

# RDF Engines

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## Introduction

Resource Description Framework  
SPARQL

## Triple Storages

Giant Triple Table  
Property Tables  
Vertically Partitioned Table  
Hexastore

## Experimental Evaluation

## Conclusion

# Resource Description Framework

RDF is an acronym for **R**esource **D**escription **F**ramework

- ▶ RDF is a data model to represent meta data.
- ▶ RDF was invented and polished by the W3C in '98.
- ▶ RDF was revisited in 2004.

further informations <http://www.w3.org/RDF/>

Nowadays there is a lot of RDF Data in the world wide web.  
For instance:

- ▶ Wikipedia
- ▶ Sience Life
- ▶ Social Networks

# Resource Description Framework

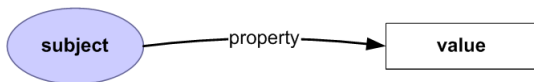
But why is the RDF data model that important?

- ▶ RDF allows to store and describe meta data in a standardized way.
- ▶ The simplicity allows to describe all kinds of data.
- ▶ Important for Semantic Web and Web 2.0

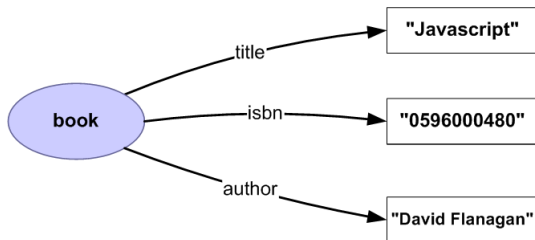
# Structuring RDF

## RDF structure

- ▶ In RDF data is represented as triples.
- ▶ The structure of such a triple is (subject,property,value).



# Example



But how to query data from such a RDF ??

# Query Language SPARQL

## What is SPARQL ?

- ▶ SPARQL is a standart language to query RDF data.
- ▶ SPARQL was recommended in the year 2008 by W3C.

*How does a SPARQL query look like ??*



# How does a SPARQL query look like ?

A simple generic SPARQL query

```
Select ?var1 ?var2 ....  
      where { pattern1. pattern2. .... }
```

# SPARQL Patterns

Triple Patterns	
p1	(s,p,o)
p2	(?s,p,o)
p3	(s,?p,o)
p4	(s,p,?o)
p5	(?s,?p,o)
p6	(s,?p,?o)
p7	(?s,p,?o)
p8	(?s,?p,?o)

Variables are characterized with questionmarks.

Variables are bound such that the triple is in the data set.

## A simple query:

```
select ?i  
  where { (book,isbn,?i). }
```

This query yields the isbn number of the book 'book' in our graph example.

Thus the result is the number 0596000480.

## Joins:

```
select ...  
  where { (?a,b,c).( ?a,e,f). ... }
```

# Triple Storages

RDF Graph is represented by set of all triples  $(s,p,o)$ . Often the triples are stored in a relational Database.

The Problem: How to store these triples in such a way that queries can be answered fast?

# Some approaches

## Four different approaches

1. Giant Triple Storage
2. Property Tables
3. Vertically Partition Tables
4. Hexastore

## Giant Triple storage

### Basic idea

- ▶ Just store all triples in a single table.
- ▶ Important for the performance is the choice of the right indexes, clustered and unclustered.

# LUBM like example

## LUBM Example:

Subj	Property	Obj
ID1	type	FullProfessor
ID1	teacherOf	'AI'
ID1	bachelorFrom	'MIT'
ID1	mastersFrom	'Cambridge'
ID1	phdFrom	'Yale'
ID2	type	AssocProfessor
ID2	worksFor	'MIT'
ID2	teacherOf	'DataBases'
ID2	bachelorsFrom	'Yale'
ID2	phdFrom	'Stanford'
ID3	type	GradStudent
ID3	advisor	ID2
ID3	teachingAssist	'AI'
ID3	bachelorsFrom	'Stanford'
ID3	mastersFrom	'Princeton'
ID4	type	GradStudent
ID4	advisor	ID1
ID4	takesCourse	'DataBases'
ID4	bachelorsFrom	'Columbia'



# Pros and Cons

## Pros

- ▶ Easy to implement
- ▶ It's suitable for a large number of properties
  - ▶ If the right indexes has been chosen

