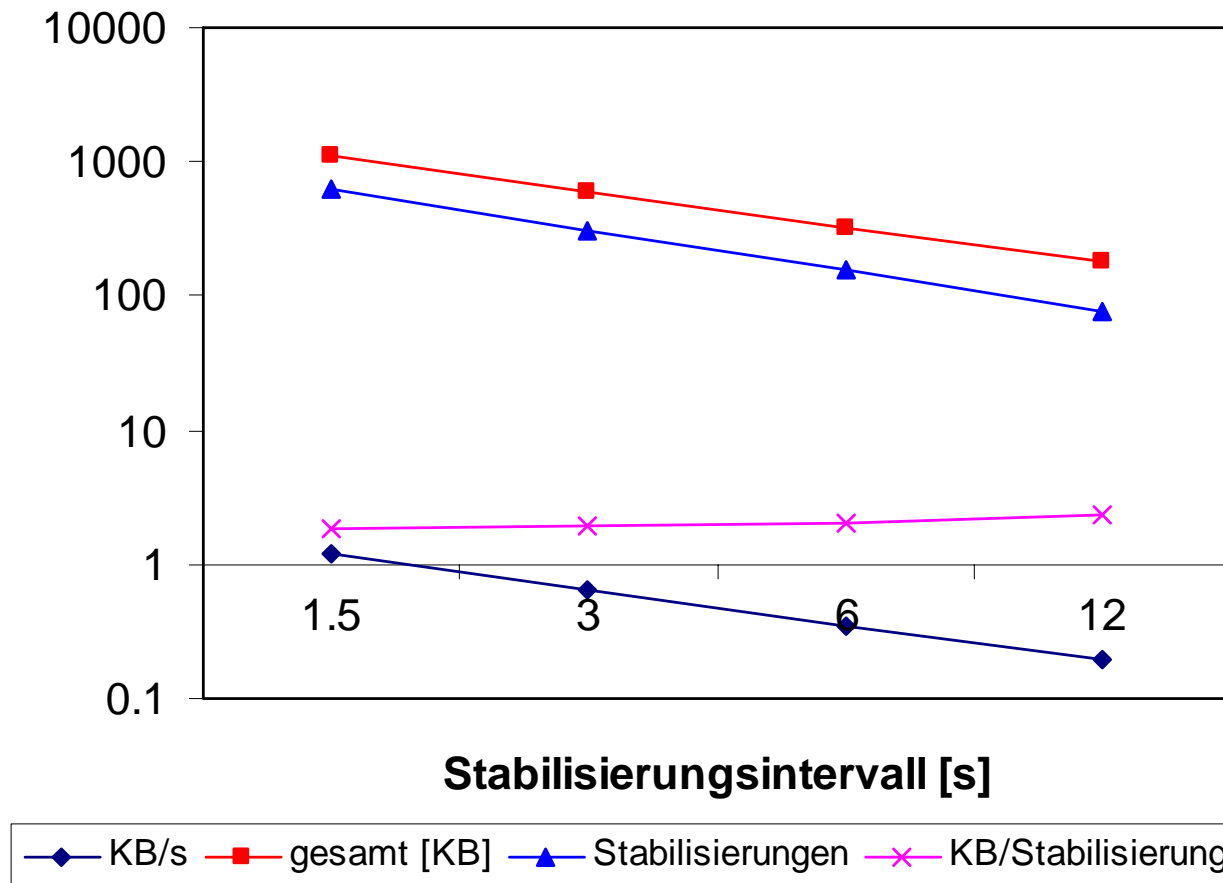
Lookup Timeouts



Data per Lookup



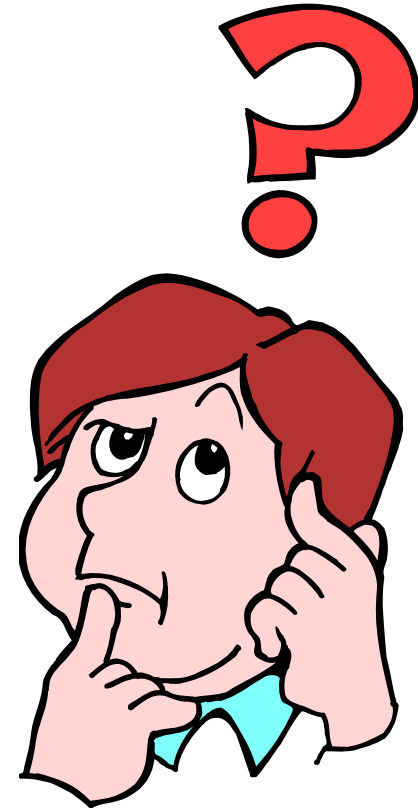
