

Preserving Peer Replicas By Rate-Limited Sampled Voting

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Peer to Peer Seminar
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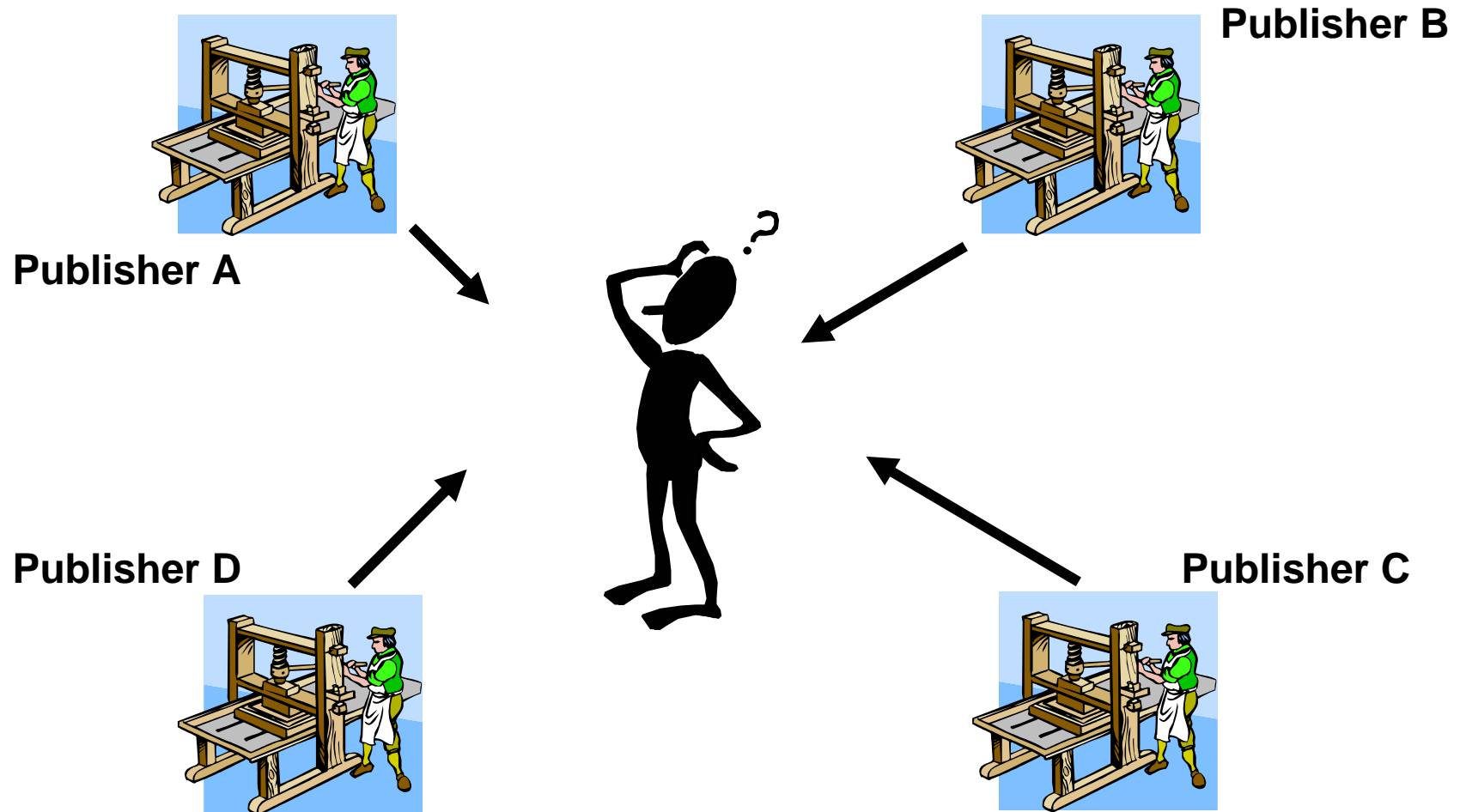


