GALANX: An Efficient Peer-to-Peer Search Engine System

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Overview

- Introduction: What is GALANX
- The Key Idea
- State of the Art
- A first DHT-Solution
- GALANX in Detail
- Implementation & Tests
- Conclusions & Future Work

Introduction: What is GALANX

- Looking at large data networks (e.g. "the Web", corporate LANs or perhaps file-sharing networks) we have a lot of data providers and even more clients searching for specific data
- How to route search queries efficient?
 - Google: centralised, fast lookups, not scalable, expensive, vulnerable, not up to date, not covering the whole web

Introduction: What is GALANX

- Scanning the whole web instantly with a centralized approach (Google-like) seems to be impossible looking at the cost
- Peer-to-Peer looks like a good option, but several open questions how to do this right
- GALANX focuses on the core of how to route a query fast and efficient to the appropriate data provider



- Each data provider is a node in the Peer-to-Peer network
- Each node publishes an Index which holds information about the data it shares

The Key Idea - Properties

- The system is adaptable to complicated queries (e.g. full text search)
- Ownership of data is respected (no data caching or moving between nodes)
- Every node has as much freedom as possible

The Key Idea – Simplify

- We assume:
 - Data objects are simple text documents
 - Requests are lists of keywords, we have a hit if all keywords appear in a file
 - Data sharing nodes (providers) are stable

Keyword ₁	Doc _{1, 1} , Doc _{1, 2} ,, Doc _{1, k1}
Keyword ₂	Doc _{2, 1} , Doc _{2, 2} ,, Doc _{2, k2}
	••••
Keyword _n	Doc _{n, 1} , Doc _{n, 2} ,, Doc _{n, kn}

Fig.1 - The Local Data Index on an arbitrary node

State of the Art in P2P routing

- Centralized Index (Napster)
 - Not scalable, vulnerable, "no real P2P"
- Flooding (Gnutella)
 - wastes bandwidth, not reliable (TTL)
- DHT Indices
 - Scalable and reliable, distributed load
 - BUT distributes data deterministic over the network by a special key-value which would imply moving of data
 - A good starting point ?

A first DHT-Solution - Idea

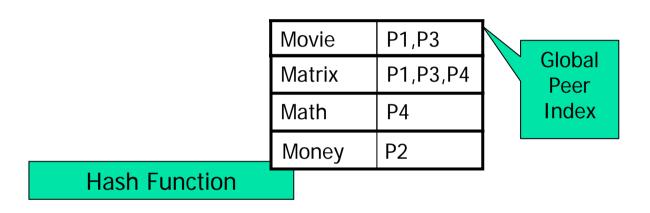
- Idea: Put a "DHT routing layer" above the real data network
- For that we need **Peer Indices**:

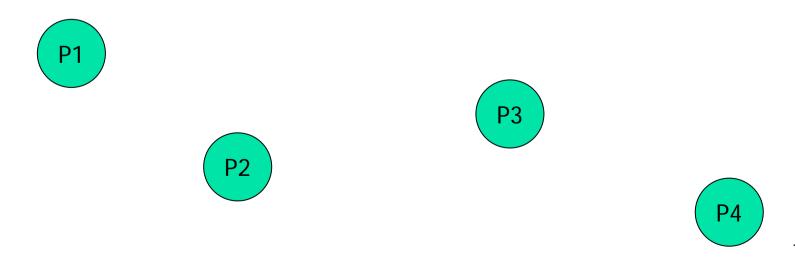
Keyword ₁	Peer _{1, 1} , Peer _{1, 2} ,, Peer _{1, k1}
Keyword ₂	Peer _{2, 1} , Peer _{2, 2} ,, Peer _{2, k2}
Keyword _n	Peer _{n, 1} , Peer _{n, 2} ,, Peer _{n, kn}

