applications

MULTENANT-DATABASES

on demand

software as a service

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SaaS

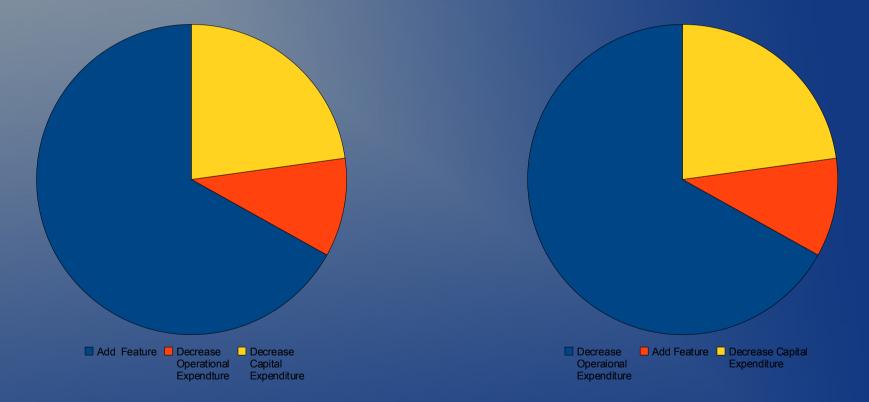
Outline

- Motivation
 - Software as a Service
 - Need for Multi-tenancy
 - Need for Multi-tenant Databases
- Multi-tenant Databases
 - Challenges
 - Design Approaches and tradeoff
 - Experiment
- Discussion
- Conclusion

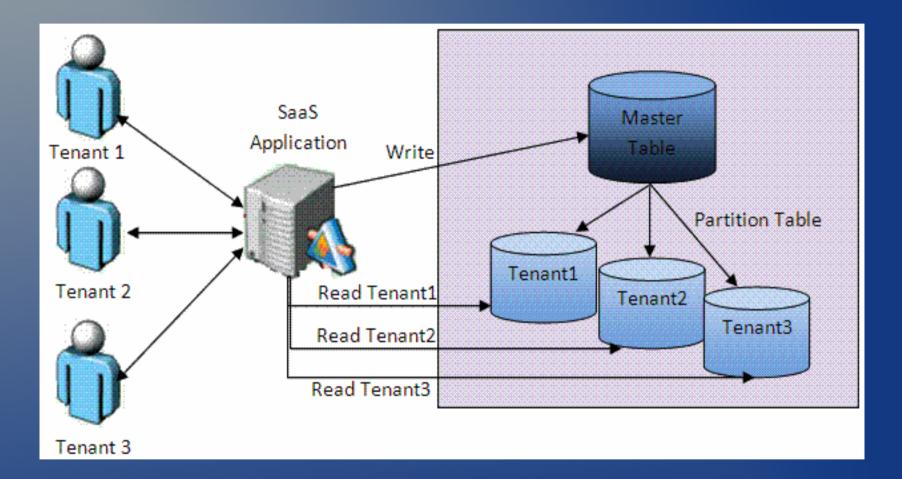
Today's Software Service Industry

• On Premise Software

• Software as a Service



Multi-tenancy



Why Multi-tenant databases ?

- Consolidating multiple databases onto same operational system
- Reduces Total Cost of Ownership

Challenges in Multi-tenant Database

- Scalability
 - Tradeoff between cost handling many tables and cost query rewriting
- Allow Schema Extensibility
 - Multiple tenant share tables
 - Need for tenant specific schema extensibility

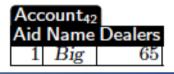
Design Approaches

Private Table

- Natural Thing to do each tenant gets a private schema
- Low cost on query transformation
- Less consolidation

Account ₁₇			A
Aid Name	Hospital	Beds	Α
1 Acme	St. Mary	135	
2 Gump	State	1042	





Extension Table

- Split off extensions into separate tables
- Higher cost on Query transformation
- Slightly better consolidation

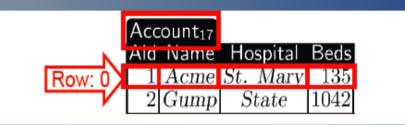
	_				Healthca	area	ccount	
Account					Tenant F	Row	Hospital	Beds
Tenant R	ow /	Aid	Name		17	0	St. Mary	135
17	0	1	Acme		17	1	State	1042
17	1	2	Gump	Ľ				
35	0	1	Ball	1	Automo	tive,	Account	
42	0	1	Big		Tenant F			
					42	0	65	