Presentation by: **Renata Dividino**

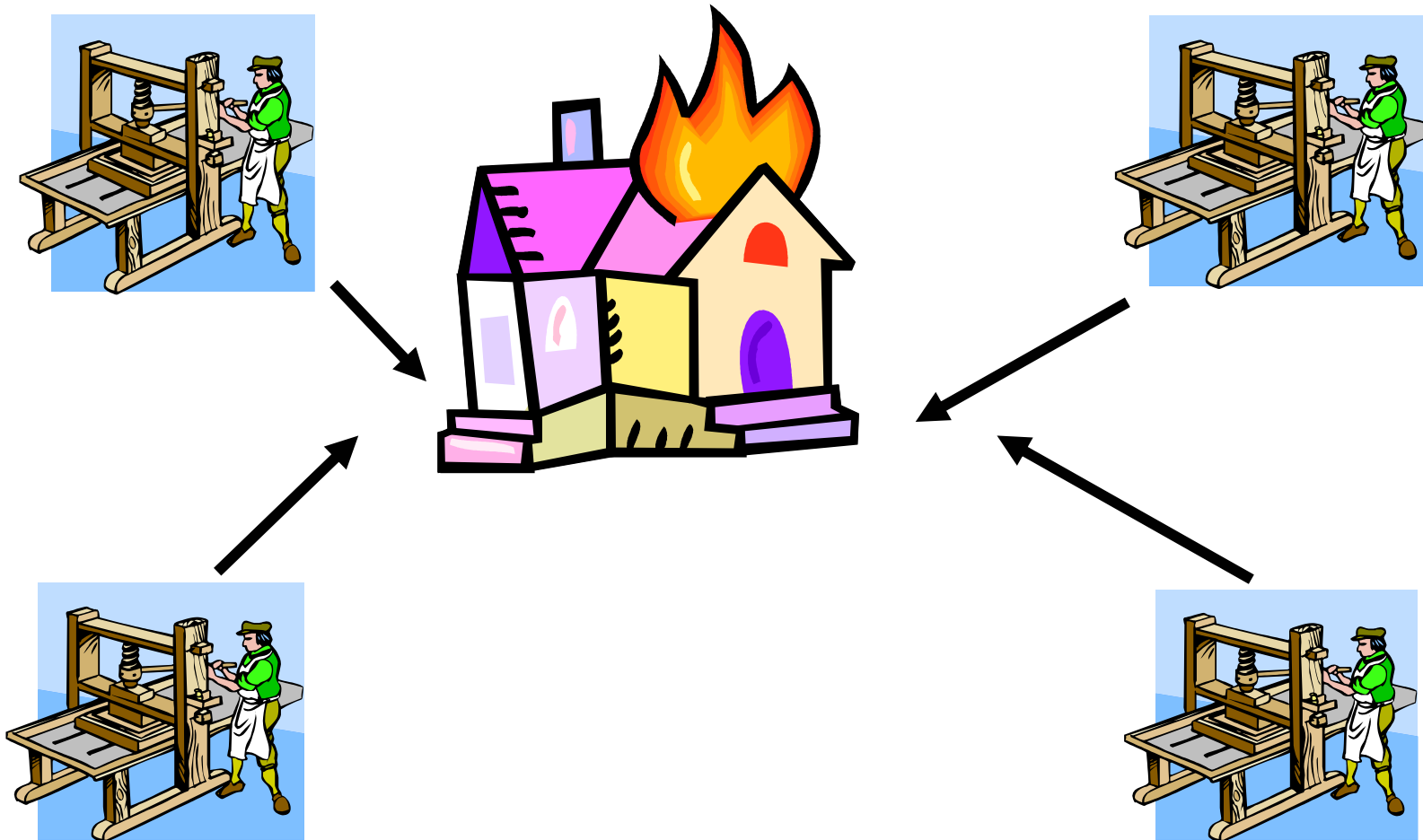
Overview

- Motivation
- Digital Preservation System
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- What's LOCKSS ?
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- Protocol Analysis
- Adversary
- Results
- Conclusion

Motivation: Why Libraries?



Motivation: Why Distributed Libraries?

