

P-Grid: A self-organizing access structure for P2P information systems

Karl Aberer

Department of Communication
Systems

Swiss Federal Institute of Technology
(EPFL)

Xiaoli Zhang

lilie0927@yahoo.de

Overview

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2. System Model and Access Structure
3. P-Grid Distributed Search Structure
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6. Conclusions

1. Introduction

- Peer-To-Peer systems
 - are driving a major paradigm shift in the era of genuinely distributed computing
 - Every node of the system plays the role of a client and a server

P-Grid (Peer-Grid)

- The resulting distributed access structure we call P-Grid.
- is a next generation peer-to-peer platform for distributed information management file-sharing
- is a virtual binary search tree that distributes replication over a community of peers and supports efficient search

P-Grid (Peer-Grid)

- P-Grid's most important properties are:
 - complete decentralization
 - self-organization
 - decentralized load balancing
 - efficient search

P-Grid (Peer-Grid)

- P-Grid's most important applications are:
 - data management functionalities (update)
 - management of dynamic IP addresses and identities

2. System Model and Access Structure

- **Assume:** address $r \in \text{ADDR}$ we define $\text{peer}(r) = a$ iff $\text{addr}(a) = r$ for $a \in P$
- Peers are online:
- with a probability online: $P \rightarrow [0,1]$
- can be reached reliably using the underlying communication infrastructure by means of their address

Access Structure

- Now we define the access structure (not formally):
 - The search space is partitioned into intervals of the form $I(k), k \in K$.
 - Every peer takes over responsibility for an interval $I(k)$.
 - The peers have the same prefix of length L , but a different value at position $L+1$.
 - A search can be started at each peer.