Bandwidth Usage - Stabilization



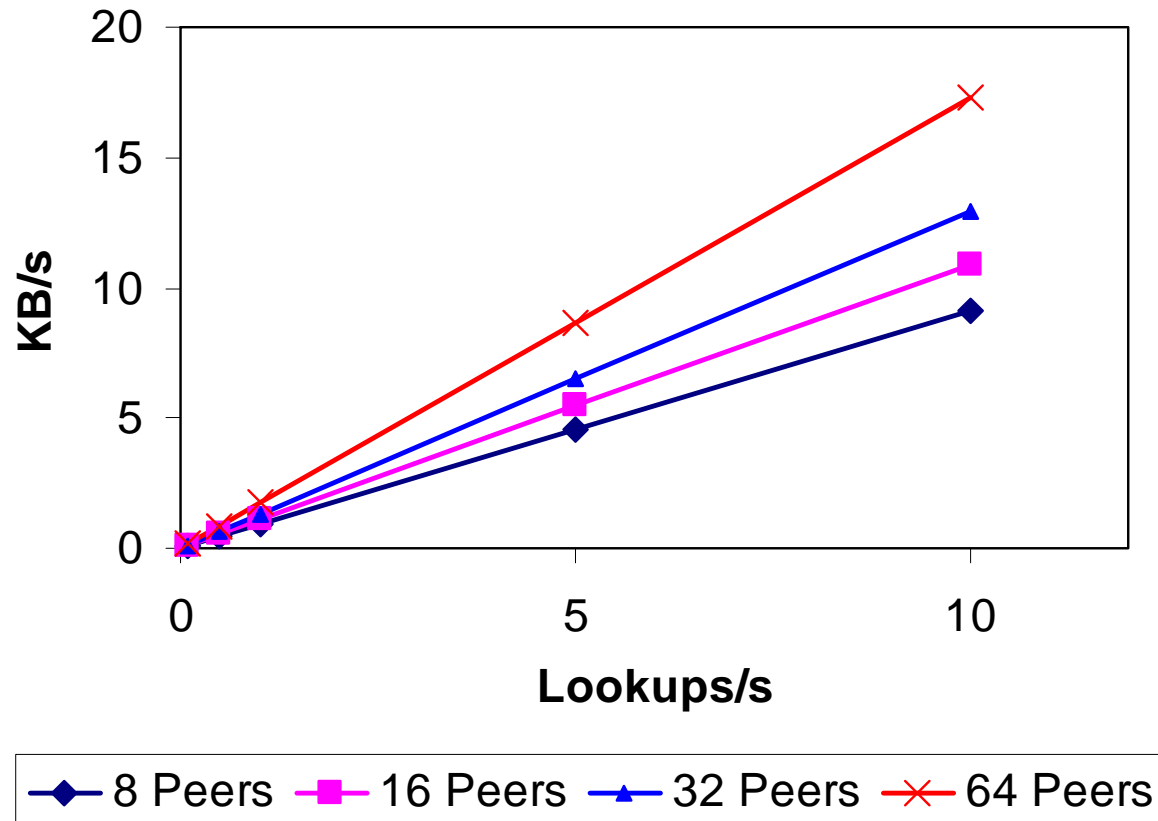
Questions & Discussion

Thank you for attention!