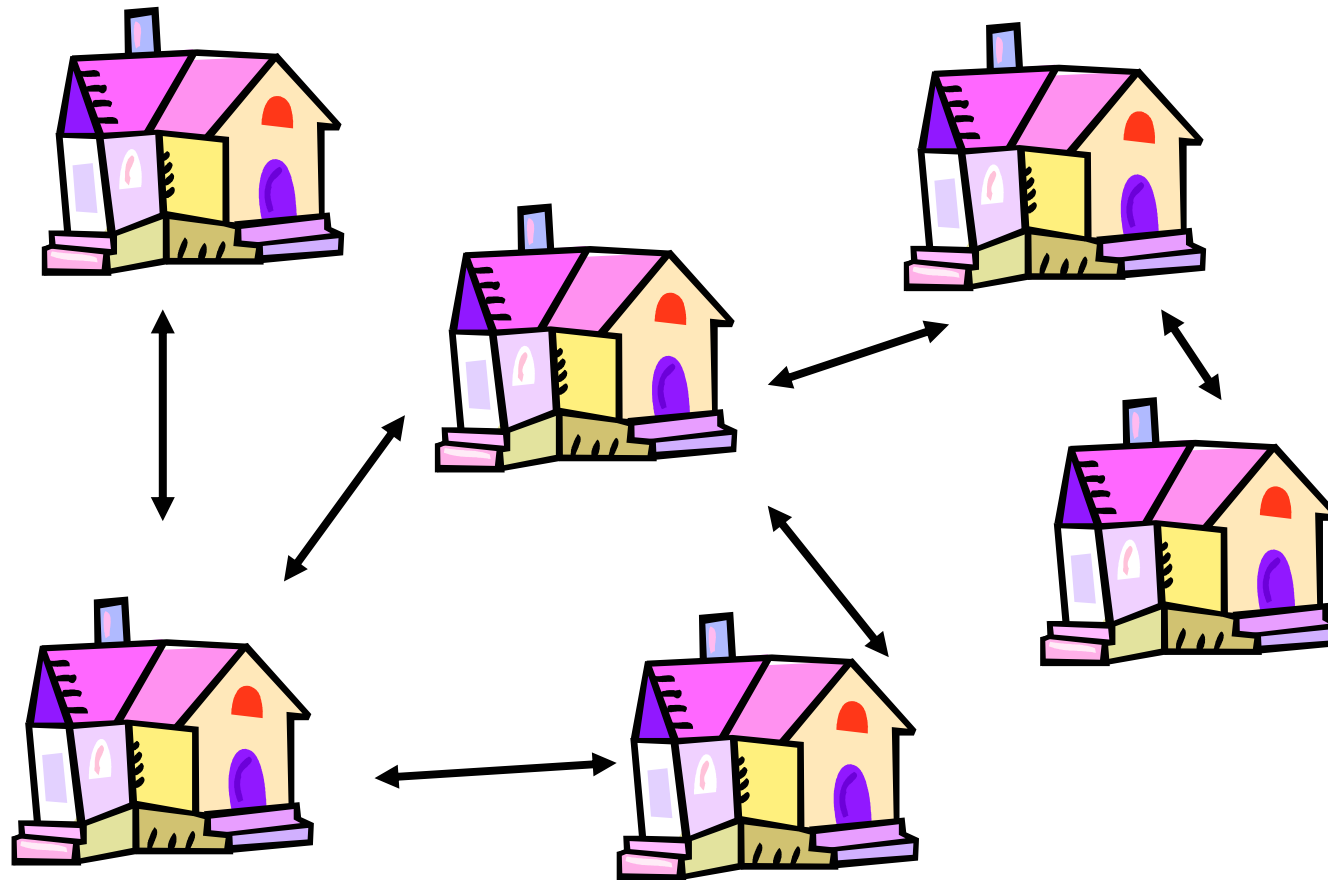


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Motivation: Distributed Digital Libraries



Digital Preservation System

- Collect material
- Distribute it to local readers
- Preserve by cooperating with others that hold the same material to detect and repair damage

Problems / Challenges

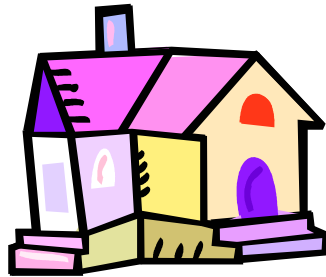
- Malign / Loyal Peers
- Sybil Effect

- Common Solution – Add extra effort involved in messages.

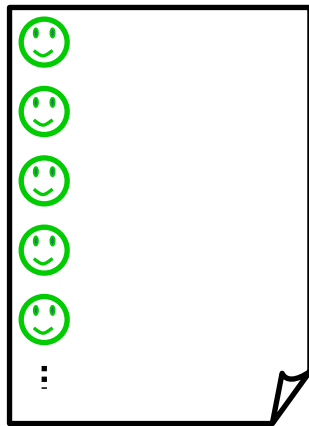
What's LOCKSS ?

- The LOCKSS (Lots Of Copies Keep Stuff Safe)
 - large number of independent, low-cost, persistent web caches,
 - protocol,
 - detect and repair damage by voting in „opinion polls“ on their documents caches.

Protocol Description (1)



Friend List



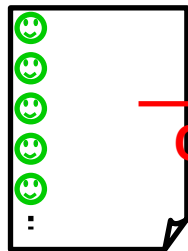
For each document (or Archival Unit - AU) a library maintains a Friend List.

Protocol Description (2)

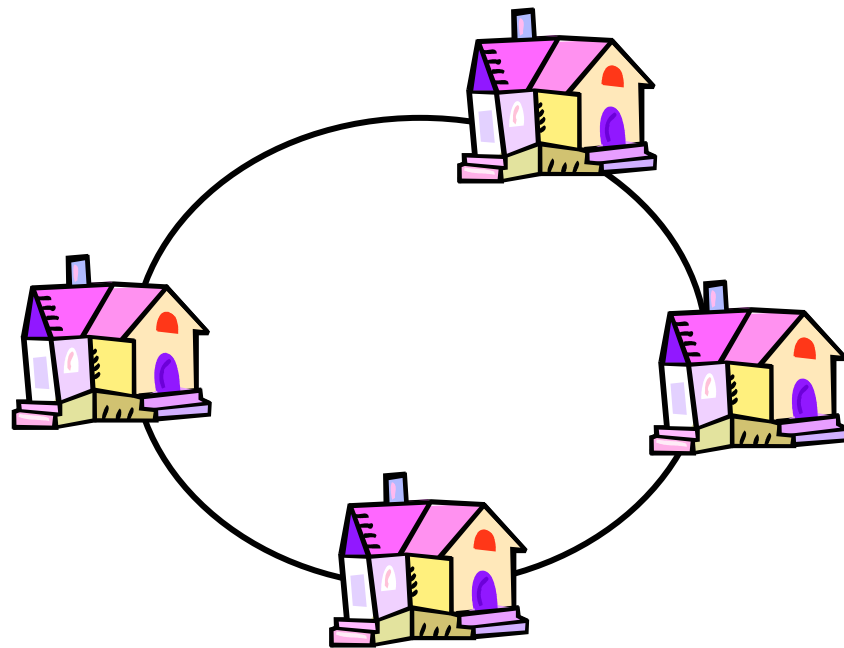
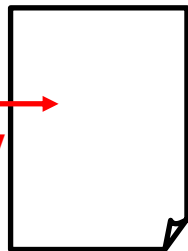