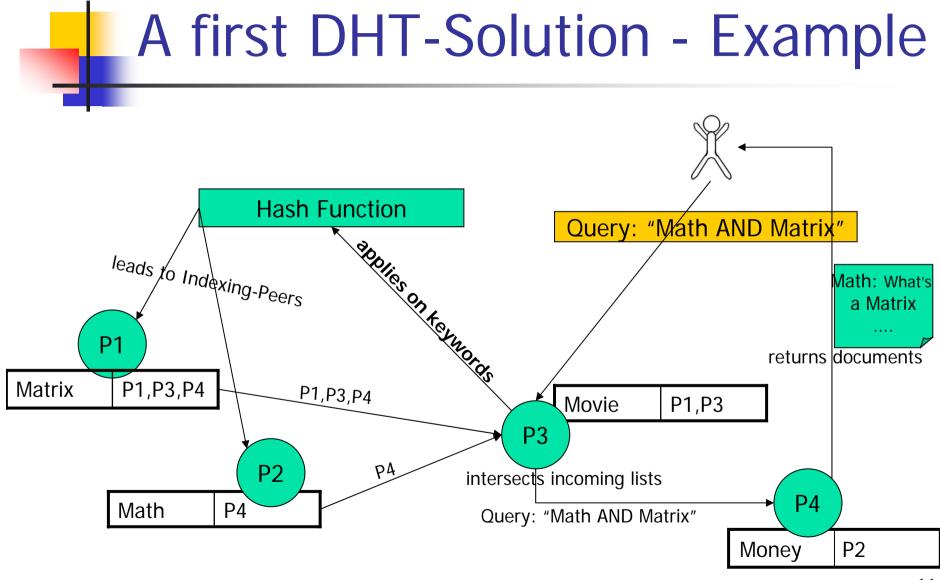
Fig.2 - A Peer Index

If such a Peer Index covers every Keyword in the network, we call it a <u>Global Peer Index</u>.

A first DHT-Solution - Example







DHT-Solution – Good Enough?

- It is reliable, theoretically scalable, Ownership respecting and adaptable
- BUT nodes are not free to decide themselves what indices to maintain
- Performance Loss through "Extra"-Lookup
- Deterministic Distribution increases Overhead

GALANX in Detail – A perfect Solution?

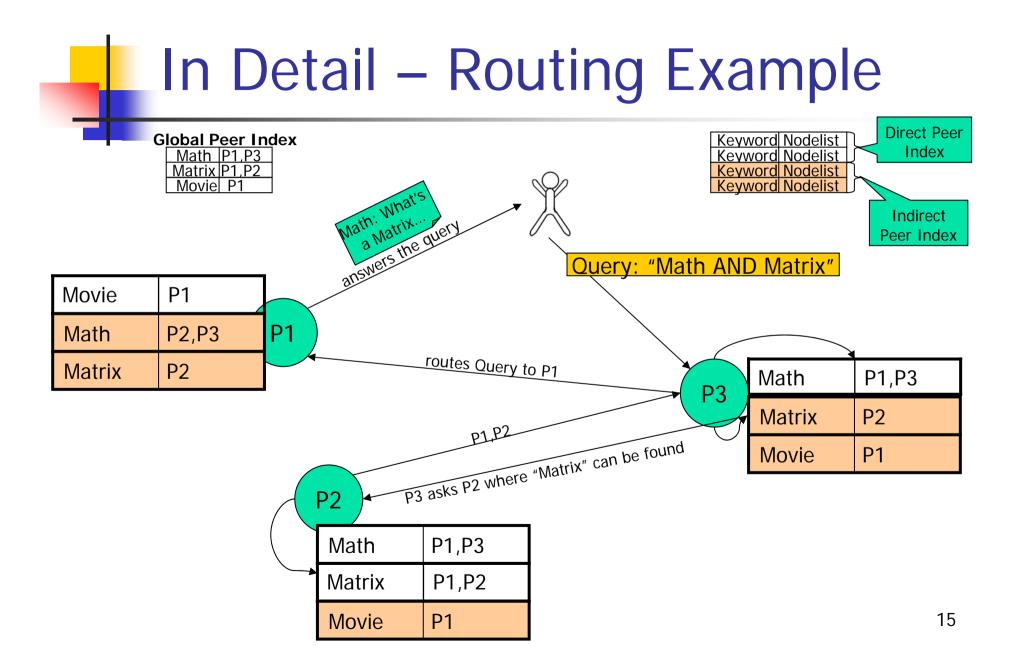
- Imagine a configuration where every node maintains the complete Global Peer Index
- In terms of query routing that would be a perfect situation, because every query could be routed directly to the correct node(s)
- BUT Updates would be extremely expensive and Peer Indices would be very large

