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## Exercise 6

## Motivation

We practise arrangements.

## Bounding box

Give an $O(n \log n)$ algorithm to compute an axis-parallel rectangle $R$ that contains all vertices in an arrangement of $n$ lines.

## Zone of segment in triangulation

Given a triangulation $\mathcal{T}$ with $n$ triangles and a segment $s:=\overline{p q}$. Compute all triangles intersected by $s$. Especially take care about degenerate situations.

## Trapezoidal decomposition

The trapezoidal decomposition of an arrangement $\mathcal{A}$ induced by a set of curves is given by drawing vertical extensions from each vertex in upward and downward direction. Such an extension is either a segment if it hits another curve, or a ray that extends to infinity.

1. Sketch an algorithm to compute this decomposition.
2. Assume that the arrangement consists of $n_{e}$ edges and they are in general position. Use a sweep argument to show that the vertical decomposition consists of $3 n_{e}+1$ trapezoids.
3. Relax the conditions for a trapezoid and show that the upper bound of $3 n_{e}+1$ trapezoids still holds for edges not being in general position.

Have fun with the solution!