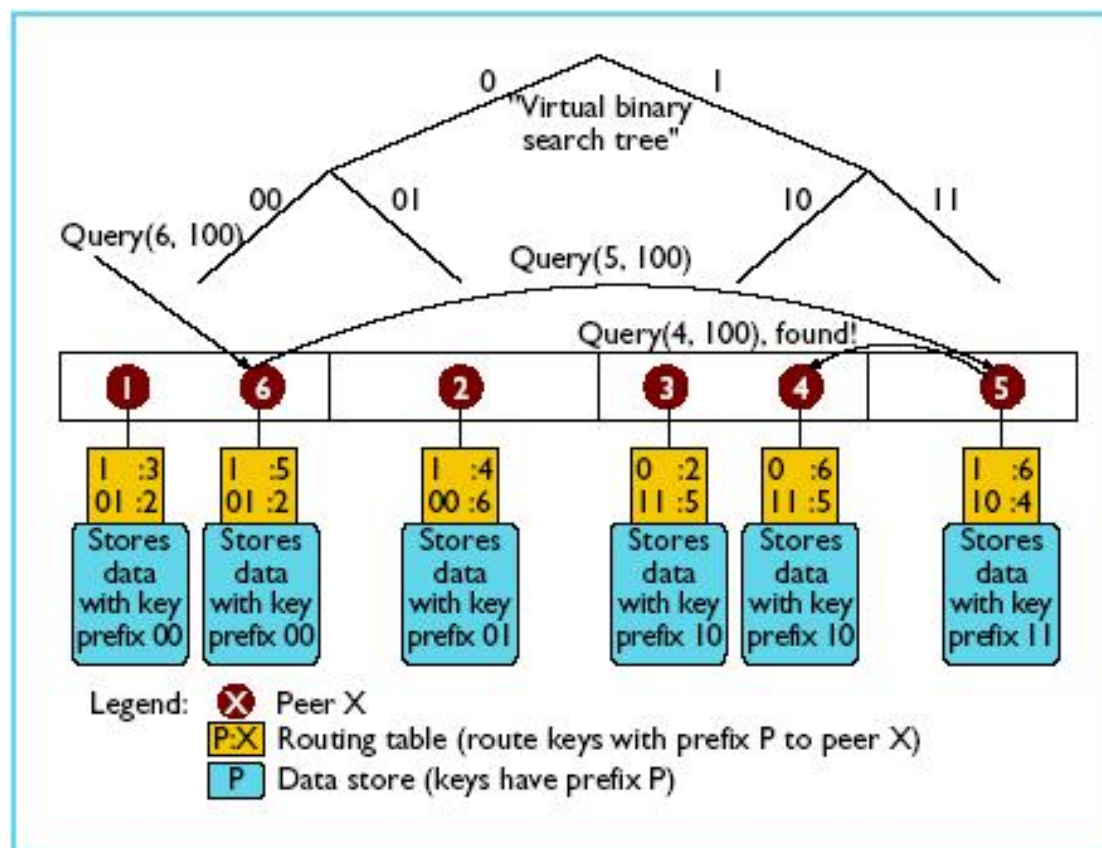
Access Structure

- Work in the following way
 - Every peer maintains a small fragment of the set of keys and a routing table
 - The routing tables are organized such that at different levels of granularity requests can be forwarded
 - Replication is used to increase robustness

Access Structure

- The data structure for peers of the P-Grid (formally):
 - Every peer $a \in P$ maintains a sequence $(p_1, R_1)(p_2, R_2) \dots (p_n, R_n)$, where $p_i \in \{0, 1\}$ and $R_i \subset \text{ADDR}$, $\text{path}(a) = p_1 \dots p_n$, $\text{prefix}(i, a) = p_1 \dots p_i$ for $0 \leq i \leq n$ and $\text{refs}(i, a) = R_i$.
 - ☆ In addition the number of references in R_i will be limited by a value [refax](#).
 - ☆ One can see that multiple peers can be responsible for the same key, we call these later also [replicas](#).
 - ☆ We bound the recursion depth up to which the exchange function is called by the value [recmax](#).

3.The P-Grid Search Structure



3.P-Grid Distributed Search Structure

- ☆ it is completely decentralized;
- ☆ all peers serve as entry points for search;
- ☆ interactions are strictly local;
- ☆ it uses randomized algorithms for access and search;
- probabilistic estimates of search request success can be given;
- search is robust against node failures; and it scales gracefully in the total number of nodes and data items.

4.P-Grid Construction

- No global control ,only by local interactions
- Meet randomly
- Initially, all peers are responsible for all the search keys.
- When two peers meet, they decide to split the search space into two parts and take over responsibility for one half each, also store the reference to other peer.
- The same happens whenever two peers meet, that are responsible for the same interval at the same level.

4.P-Grid Construction

- Namely, peers will meet
 - 1, their keys share a common prefix,
 - 2, their keys are in a prefix relationship.
- ☆ As many peers can be responsible for the same key the general problem is of how to find all those peers in case of an update. Different strategies are possible:
 - 1, depth first search
 - 2, breadth first search
 - 3, search with buddies

5. Simulation

- Implementation in Mathematica(www.wolfram.com)
- Questions:
 - 1, How many communications in terms of executing the exchange algorithm are required for building a P-Grid ?
 - 2, How balanced the P-Grid is with respect to the distribution of keys?
 - 3, How reliably updates can be performed?

