

Excercises Computational Geometry

<http://www.mpi-inf.mpg.de/departments/d1/teaching/ws13/ComputationalGeometry/>

Sheet 11

Deadline: 14.01.2014, 10:00am

Rules: Until the end of the semester you have to reach 50% of the achievable points to be admitted to the exam. 40 points correspond to 100%; you can get up to 10 bonus points.

Exercise 1 (10 pts)

What is the largest representable number as a `double`, the smallest positive number, the smallest normalized positive number?

Exercise 2 (10 pts)

Let $a, b \in \mathbb{F}$ with $\frac{1}{2} \leq \frac{a}{b} \leq 2$. Show that $a \ominus b = a - b$. This was first observed by Sterbenz.

Exercise 3 (10 pts)

Assume that point coordinates are `doubles` in $[\frac{1}{2}, 1]$. Show for three given points p, q, r

- $\text{orientation}(p, q, r) = 0$ implies $\text{float_orientation}(p, q, r) = 0$
- $\text{float_orientation}(p, q, r) \neq 0$ implies $\text{float_orientation}(p, q, r) = \text{orientation}(p, q, r)$
- What does this mean for the geometry of `float_orientation`?
- Can you find examples that make the floating point implementation of the convex hull algorithm crash when point coordinates are restricted to `doubles` in $[\frac{1}{2}, 1]$?

Exercise 4 (10 pts)

Given real numbers x, y, z . Derive (floating point) error bounds for the decision if the expression $xy - z$ is larger or smaller than zero.

Exercise 5 (10 pts)

Analyse the computational overhead of guarding predicates with (a) forward error analysis and (b) interval arithmetic.

Exercise 6 (10pts) Simulation of Simplicity: Argue why $\varepsilon(i, j) = \varepsilon^{i \cdot \delta + j}$ and $\varepsilon(i, j) = (i \cdot \delta + j) \cdot \varepsilon$ are no suitable choices for successful simulated perturbations of a matrix predicate.