In Detail – GALANX Peer Indices

Keyword ₁	Peer _{1, 1} , Peer _{1, 2} ,, Peer _{1, k1}	
Keyword ₂	Peer _{2, 1} , Peer _{2, 2} ,, Peer _{2, k2}	Direct Peer
	•••	Index
Keyword _j	Peer _{j, 1} , Peer _{j, 2} ,, Peer _{j, kj}	
Keyword _{j+1}	$Peer_{j+1, 1}, Peer_{j+1, 2}$	
Keyword _{j+2}	Peer _{j+2, 1} , Peer _{j+2, 2} , Peer _{j+2, 3}	Indirect Peer In-
		dex
Keyword _n	Peer _{n, 1} , Peer _{n, 2}	

Fig.3 – GALANX Peer Index

- The **Direct Peer Index** is a Peer Index as known
 - covering several keywords with full list of data holding peers
- The <u>Indirect Peer Index</u> has for each of its keywords a list of peers maintaining a Direct Peer Index for this keyword
- Both together cover all keywords in the network



In Detail – Properties?

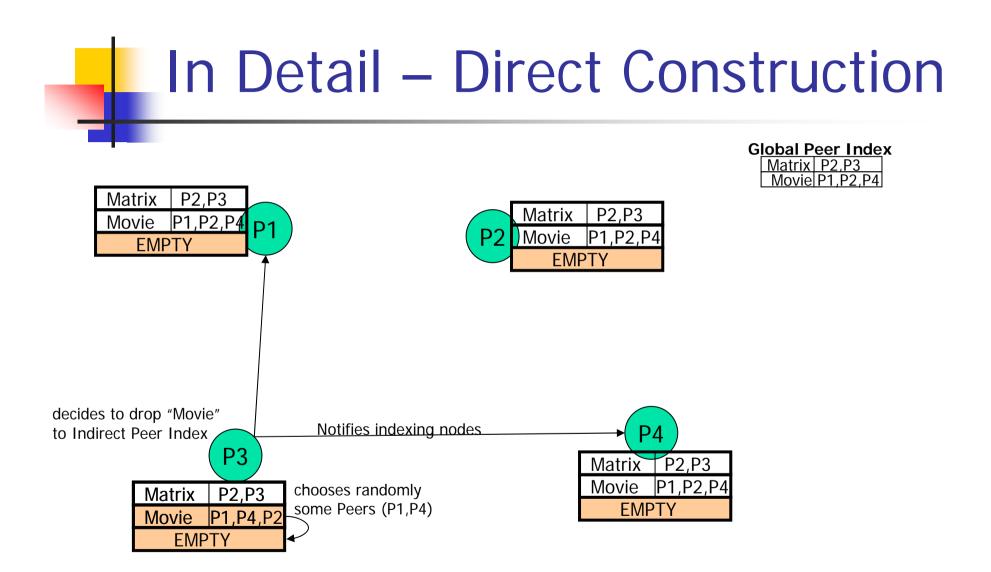
- If all keywords of a query are in the Direct Peer Index, that's perfect
- If not, the remaining keywords can probably be found at just one other node or at least at less than with pure DHT distribution (by intersecting the Peer lists wisely)
- If a node's strategy is good, most popular words can be found in Direct Peer Index