Universal Table

- Generic Structure with VARCHAR value columns
 - n-th Column of the logical table is mapped to Col-n in the universal table
 - Extensibility
- Disadvantages
 - Many Null Values
 - Not type safe
 - No Indexing

Universa			Col2	Col3	Col4	Col5	Col6
17		1		St. Mary		COIS	
17	0	2	Gump	-	1042	_	_
35	1	1	Ball	State	1042	_	_
49	1	1		65	_	_	_
42	2	1	ыg	60	_	_	—

Pivot Table



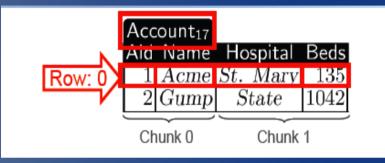
Pivot _{st}				
Tenant	: Table	Col	Row	Str
17	0	1	0	Acme
17	0	2	0	St. Mary
17	0	1	1	Gump
17	0	2	1	State
35	1	1	0	Ball
42	2	1	0	Big

Pivot _{int} Tenant	Table	Col	Row	Int
17	0	0	0	1
17	0	3	0	135
17	0	0	1	2
17	0	3	1	1042
35	1	0	0	1
42	2	0	0	1
42	2	2	0	65

Chunk Table

Generic Structure

- Suitable if data-set can be partitioned into dense subsets
- Derived from Pivot table
- Performance
 - Fewer joins for reconstruction if densely populated subsets can be extracted
 - Reduced metadata/data ratio dependent on the chunk size
 - Indexable



Chunk _{in}					
Tenant	Table Ch	nunk R	low	Int1	Str1
17	0	0	0	1	Acme
17	0	1	0	135	St. Mary
17	0	0	1	2	Gump
17	0	1	1	1042	State
35	1	0	0	1	Ball
42	2	0	0	1	Big
42	2	1	0	65	-

Row Fragmentation

- Combine different schema mappings for getting best fit
 - Mixes Extension and Chunk Tables
 - Each fragment can be stored in an optimal schema layout
- Optimal row fragmentation depends on
 - Workload
 - Data distribution
 - Data popularity

Account Tenant		Aid	Name
17	0	1	Acme
17	1	2	Gump
35	0	1	Ball
42	0	1	Big

Chunk _R					
Tenant	Table	Chunk	Row	Int1	Str1
17	0	0	0	135	St. Mary
17	0	0	1	1042	State
42	2	0	0	65	—

Query Transformation

- Reconstructing original tav requires many equi-joins
- Source Query
 - SELECT Beds
 - FROM Account17
 - WHERE Hospital = 'State'
- Collect table and column names
 - Account17 : Beds , Hospital
- Obtain chunk tables and meta-data
 - Chunk (int|str)
 - Account17 :
 - Table = 0, Tenant =17
 - Beds , Hospital :
 - Chunk =1

	count ₁₇		
Aid	Name	Hospital	Beds
1	Acme	St. Mary	135
2	Gump	State	1042

Chunk _{ir}	nt str				
Tenant	Table	Chunk	Row	Int1	Str1
17	0	0	0	1	Acme
17	0	1	0	135	St. Mary
17	0	0	1	2	Gump
17	0	1	1	1042	State
35	1	0	0	1	Ball
42	2	0	0	1	Big
42	2	1	0	65	_

Query Transformation

• Generate filter query

SELECT Str1 as Hospital , Int1 as Beds FROM Chunk (int(str)

WHERE Tenant = 17 AND Table = 0 AND Chunk = 1

Replace reference in source query

SELECT Beds FROM (SELECT Str1 as Hospital Int1 as Beds FROM Chunk (int|str) WHERE Tenant = 17 AND Table =0 AND Chunk =1) As Account17 WHERE Hospital = 'State'

Query Transformation

Structural Changes

- Additional Nesting
- Joins
- Base Table Access
- Impact on Performance
 - Nesting can be flattened by query optimizer
 - Joins are cheaper only if the cost of loading the chunks and applying index supported join are cheaper that loading wider conventional relation
 - Meta data columns in base tables have indexing support

Query Evaluation Experiment

• Goal

- Show if the query transformation can handle issues of scalibility
- Evaluate impact of Join overhead
- Evaluate impact of meta-data overhead
- Test Query

SELECT p.id, ...