## Cons

- ▶ Many self joins needed

# Some approaches

## Four different approaches

1. Giant Triple Storage
2. Property Tables
3. Vertically Partition Tables
4. Hexastore

## What is the idea?

- ▶ Get relational Schemata out of the RDF data

## How can this be achieved?

- ▶ Collect several properties in one table

## A basic example

Professors						
ID	type	teacherOf	bachelorsFrom	mastersFrom	phdFrom	worksFor
ID1	FullProfessor	'AI'	'MIT'	'Cambridge'	'Yale'	NULL
ID2	AssocProfessor	'DataBases'	'Yale'	NULL	'Stanford'	'MIT'

Students						
ID	type	advisor	bachelorsFrom	mastersFrom	teachingAssist	takesCourse
ID3	GradStudent	ID2	'Stanford'	'Princeton'	'AI'	NULL
ID4	GradStudent	ID1	'Columbia'	NULL	NULL	'DataBases'

# Pros and Cons

## Pros

- ▶ If the data is structured then we will have nice relational tables.

## Cons

- ▶ imposes structure on semistructured data
- ▶ this yields NULL entries
- ▶ the selection of properties to store in one table is non trivial

# Some approaches

## Four different approaches

1. Giant Triple Storage
2. Property Tables
3. Vertically Partition Tables
4. Hexastore

# Vertical Partitioned Table

The basic idea of the vertical partitioned table approach:

- ▶ Generate for every distinct property an own two column table.

### type

ID1	FullProfessor
ID2	AssocProfessor
ID3	GradStudent
ID4	GradStudent

### teacherOf

ID1	'AI'
ID2	'DataBases'

### bachelorsFrom

ID1	'MIT'
ID2	'Yale'
ID3	'Stanford'
ID4	'Columbia'

### mastersFrom

ID1	'Cambridge'
ID3	'Princeton'

### phdFrom

ID1	'Yale'
ID2	'Stanford'

### worksFor

ID2	'MIT'
-----	-------

### advisor

ID3	ID2
ID4	ID1

### bachelorsFrom

ID1	'MIT'
ID2	'Yale'
ID3	'Stanford'
ID4	'Columbia'

### teachingAssist

ID3	'AI'
-----	------

### takesCourse

ID4	'DataBases'
-----	-------------



# Pros and Cons

## Pros

- ▶ Excellent performance
  - ▶ Small number of properties
  - ▶ Queries are bound on just a few properties

## Cons

- ▶ Bad performance
  - ▶ When a property of a query is not bound
  - ▶ When a property is bound in the running time
  - ▶ When there exists a huge number of properties

# Some approaches

## Four different approaches

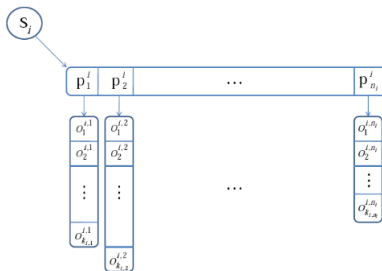
1. Giant Triple Storage
2. Property Tables
3. Vertically Partitioned Tables
4. **Hexastore**

## The basic idea of **Hexastore**

- ▶ An index for each permutation is materialized. In the Hexastore solution indexes for all orderings are generated: Namely **spo, pos, osp, sop, pso, ops**
- ▶ This index should give you a fast response time.

# The hexastore index structure

How does such an index look like in the hexastore?



## An example

The **spo** index :

Subject key  $S_i$  points to sorted vector of  $n_i$  property keys,  
 $\{p_1^i, p_2^i, \dots, p_{n_i}^i\}$ .

Each property key  $p_j^i$  is linked to a sorted list of object keys.

The object key lists are shared with the **ps** index.

# Pros and Cons

## Pros

- ▶ Fast Merge Joins in the first steps of any query.

## Cons

- ▶ Up to five fold memory usage.
- ▶ Optimized for memory based storage.
  - ▶ might have bad performance on disk based storage.

## Experimental Evaluation - Motivation

We have seen four approaches,  
but which is the best one?

# Experimental Evaluation - Introduction

How are we going to test these approaches?

1. We will build some reference implementation of the different approaches using state of the art techniques.
2. The reference implementations will be tested with real data sets or synthetic data sets.
  - ▶ LUBM - synthetic data set
  - ▶ Barton - real data set
3. We will query a small set of queries to the data sets.
4. Benchmark the response time.