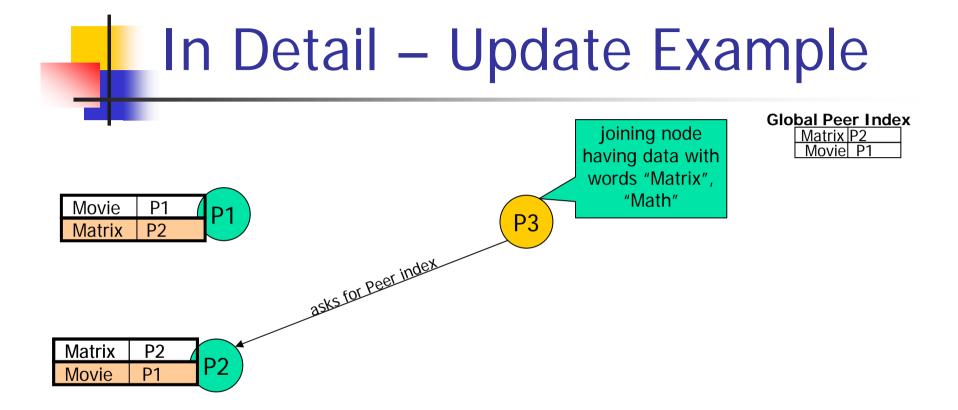
In Detail – Direct Index Construction

- Starting with few (or one) nodes
 - every node holds the complete Global Peer Index
- Joining nodes copy an arbitrary Peer Index
 - update that with own data
- Nodes can drop keywords from Direct to Indirect Peer Index by:
 - Randomly pick a few "keyword-data" holding peers
 - Put them in the Indirect Peer Index as indexing nodes
 - Notify them, so they can notify you if they drop that word

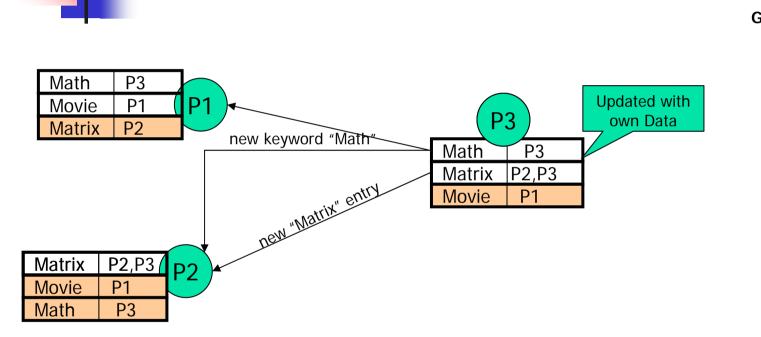
In Detail – Keyword Updates

- Keyword already exists in the network:
 - Entry in the Peer Index -> Update sent to all indexing nodes
- Keyword is completely new:
 - Update is sent to all nodes in the network
- Possible: lazy Update-strategy with wildcards