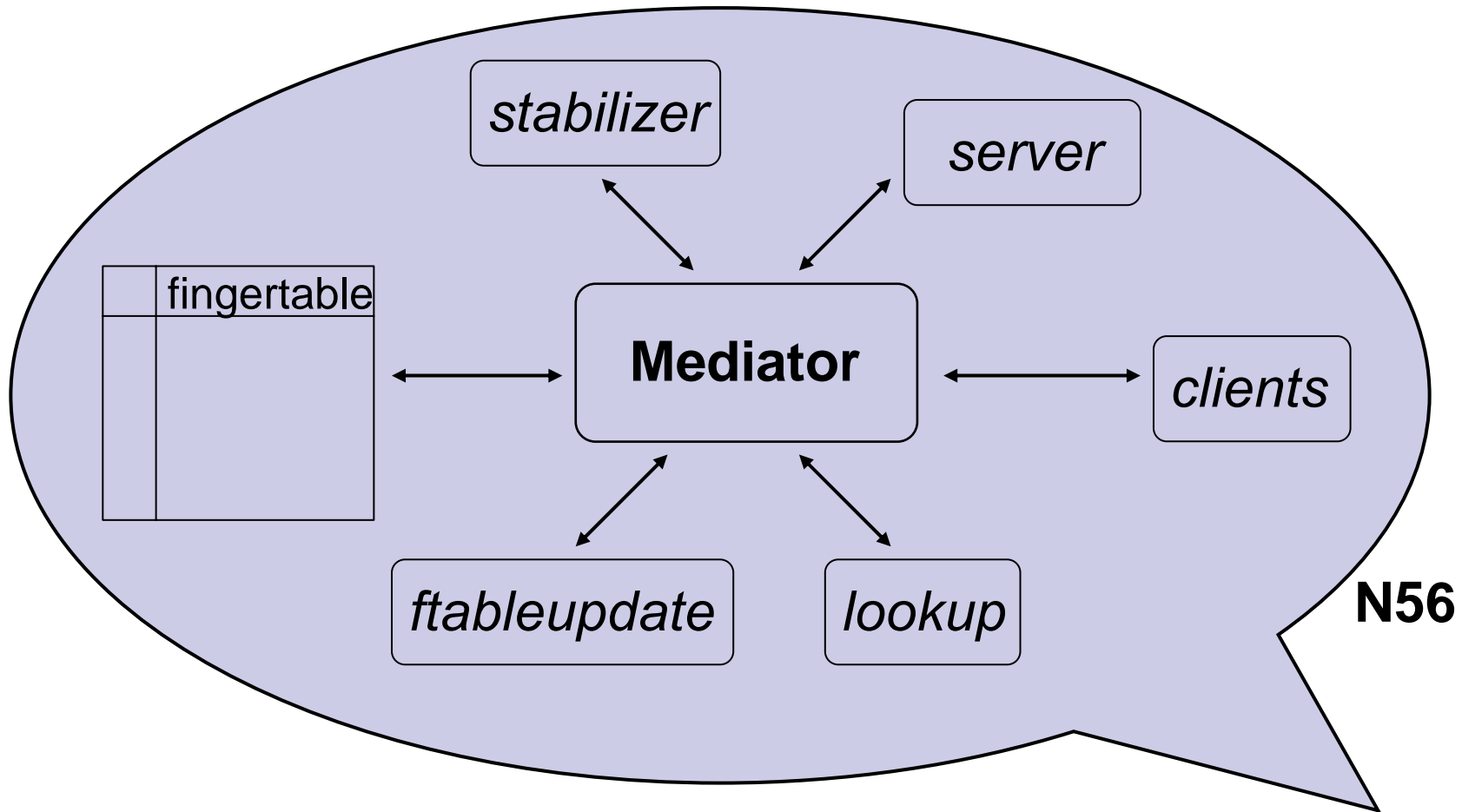
Questions?