5. Simulation

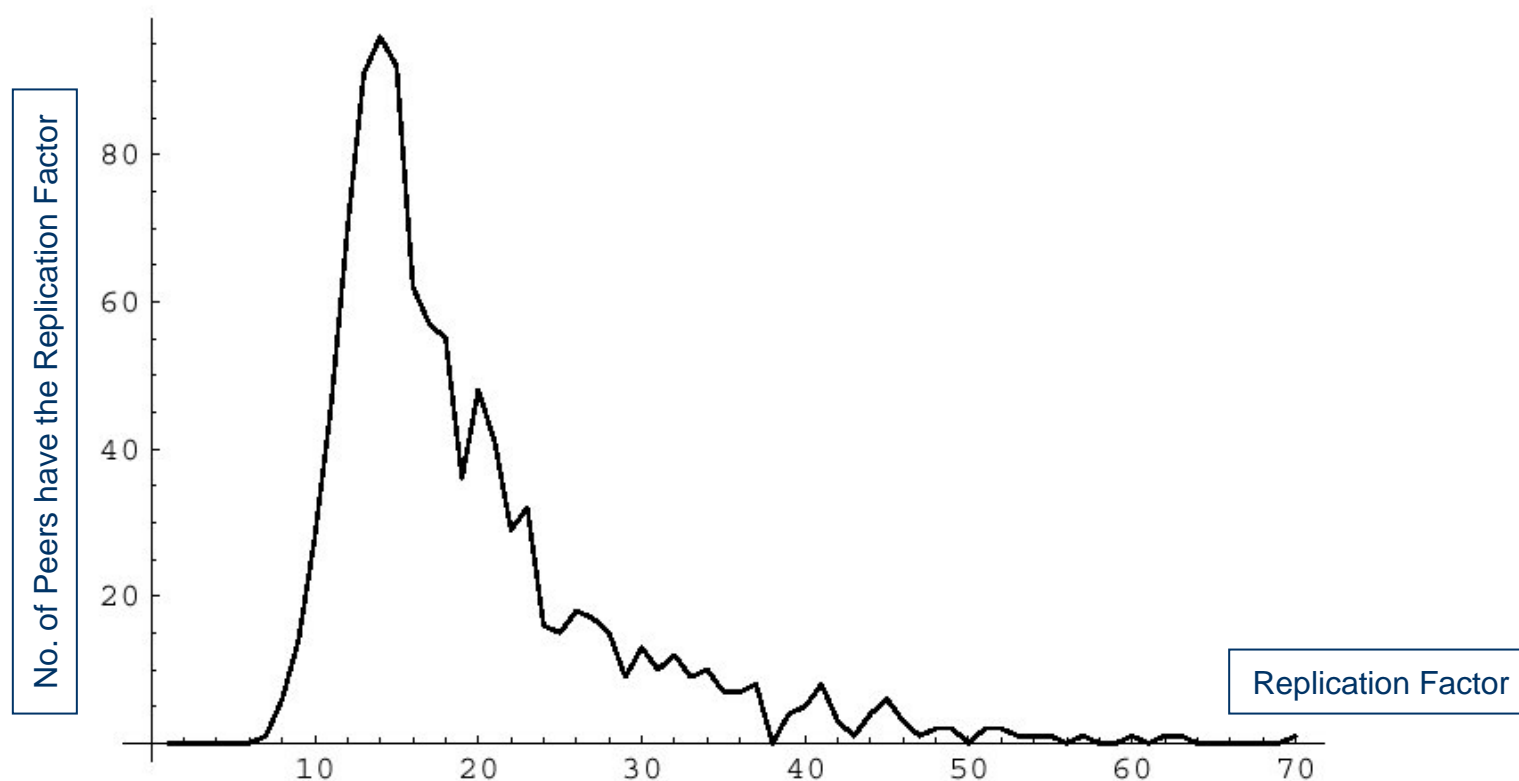
- Question 1

The results indicate that a linear relationship exists between the number of peers (N) and the total number of communications (e) needed in building the P-Grid.

N	recmax = 0		recmax=2	
	e	$\frac{e}{N}$	e	$\frac{e}{N}$
200	15942	79.71	4937	24.68
400	27632	69.08	10383	25.95
600	43435	72.39	15228	25.38
800	59212	74.01	18580	23.22
1000	74619	74.61	25162	25.16