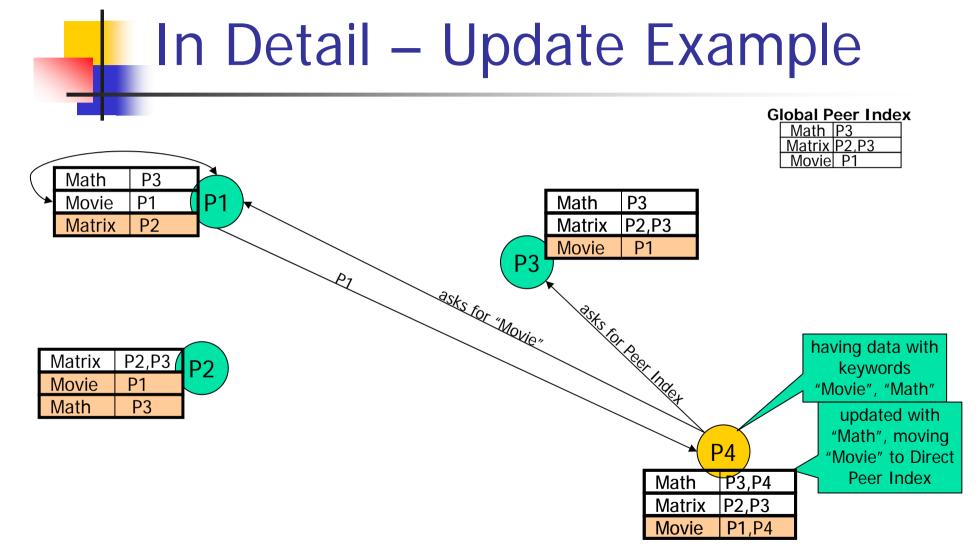


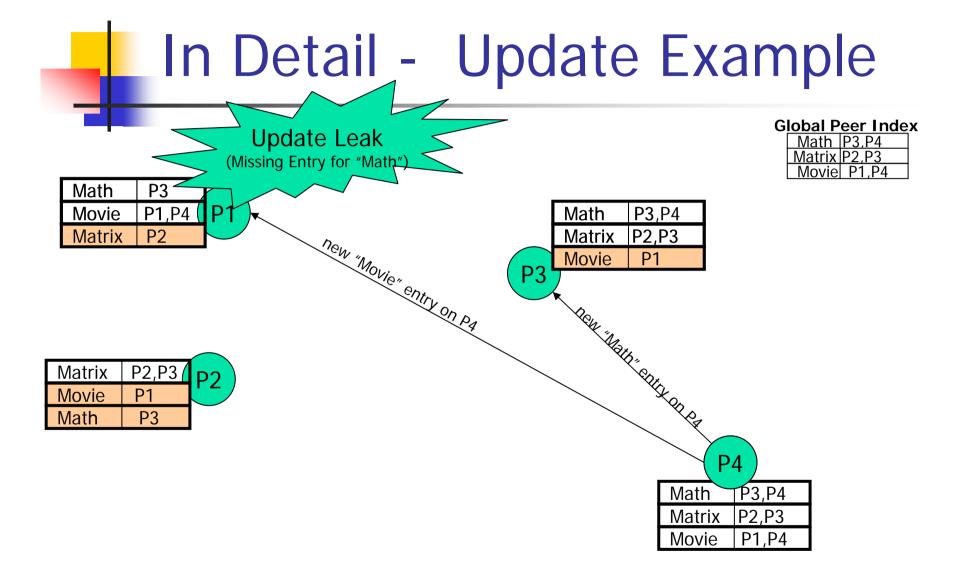
In Detail – Update Example



Global Peer Index

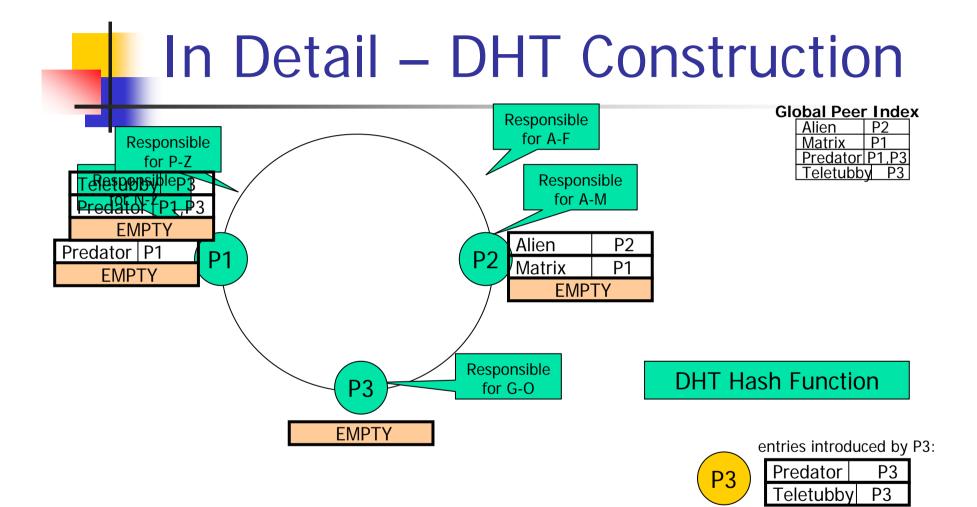
Math P3 Matrix P2,P3 Movie P1



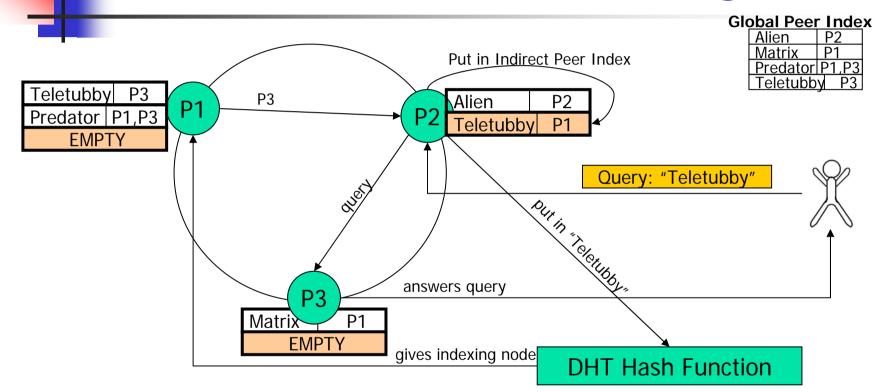


In Detail – DHT Index Construction

- Starting again with few nodes
 - The Global Peer Index is distributed over all peers via a Hash function
 - Every Peer is responsible for indexing a part of the Keyspace
- Joining nodes get "their" Keyspace by the DHT mechanism, corresponding Peer Index entries are moved
- Updates work similar as seen using DHT mechanism
- Nodes trained by Queries
 - Keywords taken in Indirect (or Direct) Peer Index if not already in Direct Peer Index)