Friend

Reference

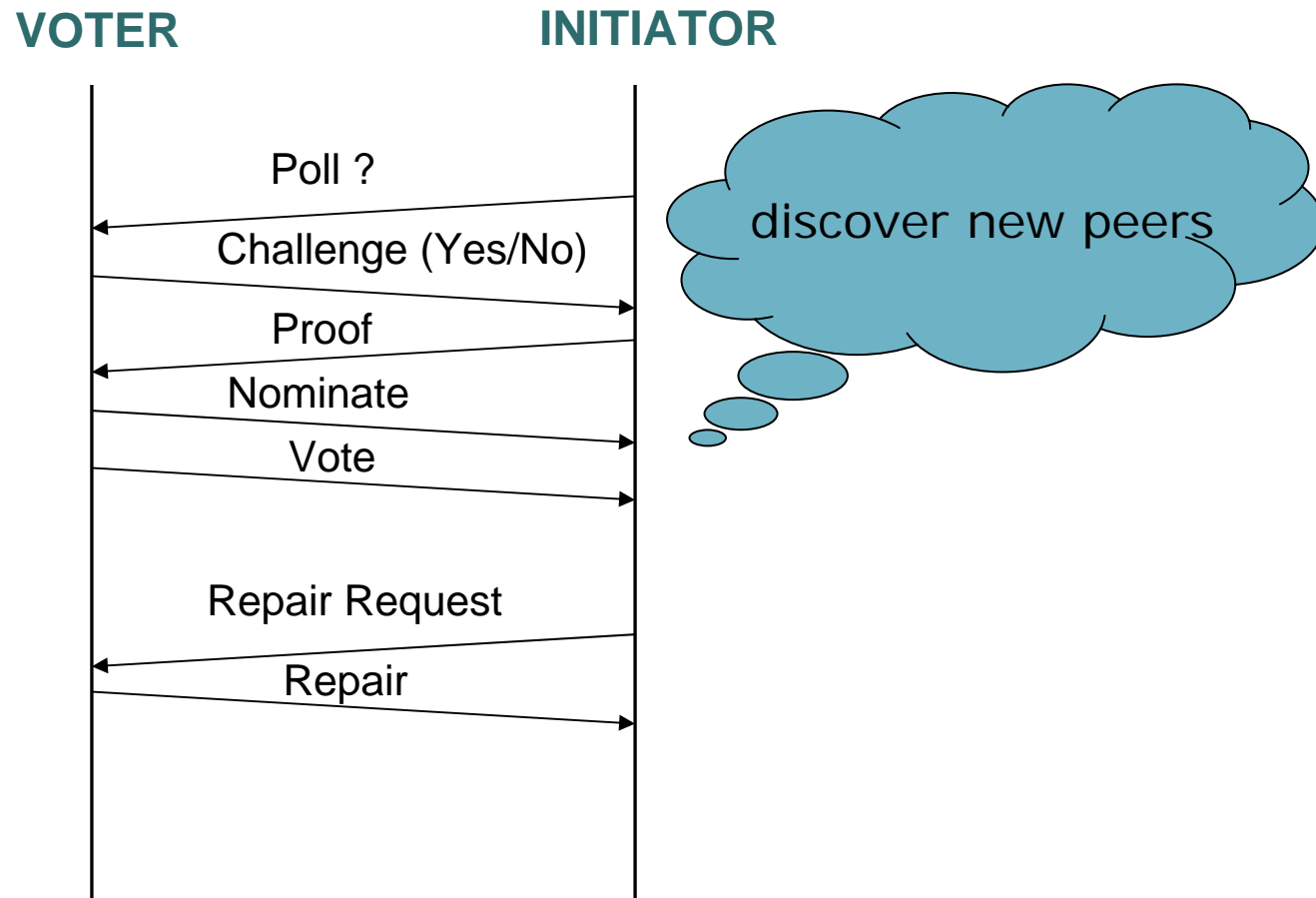


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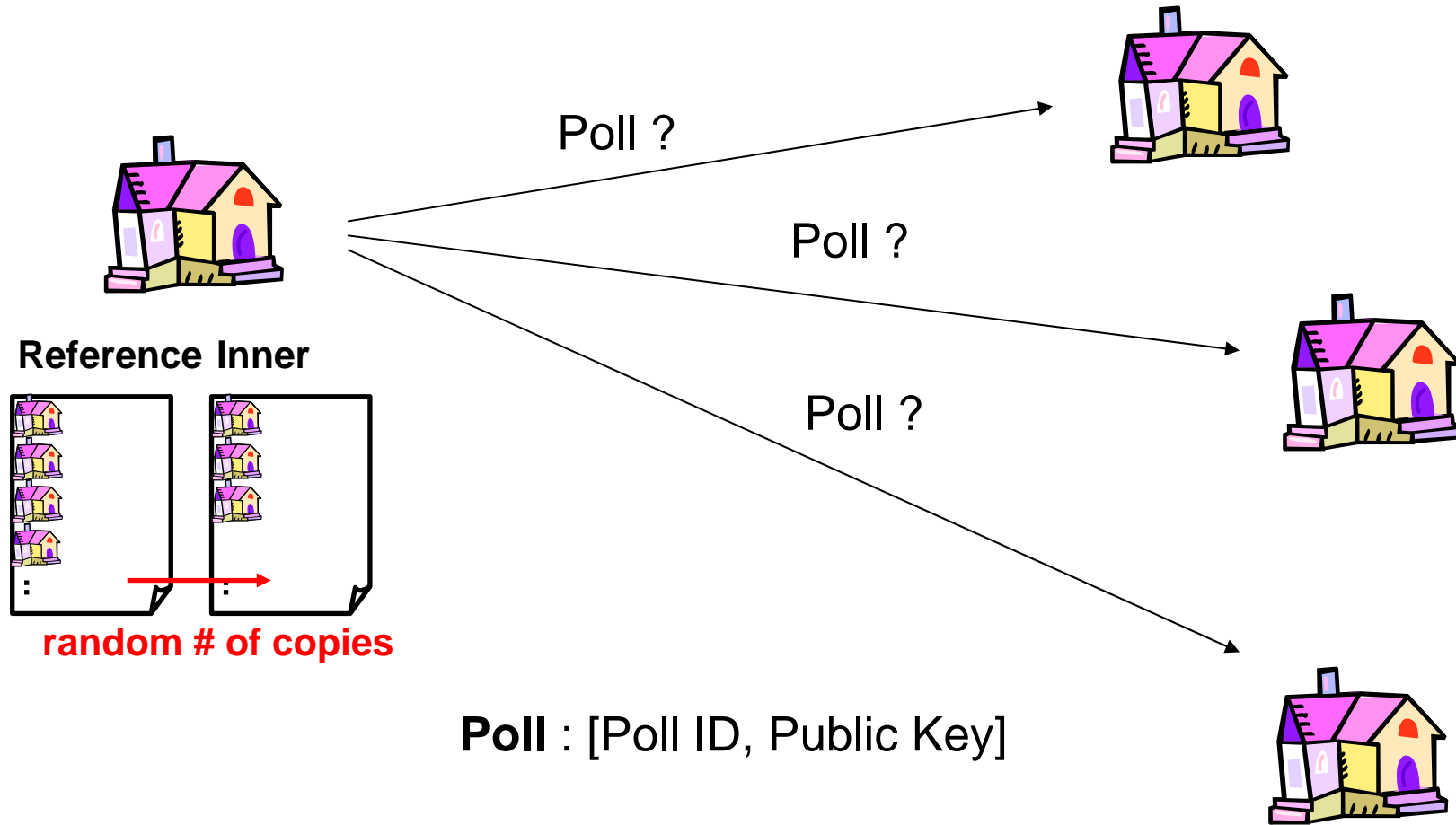


