

Computational Geometry and Geometric Computing Eric Berberich Kurt Mehlhorn Michael Sagraloff Winter 2009/2010 Discussion on December 2nd.

Exercise 6

## Motivation

We practise arrangements.

## Bounding box

Give an  $O(n \log n)$  algorithm to compute an axis-parallel rectangle R that contains all vertices in an arrangement of n lines.

## Zone of segment in triangulation

Given a triangulation  $\mathcal{T}$  with *n* triangles and a segment  $s := \overline{pq}$ . Compute all triangles intersected by *s*. Especially take care about degenerate situations.

## Trapezoidal decomposition

The *trapezoidal decomposition* of an arrangement  $\mathcal{A}$  induced by a set of curves is given by drawing vertical extensions from each vertex in upward and downward direction. Such an extension is either a segment if it hits another curve, or a ray that extends to infinity.

- 1. Sketch an algorithm to compute this decomposition.
- 2. Assume that the arrangement consists of  $n_e$  edges and they are in general position. Use a sweep argument to show that the vertical decomposition consists of  $3n_e + 1$  trapezoids.
- 3. Relax the conditions for a trapezoid and show that the upper bound of  $3n_e+1$  trapezoids still holds for edges not being in general position.

Have fun with the solution!