# First Experiment - Triple Store vs Vertically Partitioned

# First Experiment - Triple Store vs Vertically Partitioned

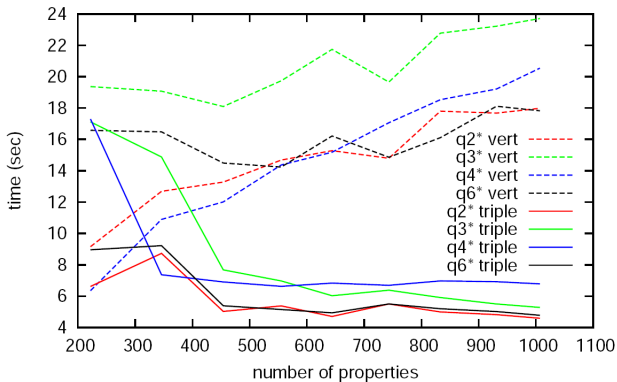
	store	cluster	time	q1	q2	q2*	q3	q3*	q4	q4*	q5	q6	q6*	q7	q8	G	G*	$\frac{G^*}{G}$
DBX	triple	SPO	real	4.29	42.61	93.11	34.86	97.92	8.02	6.12	11.70	89.11	142.10	1.34	14.47	13.2	21.1	1.6
			user	4.29	33.31	68.88	34.16	95.11	8.02	6.10	11.68	74.96	120.36	1.27	10.58	12.3	19.0	1.5
		PSO	real	1.72	40.18	38.35	45.65	67.32	3.22	2.49	10.61	57.52	63.04	1.42	12.14	9.8	13.6	1.4
			user	1.72	40.17	38.35	45.64	66.85	3.22	2.47	10.60	57.33	63.03	1.34	8.02	9.7	13.1	1.4
	vert.	SO	real	1.55	39.62	74.85	45.17	94.59	6.12	14.18	5.69	45.57	154.81	1.25	11.55	9.1	17.7	1.9
			user	1.55	39.61	74.83	45.16	94.09	6.12	14.15	5.67	45.56	153.08	1.18	7.49	9.1	17.0	1.9
MonetDB	triple	SPO	real	1.53	3.50	3.63	5.28	17.54	1.68	1.98	2.77	8.37	7.33	1.82	4.76	2.9	3.7	1.3
			user	1.36	2.73	2.91	4.33	15.40	1.41	1.65	2.30	6.20	5.70	1.65	3.75	2.4	3.1	1.3
		PSO	real	0.78	2.80	2.83	4.36	12.59	1.70	1.97	1.44	5.67	4.59	0.18	5.23	1.5	2.4	1.6
			user	0.69	2.31	2.31	3.69	10.54	1.59	1.86	1.16	3.80	3.65	0.17	3.60	1.3	2.0	1.5
	vert.	SO	real	0.79	1.50	5.50	2.64	14.01	0.50	2.57	1.29	4.65	11.51	0.06	5.05	0.9	2.0	2.2
			user	0.68	1.44	5.20	2.52	13.25	0.48	2.40	1.03	4.40	11.23	0.06	4.20	0.8	1.9	2.4
C-Store	vert.	SO	real	0.59	1.35	-	4.08	-	1.52	-	12.95	12.04	-	0.77	-	2.4	-	-
			user	0.49	1.17	-	3.45	-	1.28	-	11.67	10.49	-	0.34	-	1.9	-	-

## First Experiment - Triple Store vs Vertically Partitioned cont'd

The vertical approach outperforms the triple storage on almost all queries without a star.

What happens if the number of properties is further increased?

# First Experiment - Triple Store vs Vertically Partitioned cont'd

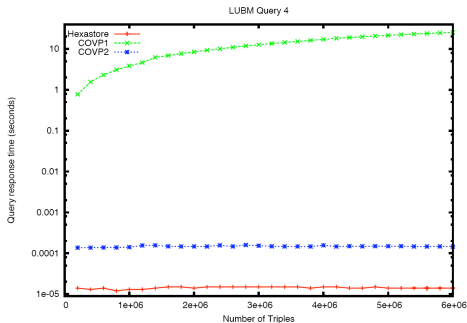


# Second Experiment - Hexastore vs Vertically Partitioned

## Second Experiment - Hexastore vs Vertically Partitioned

### Results:

- ▶ Hexastore outperforms vert part approach.



## Second Experiment - Hexastore vs Vertically Partitioned

But:

- ▶ All triples were loaded in the main memory beforehand

## Not all swans are white!

Every storage scheme has its weaknesses. The triple store with elaborated usage of the “right“ indexes seems to have the best performance regarding scalability.



Thanks for your attention!!

# References