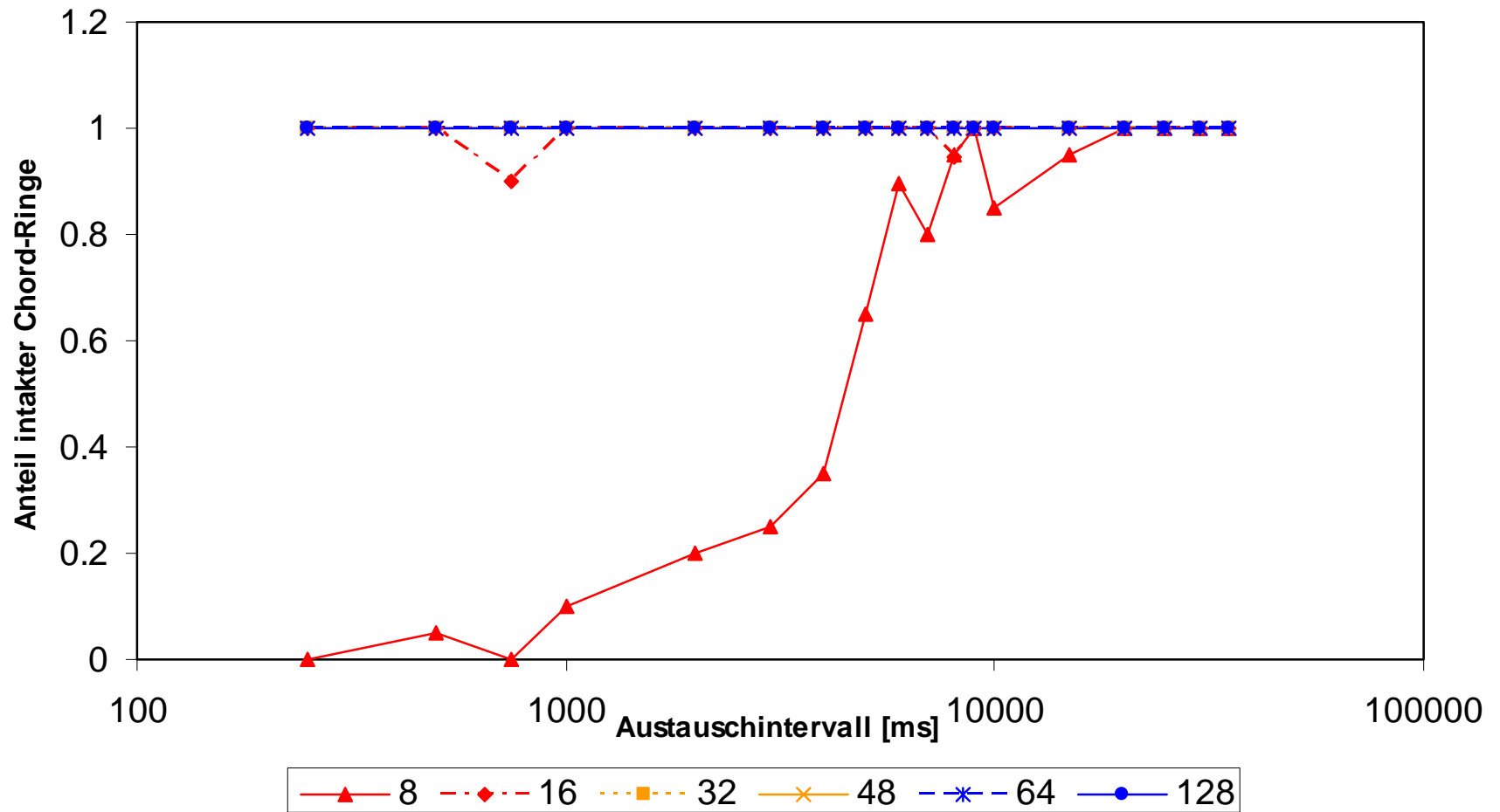
Bandwidth Usage - Lookups



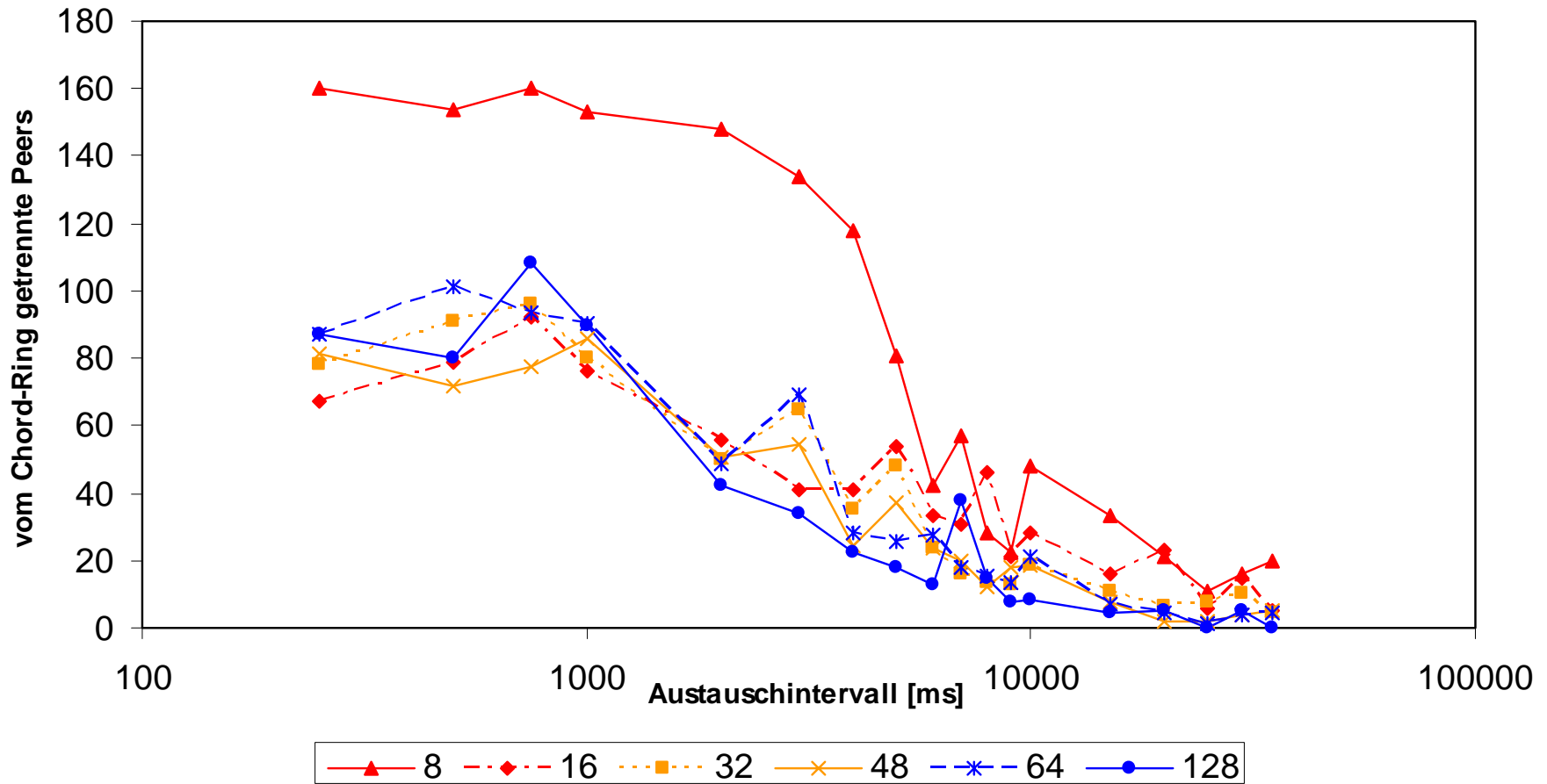
Implementation Overview

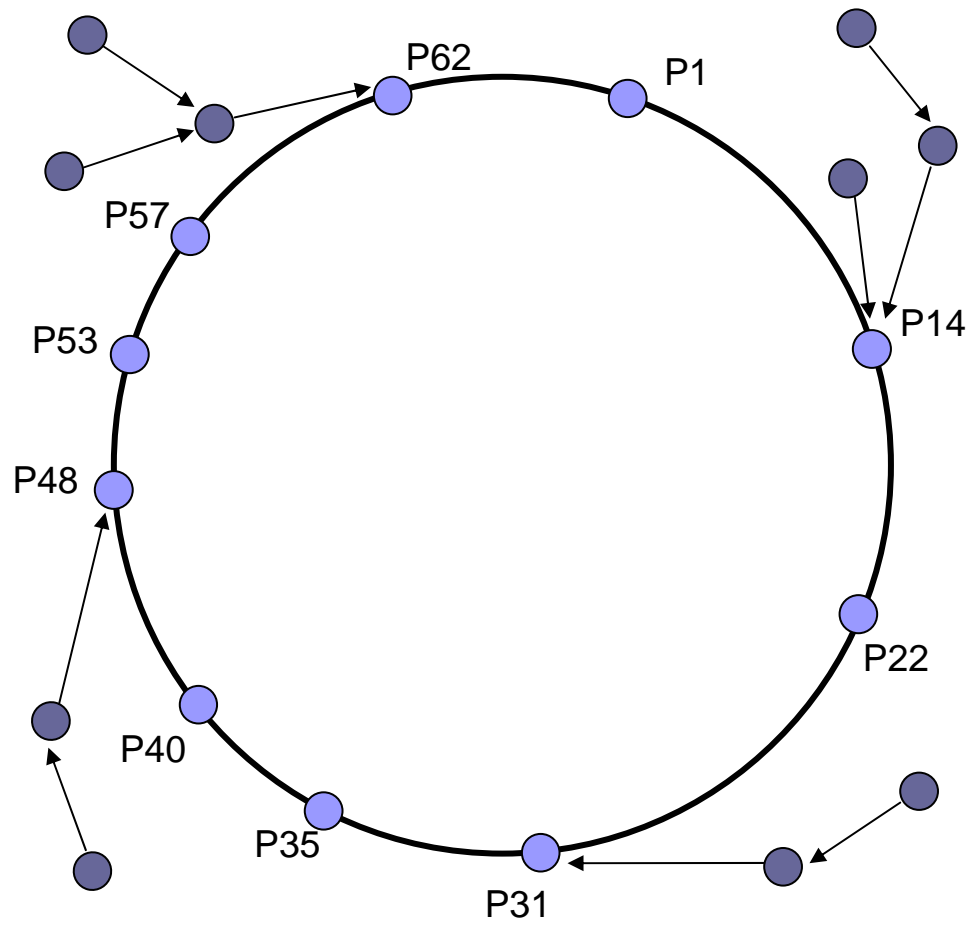


Robustness



Robustness (2)





The Chord algorithm – Node joins and stabilization

Stabilization protocol for a node x :

■ $x.stabilize()$:

- ask successor y for its predecessor p
- if $p \in (x; y]$ then p is x 's new successor

■ $x.notify()$:

- notify x 's successor p of x 's existence
- notified node may change predecessor to x