When a library joins a network,
it creates its **Reference List**

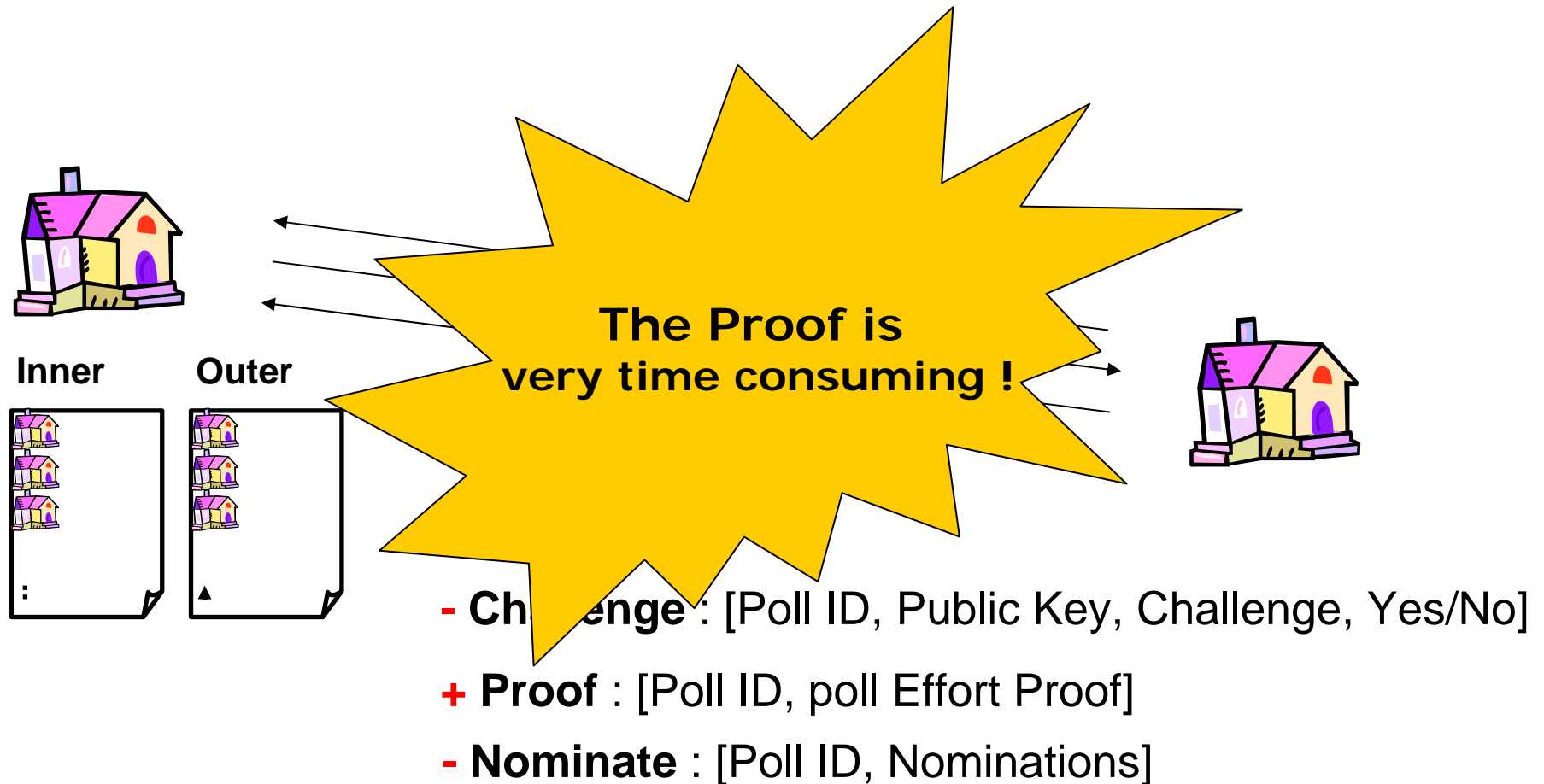
Protocol Description (3)