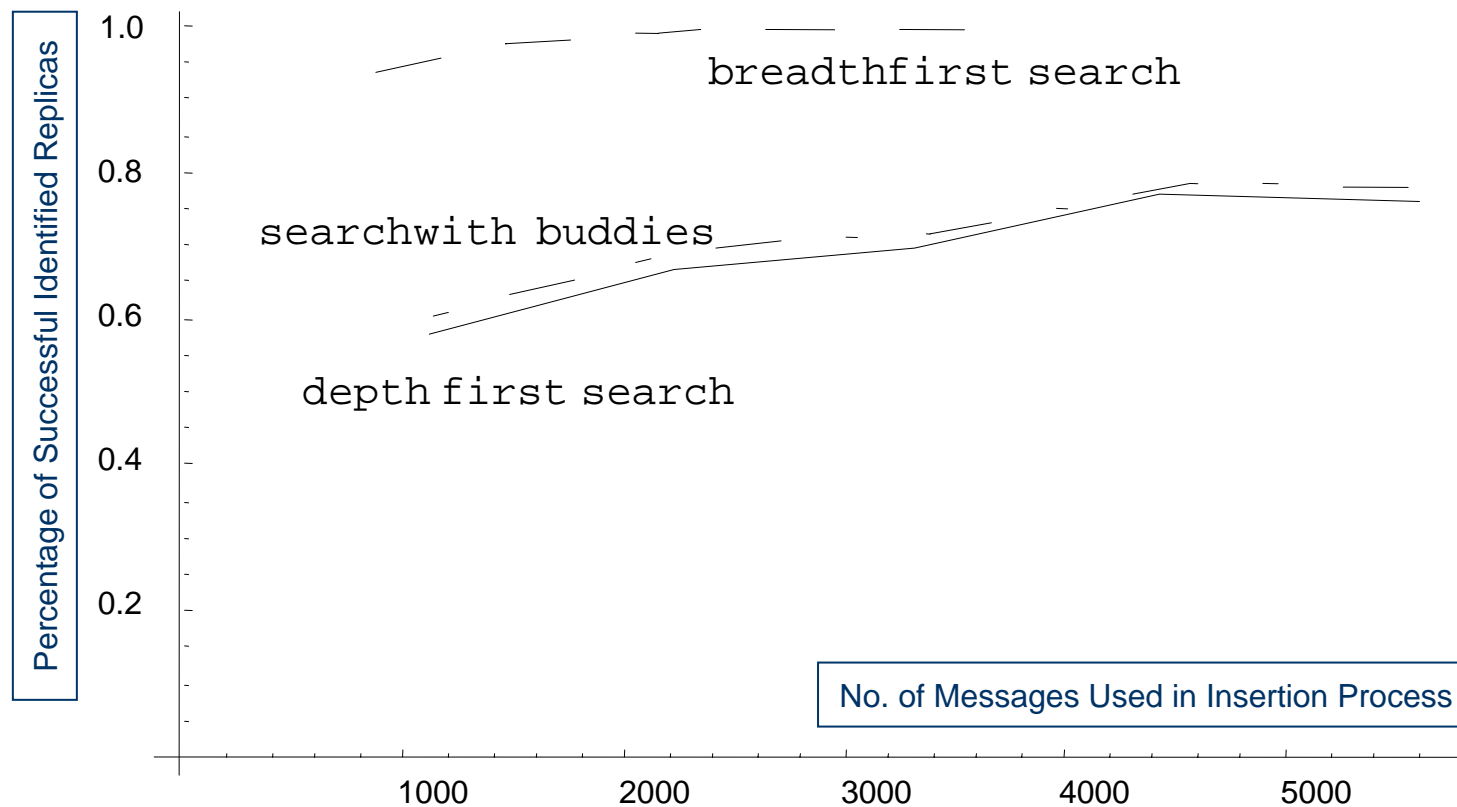
Replica distribution

- *Question 2 : The average number of replication for a peer is 19.46.*



Finding all replicas-result

Question 3 :Using breadth first searches is by far superior.



6. Conclusions

- Scalable distributed and decentralized access structures are possible
- P-Grids offer a lot of flexibility to be further exploited
- Foundation for many fully decentralized P2P applications
- Application in mobile ad-hoc networks (www.terminode.com), Swiss national research centre at EPFL

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