In Detail – DHT Training



Implementation

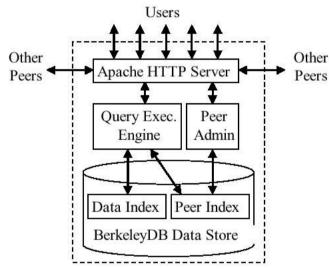


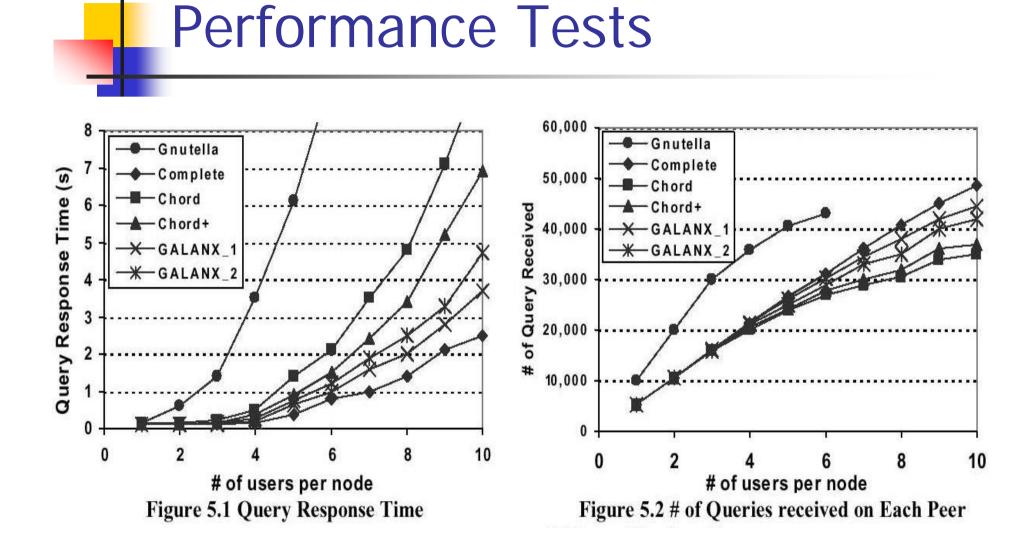
Fig.4 - GALANX Prototype Architecture

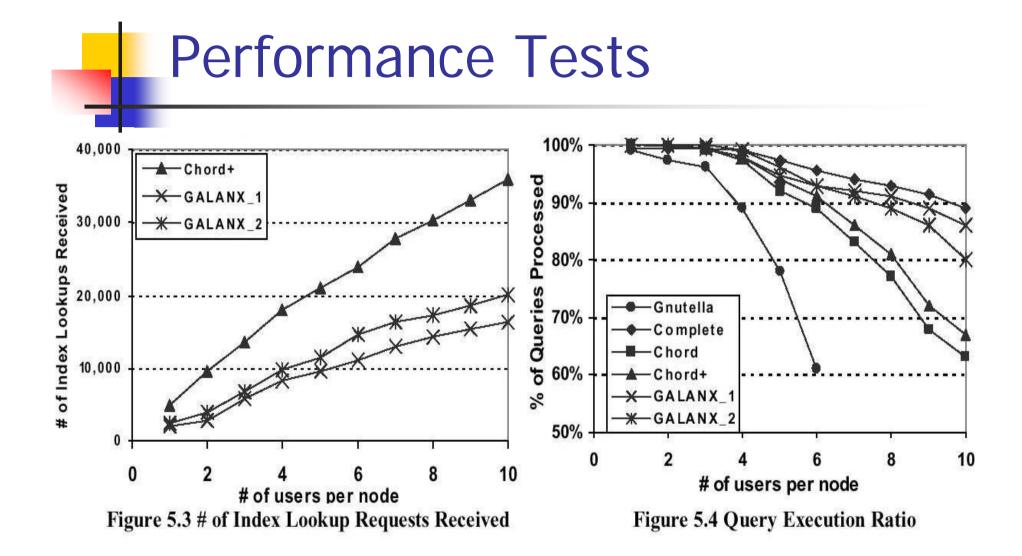
- Apache Server handles all communication
- Query Execution Engine, Peer Admin, User Interface embedded as modules, written in C
- BerkeleyDB as Database holding Data & Peer Index

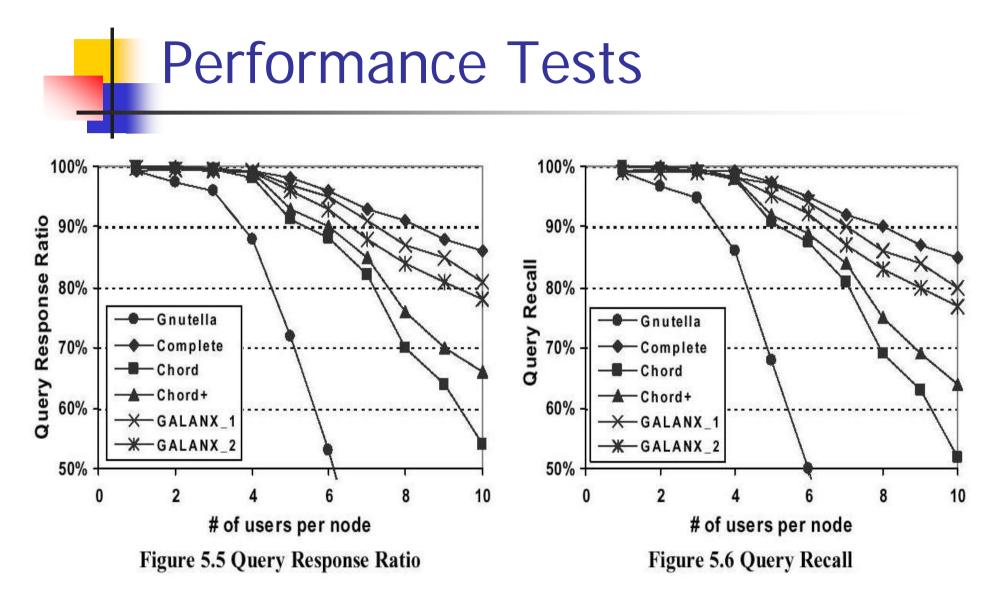
Apache Version multi-process mode doesn't work well with BerkeleyDB => **lower Performance as possible** because queries are queued in some cases