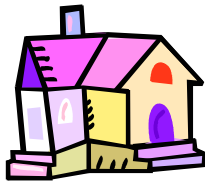
Protocol Description (4)



Protocol Description (5)



Protocol Description (6)

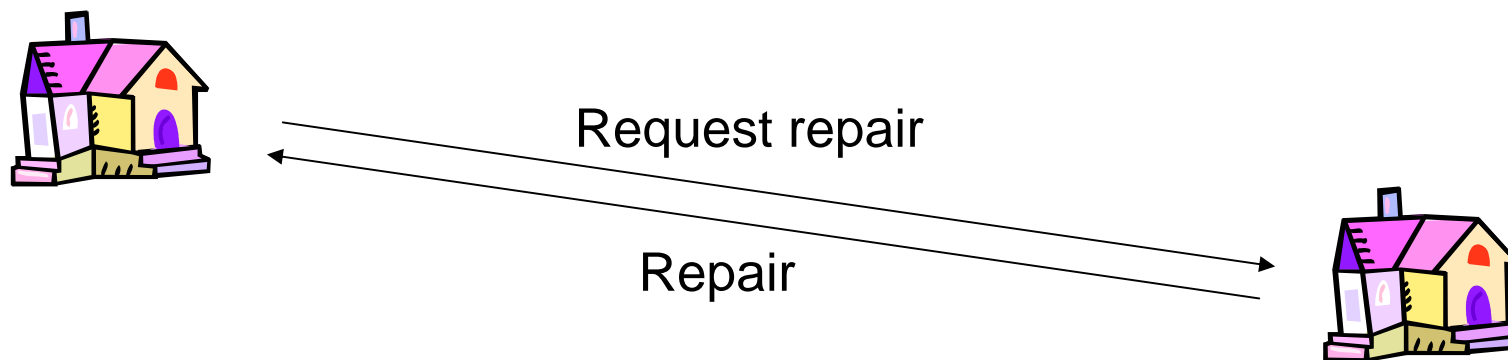


**The Vote is
very time consuming !**



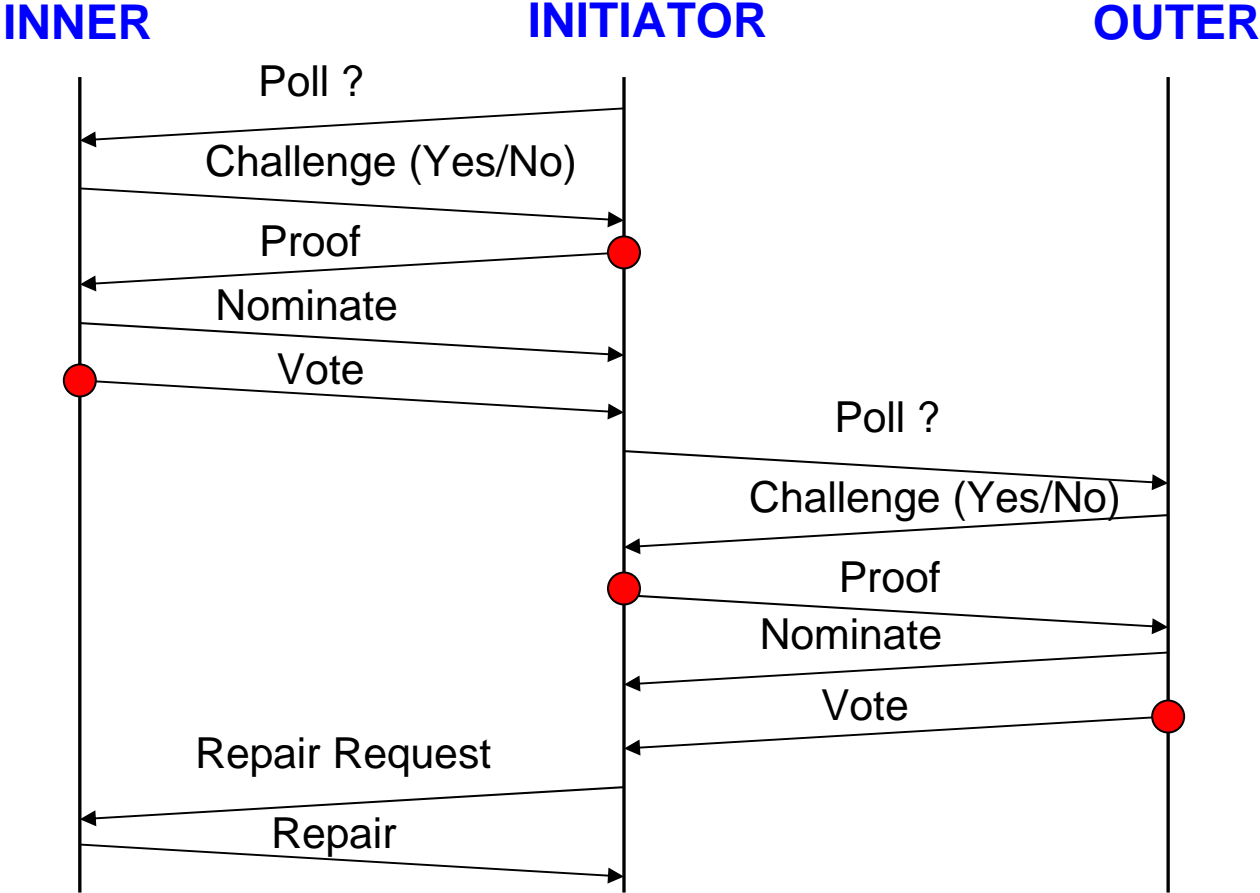
1. Invalid Vote
2. Valid Vote
 - Agreeing Vote
 - Disagreeing Vote

Protocol Description (7)



A peer supplies a repair only if the requester had previously proved with an agreeing vote that it once had the same content.

Protocol Description (8)



● **Very time consuming !**

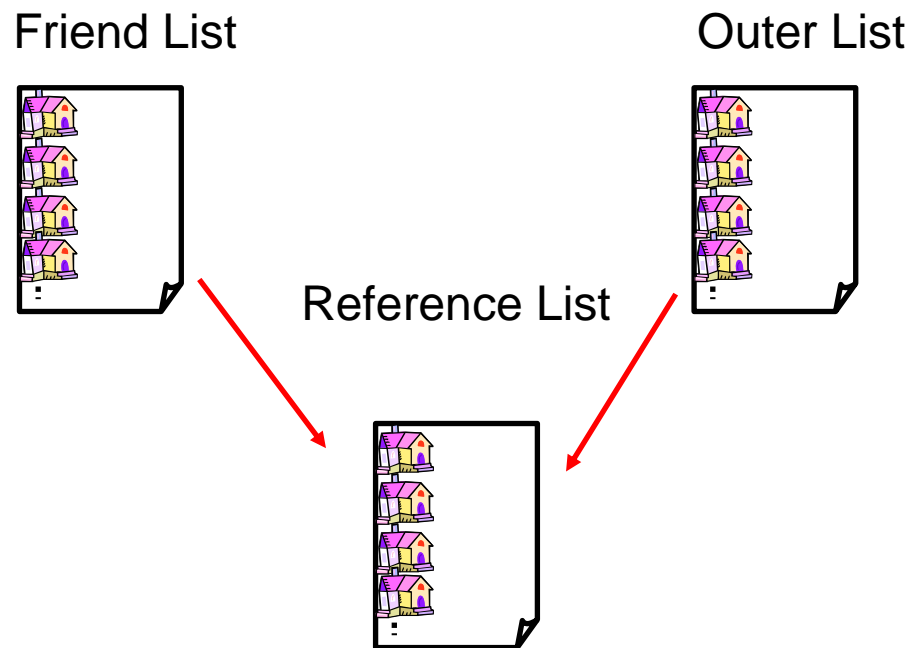
Protocol Analysis

- ❑ Reference List Churning
- ❑ Effort Sizing
- ❑ Rate Limiting
- ❑ Timeliness Of Effort
- ❑ Obfuscation of Protocol State
- ❑ Alarm

Protocol Analysis

Reference List Churning

A peer avoids depending on a fixed set of peers for maintenance of its AU.



Protocol Analysis

Effort Sizing / Rate Limiting

Large changes to a system require large effort

- Add effort involved in messages
- Rate is limited by the smaller of the adversary's effort and the effort of his victims.

Protocol Analysis

Timeliness Of Effort

Avoid that good reputation behavior is accumulated.