FROM parent p , child c WHERE p.id = c.parent AND p.id = ?.

Query Evaluation Experiments

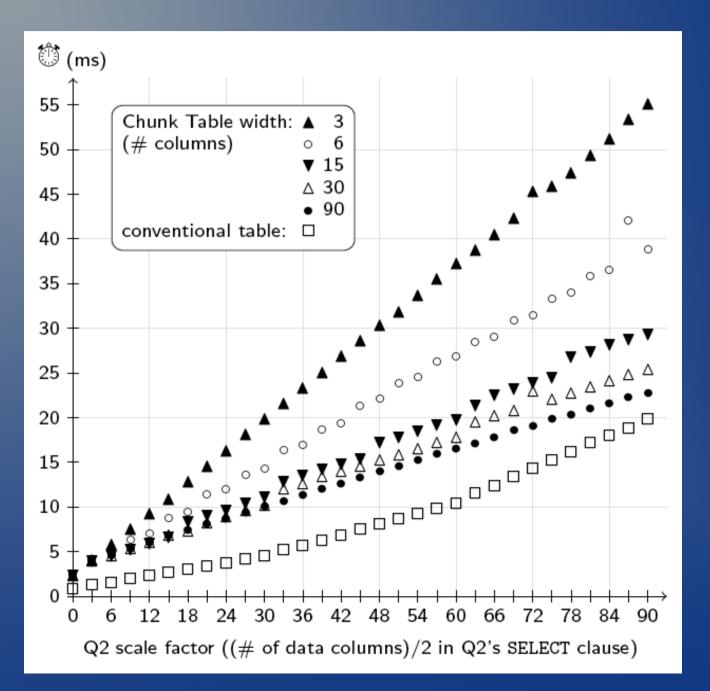
Conventional Schema

Parent id coll col2 ... col90 Child <u>id Parent_col1_co</u>l2...col90

Chunk Schema

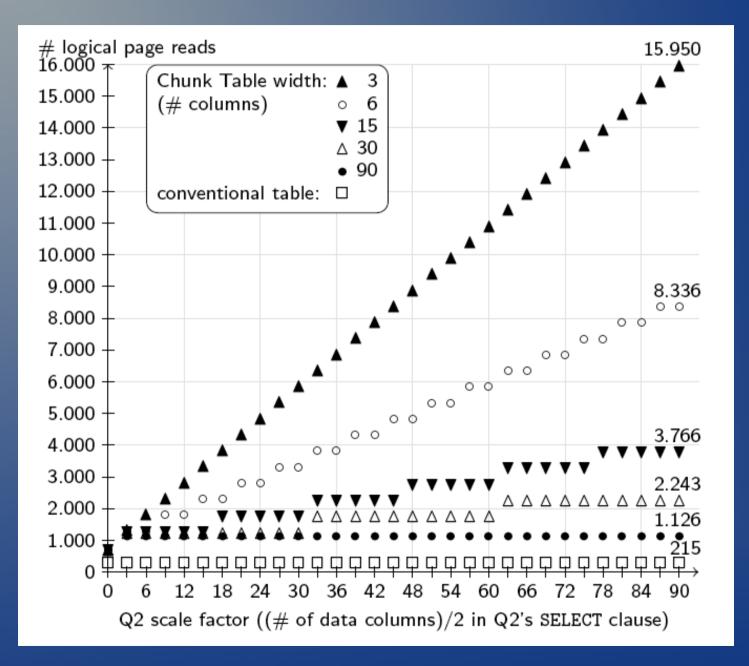
- $Chunk_{Data}$
- table chunk row int1 int2 int2 date date2
 str1 str2
- $Chunk_{Index}$
- table chunk row int1

Join Overhead Costs



17

Meta-Data Costs



Discussion

- Strengths
 - Chunk tables is a good design for trade-of extensibility and meta data usage.
 - Chunk tables gives response time improvement over vertical partitioning
- ShortComings/Future work
 - No Algorithms to design chunk tables
 - Identifying the chunks is heuristic
 - No comparative experiment done with the other schema mapping techniques proposed in the paper

Conclusion

- Is chunk tables a good approach for designing multi-tenant databases?
- How practical it is for real life systems ?
- How do companies like Salesforce.com handle it ?



QUESTIONS ?