Performance Tests

- Experimental Setup:
 - 100-node computer LAN-cluster, 933/550MHz, 1GB memory
 - 20million + HTML-pages distributed by category to nodes
 - 2.49 million keywords, average 125.000 keywords per node
 - Simulated user queries with dynamic expiration time
 - Initial expiration time 10 sec, min 5 sec, max 30 sec
 - Only queries which can be answered







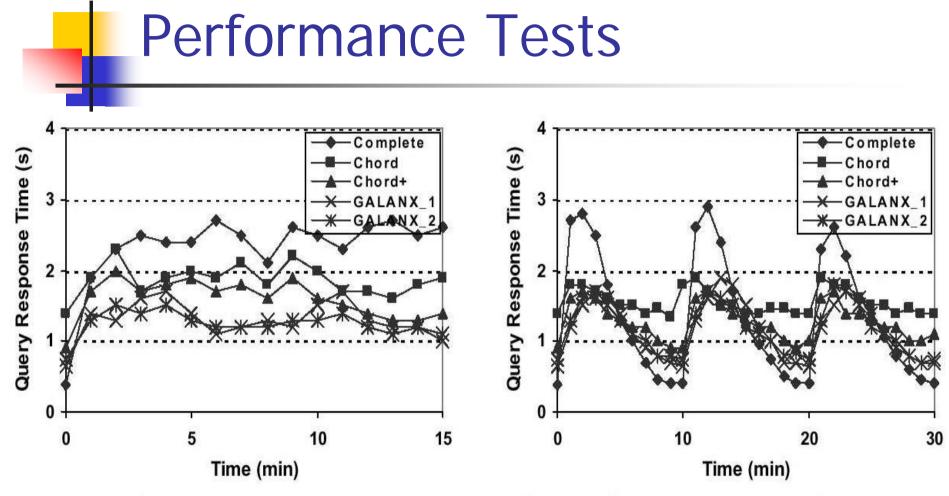


Figure 5.7 Query Response Time (with new data)

Figure 5.8 Query Response Time (with new nodes)

Conclusions & Future Work

- Fast routing Approach, IF Assumptions hold
- Just one piece of the whole; to get a really applicable system lot has to be added
- What about reliability in really dynamic Environment?
- Authors next aim:
 - Extend to more complex forms of queries
 - Introduce Information Retrieval Techniques like relevance ranking

Summary

- Want to have fast query routing system
 - Existing approaches expensive, not fast or not reliable->GALANX
- DHT-Approach
 - Global Peer Index distributed over all nodes
 - Not good enough (extra lookup, nodes not free)
- GALANX Peer Indices
 - Direct & Indirect Peer Index, Design delivers peers freedom
 - Direct Construction: complete global peer index on every node, updates sent to indexing nodes or complete network, dropping keywords to indirect Peer Index
 - DHT Construction: Having DHT-Distribution of Keyspace, nodes responsible (direct indexing) for certain Keyspace, training with queries
 - Experiments look good, although system not perfect

Thanks

for your attendanceQuestions ?...