- Ensure that only proofs recent effort can affect the system.

Protocol Analysis

Obfuscation of Protocol State

Assume a powerful adversary capable of observing traffic at many points of the network.

- ❑ Encrypt all but the first protocol message exchanged using a fresh symmetric key.
- ❑ Make all peers invited to a poll, even those who decline a vote, go through the motions of the protocol.

Protocol Analysis

Alarm

Intrusion Detection

The protocol raises an alarm when a peer:

- ❑ determines that a poll is inconclusive,
- ❑ suspects local spoofing,
- ❑ has been unable to complete a poll for a long time.

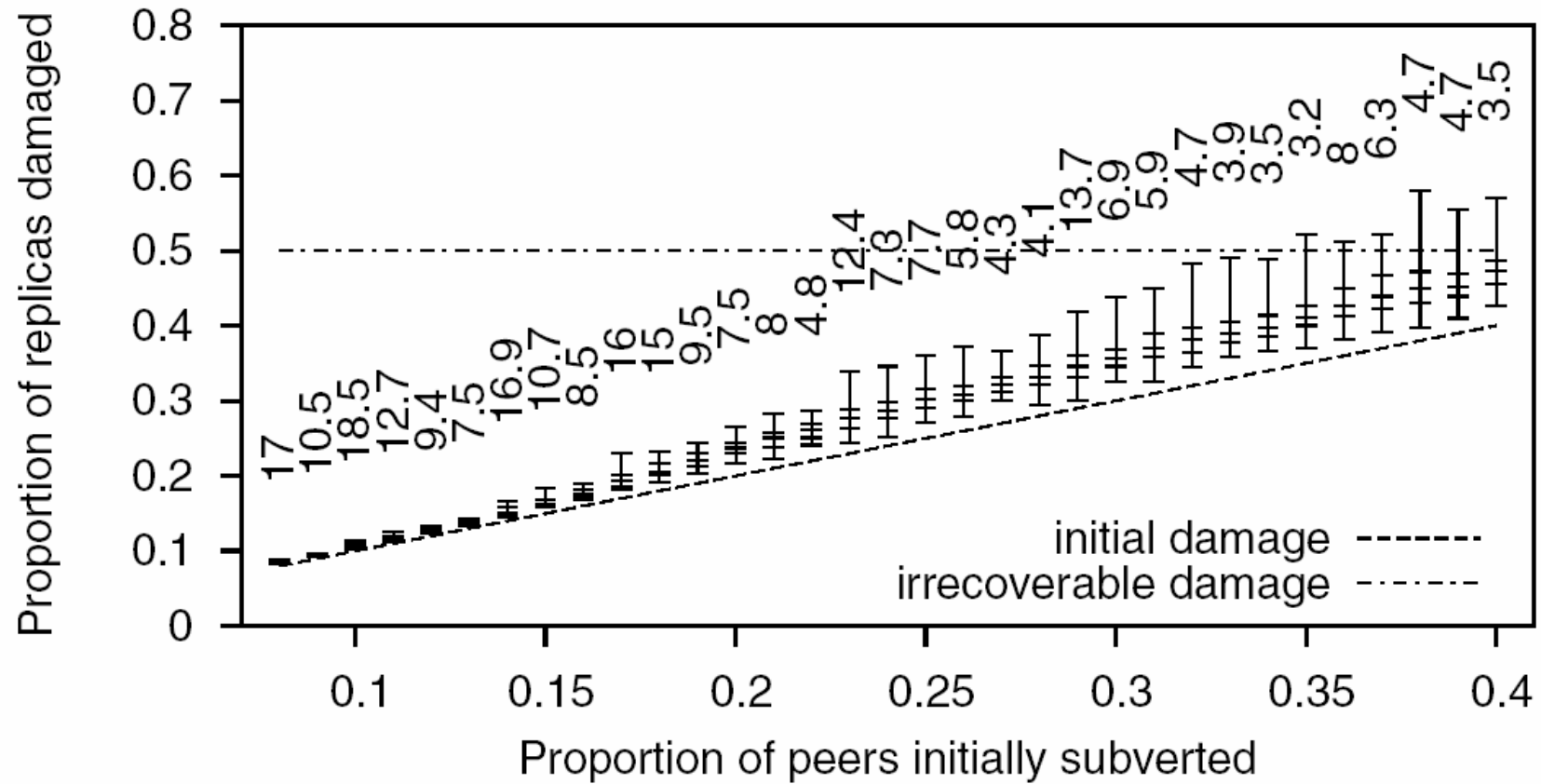
Adversary

□ Stealth Modification

- The adversary's goal is to change as many replicas of the content held by loyal peers as possible without being detected.
- Techniques:
 - Foothold in a Reference List
 - Delayed Commitment

LOCKSS: Reference List Churning, Rate-Limiting and Check consistence between repair and vote.

Results



Conclusion

- LOCKSS can produce a peer-to-peer system with remarkable ability to resist attacks over decades.
 - Massive Replicas
 - Rate Limitation
 - Intrusion Detection (->Alarm)
 - Costly Operation



□ Thank you for your attention!

□ Questions?