Chord: A Scalable Peer-to-peer Lookup Protocol for Internet Applications

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- I. Introduction
- II. The Chord Protocol and how it works
- III. Disadvantages and Improvements
- IV. Applications using Chord
- V. Summary

I. Introduction

Problem:

Decentralized network with several peers (clients) How to find specific peer that hosts desired data within this network?







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What is Chord?

- Efficient
- Simple
- Provable Performant
- Provable Correct
- Scalable
- Supporting only one operation: mapping a given key onto a node

lookup protocol

Consistent Hashing:



A first simple (but slow!) lookup algorithm



=>#messages linear in #nodes

Protocol Improvements

Introduce a routing table *m*: number of bits in the key/node identifiers *n*: ID of a node ith entry: the closest node that succeeds node *N* with ID *n* by at least 2ⁱ⁻¹ mod 2^m

 \Rightarrow Finger table could contain up to *m* entries

Each node has a successor list of its *r* successing nodes

Finger Table for scalable node localisation: not enough information to find node directly but nodes in direct neighbourhood



More efficient algorithm using finger tables

//ask node n to find the successor //of id n.find_successor(id) if (id in (n,successor]) return successor else n' = closest_preceding_node(id);

//search the local table for the
//highest predecessor of id
n.closest_preceding_node(id)
for i=m downto 1 do
if (finger[i] in (n,id))
return finger[i];
return n;

return *n*'.find_successor(*id*);





Node joins and stabilization process

Lookups are expected to be correct (though the set of participating nodes can change!)

- ⇒ Important: Any node s successor pointer is always up to date.
- Stabilization protocol running periodically in the background and updating Chord ´s finger tables and successor pointers

Stabilize()

if *n* is not its *successor*'s *predecessor* it changes its *successor*



Notify()

 N_z changes predecessor to N_x because predecessor = *nil* N_z changes predecessor to N_y because N_y between N_x and $N_{z'}$.





- a) Initial state
- b) node 26 enters the system & finds its successor
- c) Stabilize procedure updates successor of n21 to n26



Node failures

 It is provable that even if half of the nodes in the network fail, lookups are executed correctly!



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III. Disadvantages and Improvements

Why is Chord itself no search engine?

- Protocols are NO applications
- Upper layers have to coordinate purposes
- Chord only supports "exact match", cannot handle queries similar to one or more keys

Suitability for use in search engines

 Chord basically not suitable because specific information about desired data is needed to be able to compute the query key shashvalue

There exist several proposals for modifications to be able to find what is wanted without prior knowledge (meta data search extension)

III. Disadvantages and Improvements

Chord ´s disadvantages

- asymmetric lookup => lookup from node n to node p could take a different number of hops than vice versa
- => In huge Chord Rings lookup to one of near preceding nodes routed clockwise over almost complete ring.
- Approach for Improvement: **S-Chord** (->in-place notification of entry changes)
- Symmetric Structure:
 - routing entry symmetry "if p points to n then n points to p"
 - routing cost symmetrie lookup paths between to nodes very likely are equal though they are not the same! (=>no routing symmetry)

- finger table symmetry



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IV. Applications using Chord

DNS with Chord

- Host names hashed to keys, corresponding IP addresses are values
- Decentralized (no use for servers)
- Routing information and host changes can be updated dynamically

Cooperative mirroring

- Multiple providers of the same content
- Load is spread evenly over all hosts by mapping data blocks onto hosts



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V. Summary

- Simple, scalable, provable correct and performant
- Any lookup requires O(log N) messages in a N node network
- Even in unstable systems lookups are correct
- Arbitrary implementation provided for developers
 => extendable in applications



Thank you for your attention!